H8S/2623 Series, H8S/2623F-ZTATTM

Overview

HITACHI

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Preface

Hitachi's H8S family of single-chip microcomputers comprises a number of new series offering the high performance and low power consumption of the existing H8 Series, which is widely used for machine control, etc., together with significantly greater ease of use.

The H8S/2000 Series features CPU object-level compatibility with the H8/300H Series, H8/300 Series, and H8/300L Series within the H8 Series.

Series	Features	
H8S/2000	Upward-compatible with the H8/300H Series and H8/300 Series; twice the performance at the same frequency	
H8/300H	16-Mbyte linear address space; upward-compatible with the H8/300 Series; condinstruction set; powerful word-size and longword-size arithmetic instructions	
H8/300	64-kbyte address space; general register system; concise instruction set; powerf bit manipulation instructions	
H8/300L	Same CPU as the H8/300 Series; consumer application oriented supporting modules; low voltage, low power consumption	

This manual gives an overview of the H8S/2623 Series of products for single-chip applications within the H8S Series.

Intended Readership: This Overview is intended for readers who have a basic understanding of microcomputers, and are looking for information on the features and functions of the H8S/2623 Series. Readers undertaking system design using these products, or requiring more detailed information on their use, should refer to the relevant Hardware Manuals and the H8S/2600 and H8S/2000 Series Programming Manual.

Related Documents

Contents	Title	Document No.
H8S/2623 Series hardware	H8S/2623 Series, H8S/2623F-ZTAT™ Hardware Manual	ADE-602-164
H8S/2000 Series execution instructions	H8S/2600 Series and H8S/2000 Series Programming Manual	ADE-602-083A

Note: F-ZTAT (Flexible-ZTAT) is a trademark of Hitachi, Ltd.

The product specifications in this Overview are subject to change without notice. The relevant Hardware Manual must be used when undertaking product design.

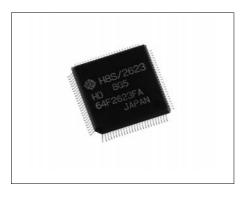
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Section 1 Overview of H8S/2623 Series

1.1 Features of H8S/2623 Series



H8S/2623 Series microcomputers are designed for faster instruction execution, using a realtime control oriented CPU with an internal 32-bit architecture, and can run programs based on the C high-level language efficiently. As well as large-capacity ROM and RAM, these microcomputers include comprehensive on-chip supporting modules needed for control systems, simplifying the implementation of sophisticated, high-performance systems.

High-performance H8S/2600 CPU

- · General-register architecture
 - Sixteen 16-bit general registers (also usable as sixteen 8-bit registers or eight 32-bit registers)
- High-speed operation suitable for realtime control
 - 20 MHz maximum operating frequency ($V_{CC} = 3.3 \text{ V} \pm 0.3 \text{ V}$, $PV_{CC} = 5 \text{ V} \pm 10\%$)
 - High-speed arithmetic operations
 - 8/16/32-bit register-register add/subtract: 50 ns
 - 16×16 -bit register-register multiply: 200 ns
 - $16 \times 16 + 42$ -bit multiply-and-accumulate: 200 ns
 - 32 ÷ 16-bit register-register divide: 1000 ns
- Instruction set suitable for high-speed operation
 - Sixty-nine types of basic instructions
 - 8/16/32-bit transfer instructions
 - Unsigned/signed multiply and divide instructions
 - Multiply-and-accumulate instruction
 - Powerful bit-manipulation instructions
- CPU operating mode
 - Advanced mode: Maximum 16-Mbyte address space

256-kbyte flash memory or 256-/128-/64-kbyte mask ROM* on-chip

Note: * In planning stage.

On-chip high-speed static RAM

12/8/4 kbytes

Bus controller on-chip

- Address space divided into 8 areas, with bus specifications settable independently for each area
- Chip select output possible for areas 0 to 7
- Selection of 8-bit or 16-bit access space for each area
- 2-state or 3-state access space can be set for each area
- Number of program wait states can be set for each area
- Burst ROM directly connectable
- External bus release function

PC break controller

- Supports debugging functions by means of PC break interrupts
- Two break channels

Data transfer controller (DTC)

- Activated by internal interrupt or software
- Multiple transfers or multiple types of transfer possible for one activation source
- Transfer possible in repeat mode, block transfer mode, etc.
- Request can be sent to CPU for interrupt that activated DTC

16-bit timer-pulse unit (TPU)

- Six-channel 16-bit timer on-chip
- Pulse I/O processing capability for up to 16 pins
- Automatic 2-phase encoder count capability

$Programmable \ pulse \ generator \ (PPG)$

- Maximum 8-bit pulse output possible with TPU as time base
- Output trigger selectable in 4-bit groups
- Non-overlap interval can be set
- Direct output or inverse output setting possible

One watchdog timer (WDT) channel

• Watchdog timer or interval timer function selectable

Three serial communication interface (SCI) channels (SCI0 to SCI2)

- Asynchronous mode or synchronous mode selectable
- Multiprocessor communication function
- Smart card interface function

One Hitachi controller area network (HCAN) channel

- CAN: Ver. 2.0B compatibility
- Buffers: 15 transmit/receive buffers, one transmit-only buffer
- Receive message filtering capability

A/D converter

- Resolution: 10 bits
- Input: 16 channels
- 13.3 µs minimum conversion time (20 MHz operation)
- Single or scan mode selectable
- Sample-and-hold function
- A/D conversion can be activated by external trigger or timer trigger

I/O ports

- 54 I/O pins
- 16 input pins

Interrupt controller

- Seven external interrupt pins (NMI, $\overline{IRQ0}$ to $\overline{IRQ5}$)
- 47 internal interrupt sources
- Eight priority levels settable

Power-down state

- Medium-speed mode
- Sleep mode
- Module stop mode
- Software standby mode
- Hardware standby mode

Four MCU operating modes

				Exterr	nal Data Bus
Mode	CPU Operating Mode	Description	On-Chip ROM	Initial Value	Maximum Value
4	Advanced	Expanded mode with on- chip ROM disabled	Disabled	16 bits	16 bits
5		Expanded mode with on- chip ROM disabled	Disabled	8 bits	16 bits
6		Expanded mode with on- chip ROM enabled	Enabled	8 bits	16 bits
7		Single-chip mode	Enabled	_	_

Clock pulse generator

- Built-in PLL circuit (×1, ×2, ×4)
- Input clock frequency: 2 to 20 MHz

Package

• 100-pin plastic QFP (FP-100B)

Product lineup (preliminary)

Product Name

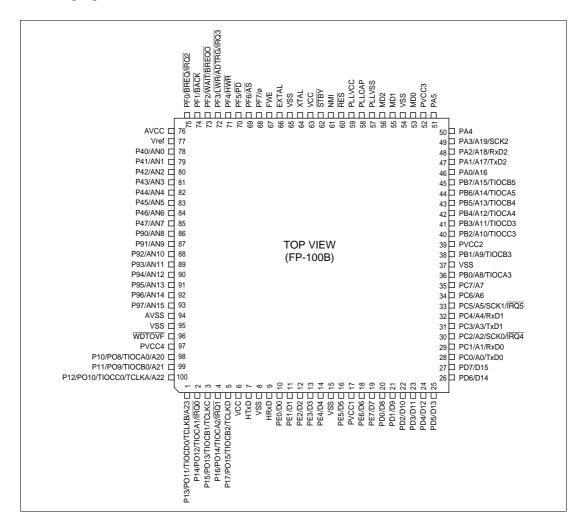
Mask ROM Version	F-ZTAT™ Version	ROM/RAM (Bytes)	Packages
HD6432623	HD64F2623	256 k/12 k	FP-100B
HD6432622	_	128 k/8 k	
HD6432621	_	64 k/4 k	

Note: F-ZTAT™ is a trademark of Hitachi, Ltd.

1.2 Pin Arrangement and Functions

H8S/2623 Series Pin Arrangement

• 100-pin plastic QFP (FP-100B)



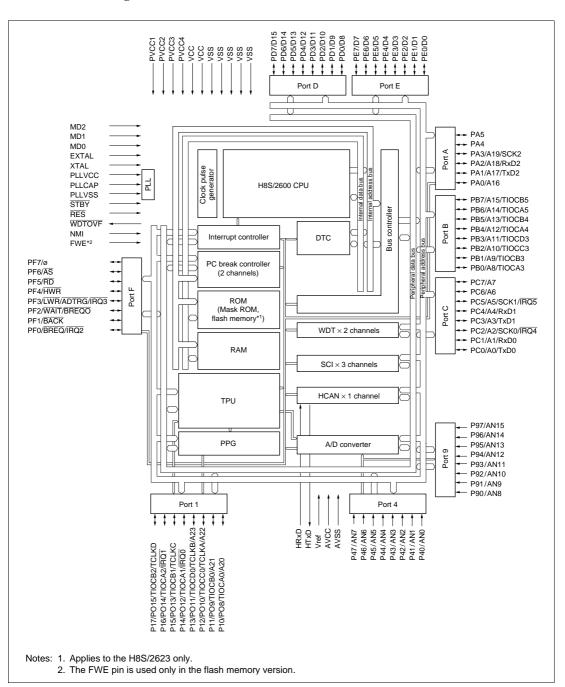
Pin Functions

Туре	Symbol	I/O	Name and Function
Power	VCC	Input	Power supply
	PVCC1	Input	Port power supply
	PVCC2	Input	Port power supply
	PVCC3	Input	Port power supply
	PVCC4	Input	Port power supply
	VSS	Input	Ground
Clock	PLLVCC	Input	PLL power supply
	PLLVSS	Input	PLL ground
	PLLCAP	Input	PLL capacitance
	XTAL	Input	Crystal
	EXTAL	Input	External clock
	Ø	Output	System clock
Operating mode control	MD2 to MD0	Input	Mode pins
System control	RES	Input	Reset input
	STBY	Input	Standby
	BREQ	Input	Bus request
	BREQO	Output	Bus request output
	BACK	Output	Bus request acknowledge
	FWE	Input	Flash write enable
Interrupt	NMI	Input	Nonmaskable interrupt
	IRQ5 to IRQ0	Input	Interrupt request 5 to 0
Address bus	A23 to A0	Output	Address bus
Data bus	D15 to D0	I/O	Data bus
Bus control	ĀS	Output	Address strobe
	RD	Output	Read
	HWR	Output	High write
	LWR	Output	Low write
	WAIT	Input	Wait

Type	Symbol	I/O	Name and Function
16-bit timer-pulse	TCLKA to TCLKD	Input	Clock input A to D
unit (TPU)	TIOCA0, TIOCB0, TIOCC0, TIOCD0	I/O	Input capture/output compare match A0 to D0
	TIOCA1, TIOCB1	I/O	Input capture/output compare match A1 and B1
	TIOCA2, TIOCB2	I/O	Input capture/output compare match A2 and B2
	TIOCA3, TIOCB3, TIOCC3, TIOCD3	I/O	Input capture/output compare match A3 to D3
	TIOCA4, TIOCB4	I/O	Input capture/output compare match A4 and B4
	TIOCA5, TIOCB5	I/O	Input capture/output compare match A5 and B5
Programmable PO15 to PO8 pulse generator (PPG)		Output	Pulse output 15 to 8
Watchdog timer (WDT)	WDTOVF	Output	Watchdog timer overflow
Serial communication	TxD2, TxD1, TxD0	Output	Transmit data
interface (SCI)/ Smart card	RxD2, RxD1, RxD0	Input	Receive data
interface	SCK2, SCK1, SCK0	I/O	Serial clock
HCAN	HTxD	Output	HCAN transmit data pin
	HRxD	Input	HCAN receive data pin
A/D converter	AN15 to AN0	Input	Analog input
	ADTRG	Input	A/D conversion external trigger input
	AVCC	Input	Analog power supply
	AVSS	Input	Analog ground
	Vref	Input	Analog reference voltage
I/O ports	P17 to P10	I/O	Port 1
	P47 to P40	Input	Port 4
	P97 to P90	Input	Port 9
	PA5 to PA0	I/O	Port A
	PB7 to PB0	I/O	Port B
	PC7 to PC0	I/O	Port C
	PD7 to PD0	I/O	Port D
	PE7 to PE0	I/O	Port E
	PF7 to PF0	I/O	Port F

1.3 Block Diagram

Internal Block Diagram of H8S/2623 Series



Section 2 CPU

2.1 Features

The H8S/2600 CPU is a high-speed central processing unit with an internal 32-bit architecture that is upward-compatible with the H8/300 and H8/300H CPUs. The H8S/2600 CPU has sixteen 16-bit general registers, can address a 16-Mbyte (architecturally 4-Gbyte) linear address space, and is ideal for realtime control.

Features

- Upward-compatible with H8/300 and H8/300H CPUs
 - Can execute H8/300 and H8/300H object programs
- General-register architecture
 - Sixteen 16-bit general registers (also usable as sixteen 8-bit registers or eight 32-bit registers)
- Sixty-nine basic instructions
 - 8/16/32-bit arithmetic and logic instructions
 - Multiply and divide instructions
 - Powerful bit-manipulation instructions
 - Multiply-and-accumulate instruction
- Eight addressing modes
 - Register direct [Rn]
 - Register indirect [@ERn]
 - Register indirect with displacement [@(d:16,ERn) or @(d:32,ERn)]
 - Register indirect with post-increment or pre-decrement [@ERn+ or @-ERn]
 - Absolute address [@aa:8, @aa:16, @aa:24, or @aa:32]
 - Immediate [#xx:8, #xx:16, or #xx:32]
 - Program-counter relative [@(d:8,PC) or @(d:16,PC)]
 - Memory indirect [@@aa:8]
- 16-Mbyte address space
 - Program: 16 Mbytes
 - Data: 16 Mbytes (4 Gbytes architecturally)

- High-speed operation
 - All frequently used instructions execute in one to two states
 - Maximum operating frequency: 20 MHz
 - 8/16/32-bit register-register add/subtract: 50 ns
 - 8×8 -bit register-register multiply: 150 ns
 - $16 \div 8$ -bit register-register divide: 600 ns
 - 16×16 -bit register-register multiply: 200 ns
 - 32 ÷ 16-bit register-register divide: 1000 ns
- CPU operating mode
 - Advanced mode
- Power-down state
 - Transition to power-down state by SLEEP instruction
 - CPU operating clock can be selected

Differences between H8S/2600 CPU and H8S/2000 CPU: The differences between the H8S/2600 CPU and the H8S/2000 CPU are as shown below.

• Register configuration

The MAC register is supported only by the H8S/2600 CPU.

• Basic instructions

The four instructions MAC, CLRMAC, LDMAC, and STMAC are supported only by the H8S/2600 CPU.

Number of execution states

The number of execution states of the MULXU and MULXS instructions is different in each CPU.

Execution States	

Instruction	Mnemonic	H8S/2600	H8S/2000
MULXU	MULXU.B Rs, Rd	3	12
	MULXU.W Rs, ERd	4	20
MULXS	MULXS.B Rs, Rd	4	13
	MULXS.W Rs, ERd	5	21

In addition, there are differences in address space, CCR and EXR register functions, power-down modes, etc., depending on the model.

Differences from H8/300 CPU: In comparison to the H8/300 CPU, the H8S/2600 CPU has the following enhancements.

• More general registers and control registers

- Eight 16-bit extended registers, and one 8-bit control register, have been added.
- Address space
 - The 16-Mbyte address space can be used effectively.
- Enhanced instructions
 - Addressing modes of bit-manipulation instructions have been enhanced.
 - Signed multiply and divide instructions have been added.
 - A multiply-and-accumulate instruction has been added.
 - Two-bit shift instructions have been added.
 - Instructions for saving and restoring multiple registers have been added.
 - A test and set instruction has been added.
- · Higher speed
 - Basic instructions execute twice as fast.

Differences from H8/300H CPU: In comparison to the H8/300H CPU, the H8S/2600 CPU has the following enhancements.

- · Additional control registers
 - One 8-bit and two 32-bit control registers have been added.
- Enhanced instructions
 - Addressing modes of bit-manipulation instructions have been enhanced.
 - A multiply-and-accumulate instruction has been added.
 - Two-bit shift instructions have been added.
 - Instructions for saving and restoring multiple registers have been added.
 - A test and set instruction has been added.
- Higher speed
 - Basic instructions execute twice as fast.

2.2 Register Configuration

The H8S/2600 CPU has general registers and control registers. The eight 32-bit general registers all have identical functions and can be used as either address registers or data registers. The control registers are the 24-bit program counter (PC), 8-bit extend register (EXR), 8-bit condition code register (CCR), and 64-bit multiply-and-accumulate register (MAC).

CPU Internal Registers

	15	07	07	0
ER0	E0		R0H	R0L
ER1	E1		R1H	R1L
ER2	E2		R2H	R2L
ER3	E3		R3H	R3L
ER4	E4		R4H	R4L
ER5	E5		R5H	R5L
ER6	E6		R6H	R6L
ER7 (SF	P) E7		R7H	R7L
				7 6 5 4 3 2 1 0
			0/	7 6 5 4 3 2 1 0 CR I UI H U N Z V C
	63		4	
MA	C Sign e	xtension	<u> </u>	MACH
	31	MAC	JL	0
	31			0
Legend:				
SP:	Stack pointer	H:	Half-carry flag	
PC: Program counter		U: N:	User bit	
	Extend register Trace bit	N: Z:	Negative flag Zero flag	
EXR:			Overflow flag	
EXR: T:	Interrupt mask bits	V:		
EXR: T: I2 to I0:	Interrupt mask bits Condition-code register	C:	Carry flag	
EXR: T:		C: MAC:	Carry flag Multiply-and-acc	umulate register

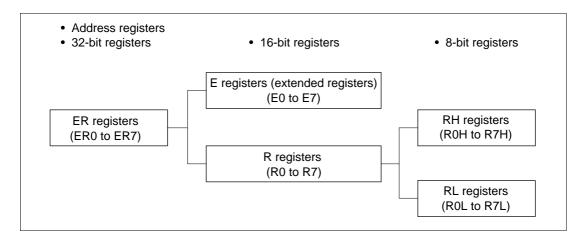
General Registers: The CPU has eight 32-bit general registers. These general registers are all functionally alike and can be used as both address registers and data registers. When a general register is used as a data register, it can be accessed as a 32-bit, 16-bit, or 8-bit register. When the general registers are used as 32-bit registers or address registers, they are designated by the letters ER (ER0 to ER7).

The ER registers divide into 16-bit general registers designated by the letters E (E0 to E7) and R (R0 to R7). These registers are functionally equivalent, providing a maximum sixteen 16-bit registers. The E registers (E0 to E7) are also referred to as extended registers.

The R registers divide into 8-bit general registers designated by the letters RH (R0H to R7H) and RL (R0L to R7L). These registers are functionally equivalent, providing a maximum sixteen 8-bit registers.

The figure below illustrates the usage of the general registers. The usage of each register can be selected independently.

Usage of General Registers



Control Registers: The control registers are the 24-bit program counter (PC), 8-bit extend register (EXR), 8-bit condition-code register (CCR), and 64-bit multiply-and-accumulate register (MAC).

• Program Counter (PC)

This 24-bit counter indicates the address of the next instruction the CPU will execute. The length of all CPU instructions is 2 bytes (one word), so the least significant PC bit is ignored. (When an instruction code is read, the least significant bit is regarded as 0.)

• Extend Register (EXR)

This 8-bit register comprises a trace bit (T) and interrupt mask bits (I2 to I0).

— Bit 7—Trace Bit (T)

Specifies whether or not trace mode is set. When this bit is cleared to 0, instructions are executed sequentially. When set to 1, trace exception handling is started each time an instruction is executed.

— Bits 6 to 3—Reserved

These bits are always read as 1.

— Bits 2 to 0—Interrupt Mask Bits (I2 to I0)

These bits specify the interrupt request mask level (0 to 7). See section 2.9, Interrupts, for details.

EXR can be manipulated by the LDC, STC, ANDC, ORC, and XORC instructions. Except in the case of STC, interrupts (including NMI) are not accepted for 3 states after the instruction is executed.

• Condition-Code Register (CCR)

This 8-bit register contains internal CPU status information, including an interrupt mask bit (I) and half-carry (H), negative (N), zero (Z), overflow (V), and carry (C) flags.

— Bit 7—Interrupt Mask Bit (I)

Masks interrupts other than NMI when set to 1. (NMI is accepted regardless of the I bit setting.) The I bit is set to 1 when exception handling execution is started. For details, refer to section 2.9, Interrupts.

— Bit 6—User Bit or Interrupt Mask Bit (UI)

Can be written and read by software using the LDC, STC, ANDC, ORC, and XORC instructions. In this series, this bit cannot be used as an interrupt mask bit.

— Bit 5—Half-Carry Flag (H)

When the ADD.B, ADDX.B, SUB.B, SUBX.B, CMP.B, or NEG.B instruction is executed, this flag is set to 1 if there is a carry or borrow at bit 3, and cleared to 0 otherwise. When the ADD.W, SUB.W, CMP.W, or NEG.W instruction is executed, the H flag is set to 1 if there is a carry or borrow at bit 11, and cleared to 0 otherwise. When the ADD.L, SUB.L, CMP.L, or NEG.L instruction is executed, the H flag is set to 1 if there is a carry or borrow at bit 27, and cleared to 0 otherwise.

— Bit 4—User Bit (U)

Can be written and read by software using the LDC, STC, ANDC, ORC, and XORC instructions.

— Bit 3—Negative Flag (N)

Stores the value of the most significant bit (sign bit) of data.

— Bit 2—Zero Flag (Z)

Set to 1 to indicate zero data, and cleared to 0 to indicate non-zero data.

— Bit 1—Overflow Flag (V)

Set to 1 when an arithmetic overflow occurs, and cleared to 0 at other times.

— Bit 0—Carry Flag (C)

Set to 1 when a carry occurs, and cleared to 0 otherwise. Used by:

- Add instructions, to indicate a carry
- Subtract instructions, to indicate a borrow
- Shift and rotate instructions, to store the carry

The carry flag is also used as a bit accumulator by bit-manipulation instructions.

• Multiply-and-Accumulate Register (MAC)

This 64-bit register stores the result of a multiply-and-accumulate operation. It consists of two 32-bit registers, MACH and MACL. Only the lower 10 bits of MACH are valid; the upper bits are sign-extended.

2.3 Data Formats

The CPU can process 1-bit, 4-bit (BCD), 8-bit (byte), 16-bit (word), and 32-bit (longword) data. Bit-manipulation instructions operate on 1-bit data by accessing bit n (n = 0, 1, 2, ..., 7) of byte operand data. The DAA and DAS decimal-adjust instructions treat byte data as two digits of 4-bit BCD data.

General Register Data Formats

Data Type	Register Number	Data Format
1-bit data	RnH	7 0 7 6 5 4 3 2 1 0 Don't care
1-bit data	RnL	7 0 Don't care 7 6 5 4 3 2 1 0
4-bit BCD data	RnH	7 4 3 0 Upper Lower Don't care
4-bit BCD data	RnL	7 4 3 0 Don't care Upper Lower
Byte data	RnH	7 0 Don't care MSB LSB
Byte data	RnL	7 0 Don't care MSB LSB

General Register Data Formats (cont)

Data Type	Register Number	Data Format
Word data	Rn	15 0 MSB LSB
Word data	En	WOD EOD
15 MSB	L	0 SB
Longword data	ERn	
31 MSB	En	16 15 0 Rn LSB
Legend: ERn: General re En: General re Rn: General re RnH: General re RnL: General re MSB: Most signi LSB: Least sign	egister E egister R egister RH egister RL ficant bit	

Memory Data Formats

Address 7 1-bit data Address L 7			_	_			
<u> </u>	,						
1-bit data Address L 7							0
	6	5 5	4	3	2	1	0
	-						1
Byte data Address L MS	SB¦	-					LSB
		:					1
Word data Address 2M MS	BB.	-					
Address 2M + 1							LSB
		-					<u> </u>
Longword data Address 2N MS	BB;	-					-
Address 2N + 1	-						1
Address 2N + 2							
Address 2N + 3		-					LSB

2.4 Addressing Modes

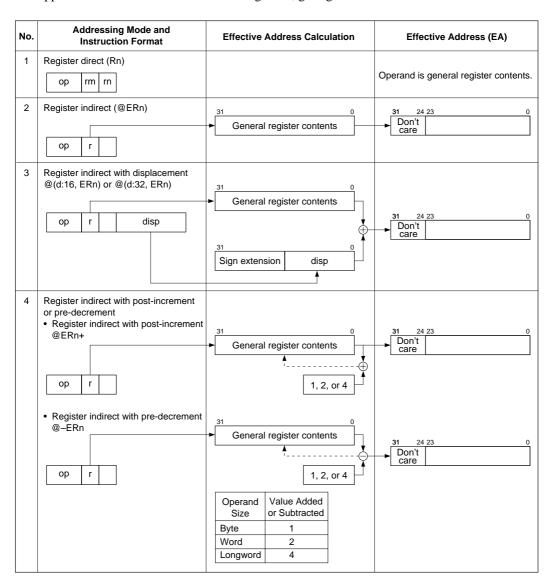
The H8S/2600 CPU supports eight addressing modes.

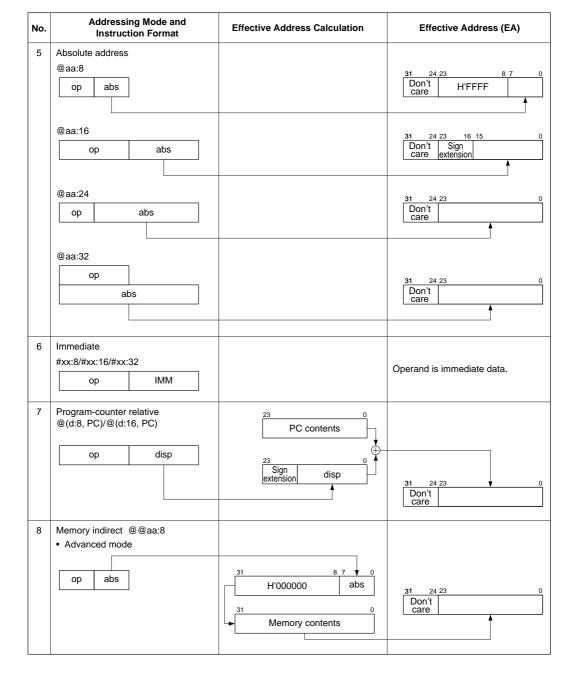
Addressing Modes

No.	Addressing Mode	Symbol
1	Register direct	Rn
2	Register indirect	@ERn
3	Register indirect with displacement	@(d:16,ERn)/@(d:32,ERn)
4	Register indirect with post-increment Register indirect with pre-decrement	@ERn+ @-ERn
5	Absolute address	@aa:8/@aa:16/@aa:24/@aa:32
6	Immediate	#xx:8/#xx:16/#xx:32
7	Program-counter relative	@(d:8,PC)/@(d:16,PC)
8	Memory indirect	@ @aa:8

Effective Address Calculation

The upper 8 bits of the effective address are ignored, giving a 16-bit address.





2.5 Instruction Set

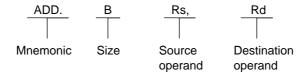
The H8S/2600 CPU has 69 types of instructions.

Features

- Upward-compatible at object level with H8/300H and H8/300 CPUs
- General register architecture
- 8/16/32-bit transfer instructions and arithmetic and logic instructions
 - Byte (B), word (W), and longword (L) formats for transfer instructions and basic arithmetic and logic instructions
- Unsigned and signed multiply and divide instructions
- Powerful bit-manipulation instructions
- · Instructions for saving and restoring multiple registers

Assembler Format

The ADD instruction format is shown below as an example.



• Data transfer instructions

			Ad	dres	sing		le/In 3yte		ction	Len	gth								
	Mnemonic	Operand Size	#xx	Rn	@ERn	@(d,ERn)	@-ERn/@ERn+	@ aa	@(d,PC)	@ @aa	ı	Operation	С	on H				ode / C	of States*
MOV	MOV.B #xx:8,Rd	В	2									#xx:8→Rd8	+		1	t			- 1
	MOV.B Rs,Rd	В		2								Rs8→Rd8	1_	_	İ	İ	() _	- 1
	MOV.B @ERs,Rd	В			2							@ERs→Rd8	1-	E	1	1	() _	- 2
	MOV.B @(d:16,ERs),Rd	В				4						@(d:16,ERs)→Rd8	1-	E	1	1	() _	- 3
	MOV.B @(d:32,ERs),Rd	В				8						@(d:32,ERs)→Rd8	_	_	1	‡	() -	- 5
	MOV.B @ERs+,Rd	В					2					@ERs→Rd8,ERs32+1→ERs32	_	_	‡	‡	() -	- 3
	MOV.B @aa:8,Rd	В						2				@aa:8→Rd8	T-	F	‡	‡	() -	- 2
	MOV.B @aa:16,Rd	В						4				@aa:16→Rd8	<u> </u>	F	#	1	() -	- 3
	MOV.B @aa:32,Rd	В						6				@aa:32→Rd8	-	_	‡	‡	() -	- 4
	MOV.B Rs,@ERd	В			2							Rs8→@ERd	-	-	1	‡	() -	- 2
	MOV.B Rs,@(d:16,ERd)	В				4						Rs8→@(d:16,ERd)	T-	F	‡	‡	() -	- 3
	MOV.B Rs,@(d:32,ERd)	В				8						Rs8→@(d:32,ERd)	E	E	#	‡	() [-	- 5
	MOV.B Rs,@-ERd	В					2					ERd32-1→ERd32,Rs8→@ERd	E	E	1	‡	() -	- 3
	MOV.B Rs,@aa:8	В						2				Rs8→@aa:8		E	‡	‡	() [-	- 2
	MOV.B Rs,@aa:16	В						4				Rs8→@aa:16	E	E	1				- 3
	MOV.B Rs,@aa:32	В						6				Rs8→@aa:32	-	_	1	‡	() -	- 4
	MOV.W #xx:16,Rd	W	4									#xx:16→Rd16	_	_	1	‡	() -	- 2
	MOV.W Rs,Rd	W		2								Rs16→Rd16	-	-	1	‡	() -	- 1
	MOV.W @ERs,Rd	W			2							@ERs→Rd16	1-	F	‡	‡	() -	- 2
	MOV.W @(d:16,ERs),Rd	W				4						@(d:16,ERs)→Rd16	-	_	#	1	() -	- 3
	MOV.W @(d:32,ERs),Rd	W				8						@(d:32,ERs)→Rd16	-	_	1	‡	() -	- 5
	MOV.W @ERs+,Rd	W					2					@ERs→Rd16,ERs32+2→ERs32	-	-	1	‡	() -	- 3
	MOV.W @aa:16,Rd	W						4				@aa:16→Rd16	-	-	1	1	() -	- 3
	MOV.W @aa:32,Rd	W						6				@aa:32→Rd16	-	_	#	‡	() -	- 4
	MOV.W Rs,@ERd	W			2							Rs16→@ERd		L	‡	‡	() -	- 2
	MOV.W Rs,@(d:16,ERd)	W				4						Rs16→@(d:16,ERd)	_	_	. ‡	‡	() -	- 3
	MOV.W Rs,@(d:32,ERd)	W				8						Rs16→@(d:32,ERd)	<u> -</u>	_	#	‡) -	- 5
	MOV.W Rs,@-ERd	W					2					ERd32-2→ERd32,Rs16→@ERd			#	1	_	_	- 3
	MOV.W Rs,@aa:16	W						4				Rs16→@aa:16	上	L	1	‡	_	_	- 3
	MOV.W Rs,@aa:32	W						6				Rs16→@aa:32	1-	L	1	1	_	_	- 4
	MOV.L #xx:32,ERd	L	6									#xx:32→ERd32	-	-	1	1	_	_	- 3
	MOV.L ERs,ERd	L		2								ERs32→ERd32			#	‡	_	_	- 1
	MOV.L @ERs,ERd	L			4							@ERs→ERd32	上	L	1	‡		<u> </u>	- 4
	MOV.L @(d:16,ERs),ERd	L				6						@(d:16,ERs)→ERd32	1-	L	1	‡	_	_	- 5
	MOV.L @(d:32,ERs),ERd	L				10						@(d:32,ERs)→ERd32	-	-	1	1	_	_	- 7
	MOV.L @ERs+,ERd	L					4					@ERs→ERd32,ERs32+4→ERs32	上		1	1	_	-	- 5
	MOV.L @aa:16,ERd	L						6				@aa:16→ERd32	上		. ‡	1		_	- 5
	MOV.L @aa:32,ERd	L						8				@aa:32→ERd32	1-	E	1	1	-	_	- 6
	MOV.L ERs,@ERd	L			4							ERs32→@ERd	1-	-	1	‡		_	- 4
	MOV.L ERs,@(d:16,ERd)	L				6						ERs32→@(d:16,ERd)	上		1	1	_	_	- 5
	MOV.L ERs,@(d:32,ERd)	L				10						ERs32→@(d:32,ERd)	1			1	_	_	- 7
	MOV.L ERs,@-ERd	L					4					ERd32-4→ERd32,ERs32→@ERd	1-	-	1	1	_	_	- 5
	MOV.L ERs,@aa:16	L						6				ERs32→@aa:16			1	‡	-	_	- 5
	MOV.L ERs,@aa:32	L						8				ERs32→@aa:32	上		1	1	_	-	- 6
POP	POP.W Rn	W									2	@SP→Rn16,SP+2→SP	1-		1	1	_	_	- 3
	POP.L ERn	L									4	@SP→ERn32,SP+4→SP	1-	-	1		_	_	- 5
PUSH	PUSH.W Rn	W										SP-2→SP,Rn16→@SP	-	-	1	-	_	_	- 3
	PUSH.L ERn	L									4	SP-4→SP,ERn32→@SP	F		1	1	(-	- 5
LDM	LDM @SP+,(ERm-ERn)	L									4	(@SP→ERn32,SP+4→SP) Repeated for each register restored			F	-	ŀ		7/9/11 [1]
STM	STM (ERm-ERn),@-SP	L									4	(SP-4→SP,ERn32→@SP) Repeated for each register saved	-			-	-	\perp	7/9/11 [1]
	MOVFPE @aa:16,Rd	Can	not b	e us	ed w	ith th	is se	eries											[2]
MOV/TPE	MOVTPE Rs,@aa:16																		[2]

Arithmetic instructions

			Add	dres	sing		le/In: Byte:		ction	Len	gth									
	Mnemonic	Operand Size	*x#	Rn	@ERn	@(d,ERn)	@-ERn/@ERn+	@aa	@(d,PC)	@ @ aa	ı	Operation	C	H			_			Number of States ^{*1} Advanced
ADD	ADD.B #xx:8,Rd	В	2									Rd8+#xx:8→Rd8	_	‡	‡	1		1	‡	1
	ADD.B Rs,Rd	В		2								Rd8+Rs8→Rd8	_	‡	‡	1		‡ [‡	1
	ADD.W #xx:16,Rd	W	4									Rd16+#xx:16→Rd16	_	[3]	‡	1		1	‡	2
	ADD.W Rs,Rd	W		2								Rd16+Rs16→Rd16	_	[3]	1	1		1	‡	1
	ADD.L #xx:32,ERd	L	6									ERd32+#xx:32→ERd32	_	[4]		1		1	‡	3
	ADD.L ERs,ERd	L		2								ERd32+ERs32→ERd32	_	[4]		- '	_	1	‡	1
ADDX	ADDX #xx:8,Rd	В	2									Rd8+#xx:8+C→Rd8	┢	‡	1	1-		1	1	1
	ADDX Rs,Rd	В		2								Rd8+Rs8+C→Rd8	_	‡	1	[5]	4	ļ	1
ADDS	ADDS #1,ERd	L		2								ERd32+1→ERd32	_	_	F	₽	- -	4		1
	ADDS #2,ERd	L		2								ERd32+2→ERd32		=	F	#	#	#		1
	ADDS #4,ERd	L		2			_					ERd32+4→ERd32	\vdash	=	Ļ	1	4	#	=	1
INC	INC.B Rd	В	_	2			_					Rd8+1→Rd8	_	=	Ţ	H	+	∐-	=	1
	INC.W #1,Rd	W		2								Rd16+1→Rd16	\vdash		. ‡	1	+	-	_	1
	INC.W #2,Rd	W		2								Rd16+2→Rd16	F	=	. ‡	H	+:	-	=	1
	INC.L #1,ERd	L		2							_	ERd32+1→ERd32	=	=	1	+	+	-	=	1
D 4 4	INC.L #2,ERd DAA Rd	L B		2								ERd32+2→ERd32	F	*	1	H	١,	*	+	1
DAA SUB	SUB.B Rs,Rd	В		2								Rd8 decimal adjust → Rd8 Rd8-Rs8→Rd8	F	t	i	+ :	+	+	†	1
SUB	SUB.W #xx:16,Rd	W	4									Rd16-#xx:16→Rd16	F	[3]		÷	H	+	†	2
	SUB.W Rs,Rd	W	-	2								Rd16-Rs16→Rd16	E	[3]		+	۲	+	†	1
	SUB.L #xx:32,ERd	L	6									ERd32-#xx:32→ERd32	E	[4]		+	۲	+	†	3
	SUB.L ERs,ERd	L	0	2								ERd32-ERs32→ERd32		[4]	-	+	H	+	†	1
SUBX	SUBX #xx:8,Rd	В	2				\vdash					Rd8-#xx:8-C→Rd8		1	i	[5	1	+	†	1
CODA	SUBX Rs,Rd	В	-	2			\vdash					Rd8-Rs8-C→Rd8		i	i	1	-	+	†	1
SUBS	SUBS #1,ERd	L		2								ERd32-1→ERd32		Ė	Ľ		-		_	1
0000	SUBS #2.ERd	L		2								ERd32-2→ERd32			t	t	+	#		1
	SUBS #4,ERd	L		2								ERd32-4→ERd32			t	t	1	#		1
DEC	DEC.B Rd	В		2								Rd8-1→Rd8			t	İ	t	t -		1
	DEC.W #1,Rd	W		2								Rd16-1→Rd16	_		1	1		1 -	_	1
	DEC.W #2,Rd	W		2								Rd16-2→Rd16	_	=	. 1	1		1 -	=	1
	DEC.L #1,ERd	L		2								ERd32-1→ERd32	_	_	1	1	1	1 -	_	1
	DEC.L #2,ERd	L		2								ERd32-2→ERd32	_	_	. ‡	1	1	ţ -	_	1
DAS	DAS Rd	В		2								Rd8 decimal adjust → Rd8	_	*	‡	1	1	k _	_	1
MULXU	MULXU.B Rs,Rd	В		2								Rd8×Rs8→Rd16 (unsigned multiplication)	_	_	Ŀ	-		_[-	_	12
	MULXU.W Rs,ERd	W		2								Rd16×Rs16→ERd32 (unsigned multiplication)	_	_	╘	-	- -	_ -	_	20
MULXS	MULXS.B Rs,Rd	В		4								Rd8×Rs8→Rd16 (signed multiplication)	_	_	. ‡	1	-	_ -	_	13
	MULXS.W Rs,ERd	W		4								Rd16×Rs16→ERd32 (signed multiplication)	_	_	†			_ -	_	21
DIVXU	DIVXU.B Rs,Rd	В		2								Rd16÷Rs8→Rd16 (RdH: remainder,	_	_	[6	[7] -	- -	-	12
												RdL: quotient) (unsigned division)			L	╙	1	4		
	DIVXU.W Rs,ERd	W		2								ERd32÷Rs16→ERd32 (Ed: remainder, Rd: quotient) (unsigned division)] [7				20
DIVXS	DIVXS.B Rs,Rd	В		4								Rd16+Rs8→Rd16 (RdH: remainder, RdL: quotient) (signed division)	_	_	8]	[7] -	- -		13
	DIVXS.W Rs,ERd	W		4								ERd32÷Rs16→ERd32 (Ed: remainder, Rd: quotient) (signed division)	_		[8	[7] -	- -		21
CMP	CMP.B #xx:8,Rd	В	2									Rd8-#xx:8	_	‡	‡	1		‡	‡	1
	CMP.B Rs,Rd	В		2								Rd8-Rs8	E	‡	‡	1]	‡	‡	1
	CMP.W #xx:16,Rd	W	4									Rd16-#xx:16		[3]	1	1	I	‡	‡	2
	CMP.W Rs,Rd	W		2								Rd16-Rs16	Ē	[3]		-	-	‡	‡	1
	CMP.L #xx:32,ERd	L	6									ERd32-#xx:32	_	[4]	+ +	1	-	‡	‡	3
	CMP.L ERs,ERd	L		2								ERd32-ERs32	F	[4]	-	1	_	1	‡	1
NEG	NEG.B Rd	В		2								0-Rd8→Rd8	\vdash	‡		1		•	‡	1
	NEG.W Rd	W		2								0-Rd16→Rd16	<u> -</u>	‡		1		-	‡	1
	NEG.L ERd	L		2								0-ERd32→ERd32	_	‡		1		‡	‡	1
EXTU	EXTU.W Rd	W		2								0 → (<bits 15="" 8="" to=""> of Rd16)</bits>	-			1			-	1
	EXTU.L ERd	L		2								0 → (<bits 16="" 31="" to=""> of ERd32)</bits>	_	_	0	‡	() -	_	1

• Arithmetic instructions (cont)

			Add	dres	sing		le/In: 3yte:		ction	Len	gth								
	Mnemonic	Operand Size			ERn	@(d,ERn)	@-ERn/@ERn+	33	@(d,PC)	@aa		Operation	С	one	diti	on	Co	de	Number of States*1
		g	#XX	윤	0	0	ė	@ aa	(e)	(e)	1		ı	н	N	z	v	С	Advanced
EXTS	EXTS.W Rd	W		2								(<bit 7=""> of Rd16) → (<bits 15="" 8="" to=""> of Rd16)</bits></bit>		F	‡	‡	0	_	1
	EXTS.L ERd	L		2								(<bit 15=""> of ERd32) → (<bits 16="" 31="" to=""> of ERd32)</bits></bit>	-	F	‡	‡	0	-	1
TAS	TAS @ERd	В			4							@ERd-0 → CRR set, (1) → (<bit 7=""> of @ERd)</bit>		-	‡	‡	0	_	4
MAC	MAC @ERn+,@ERm+	-					4					@ERn×@ERm+MAC→MAC (signed multiplication) @ERn+2→ERn, ERm+2→ERm		-	[11]	[11]	[11]	_	4
CLRMAC	CLRMAC	_									2	0→MACH, MACL	Е	1=	F	F	F	F	2 [12]
LDMAC	LDMAC ERs,MACH	L		2								ERs→MACH	1	-	_	_	-	_	2 [12]
	LDMAC ERs,MACL	L		2								ERs→MACL	-	-	_	_	-	_	2 [12]
STMAC	STMAC MACH,ERd	L		2								MACH→ERd	-	-	‡	1	‡	_	1 [12]
	STMAC MACL,ERd	L		2								MACL→ERd	E	Ē	1	1	1	_	1 [12]

• Logical instructions

			Add	iress	sing		le/Ins Bytes		tion	Len	gth									
	Mnemonic	Operand Size	*xx	Rn	@ERn	@ (d,ERn)	@-ERn/@ERn+	@ aa	@(d,PC)	@ @aa		Operation		on						Number of States*1
AND	AND.B #xx:8,Rd	В	2	_	_	_		_		_	H.	Rd8∧#xx:8→Rd8	Ŀ		t	1	. '	4		1
AND	AND.B Rs,Rd	В		2								Rd8∧Rs8→Rd8	E		H	H	_) -		1
	AND.W #xx:16.Rd	w	4									Rd16∧#xx:16→Rd16	<u> </u>	L	i	i		-		2
	AND.W Rs,Rd	w	Ė	2								Rd16∧Rs16→Rd16	<u> </u>		i	Ħ	+) -	_	1
	AND.L #xx:32.ERd	L	6	_								ERd32∧#xx:32→ERd32	1=		i	Ì	1	-	_	3
	AND.L ERs.ERd	Ī	_	4								ERd32∧ERs32→ERd32	1=		li	İ	_	5 -	_	2
OR	OR.B #xx:8,Rd	В	2									Rd8√#xx:8→Rd8	<u> </u>	_	İ	1	(5 -	_	1
	OR.B Rs,Rd	В		2								Rd8∨Rs8→Rd8	<u> </u>	_	1	1	() -	=	1
	OR.W #xx:16,Rd	W	4									Rd16√#xx:16→Rd16	<u> </u>	_	1	1	() -	_	2
	OR.W Rs,Rd	W		2								Rd16∨Rs16→Rd16	<u> </u>	_	1	1	(5 -	=	1
	OR.L #xx:32,ERd	L	6									ERd32√#xx:32→ERd32	-	_	1	1	(5 -	_	3
	OR.L ERs,ERd	L		4								ERd32√ERs32→ERd32	_	_	1	1	() -	-1	2
XOR	XOR.B #xx:8,Rd	В	2									Rd8⊕#xx:8→Rd8	-	_	‡	1	() -	-[1
	XOR.B Rs,Rd	В		2								Rd8⊕Rs8→Rd8	-	-	1	ļ	(5 -	-1	1
	XOR.W #xx:16,Rd	W	4									Rd16⊕#xx:16→Rd16	_	_	. ‡	1	() -	=	2
	XOR.W Rs,Rd	W		2								Rd16⊕Rs16→Rd16		Ŀ	‡	_ţ	() -	=	1
	XOR.L #xx:32,ERd	L	6									ERd32⊕#xx:32→ERd32	-	_	‡	1	() -	-	3
	XOR.L ERs,ERd	L		4								ERd32⊕ERs32→ERd32	-	_	1	1	() <u> </u>	-	2
NOT	NOT.B Rd	В		2								¬Rd8→Rd8	-	-	1	1	() -	-[1
	NOT.W Rd	W		2								¬Rd16→Rd16	_		1	1	(o [-	=[1
	NOT.L ERd	L		2								¬ERd32→ERd32	-	1-	1	1	() -	-	1

• Shift instructions

				Add	dres	sing		le/In: Byte:		ction	Len	gth								
SHALL BRALL BR Rd B 2 2		Mnemonic	perand Size	×	_	ERn	(d,ERn)	-ERn/@ERn+	aa	(d,PC)	@aa		Operation					ı		Number of States*1
SHALB #2.Rd B 2 2 SHALW Rd W 2 2 SHALW Rd W 2 2 SHALW Rd W 2 2 SHALW Rd W 2 2 SHALW Rd W 2 2 SHALW 2.Rd L 2 2 SHAR B #2.Rd B 2 2 SHAR W #2.Rd W 2 2 SHAR W #2.Rd W 2 2 SHAR W #2.Rd W 2 2 SHAR B #3.Rd B B 2 2 SHAR W #3.Rd B B 2 3 SHAR W #3.Rd		T	-	#		0	0	0	0	(a)	0	1		ı	н	N	z	-	-	
SHALW #2.Rd W 2 2	SHAL		_		_										_	1	ļ	,	+ :	
SHALL #2,FRd		,	_	_											_	1	ļ	ļ	+:	
SHALL ERD			-		_									_	_	1	1	ļ	1	
SHALL #2,ERd													C MSB ← LSB		_	‡	ļ	1	1	
SHAR BHARB Rd B 2 2			_												_	1	ļ		+ :	
SHARB #2,Rd		· · · · · · · · · · · · · · · · · · ·	-	_	_										_	1	ļ	,	1	
SHAR.W Rd W 2 SHARL W #2,Rd W 2 SHARL ERG L 2 SHARL ERG B C SHARL BRG SHARL BRG SHARL BRG SHARL BRG SHARL BRG SHARL BRG SHARL W 2 SHARL BRG SHARL	SHAR		_		_									_	_	‡	1		1	
SHAR.W #2,Rd					_										_	-				
SHARL ERD SHARL #2,ERD LL SHLLB RD SHLLB #2,RD SHLLW RD SHLLW RD SHLLW RD SHLLW RD SHLLW RD SHLLW RD SHLLW RD SHLLW RD SHLLW RD SHLLW RD SHLLW RD SHLLW RD SHLLW RD SHLLW RD SHLLW RD SHLLW RD SHLR BRD SHLR BRD SHLR BRD SHLRW RD SH SHLW RD SH SHLW RD SH SH SH SH SH SH SH SH SH SH SH SH SH			_											_	_	-	-	-	+ -	
SHARL ERD SHARL #2,ERD SHLL BRD SHLLB #2,RD SHLLW RD SHLLW RD SHLLW RD SHLLW RD SHLLW RD SHLLW RD SHLLW RD SHLLW RD SHLLW RD SHLLW RD SHLLW RD SHLLW RD SHLLW RD SHLLW RD SHLLW RD SHLLW RD SHLLW RD SHLL #2,ERD SHLR BRD SHLL BRD SHR BRD SHR BRD SHR BRD SHR BRD SHR BRD SHR BRD SHR BRD SHR BRD SHR BRD SHR BRD SHR BRD SHR BRD SHR BRD SHR BRD SHR BRD SHR BRD		SHAR.W #2,Rd	W		_								MSB — ISB C		_	_				
SHLL SHLL.B #2,Rd B 2 SHLL.B #2,Rd B 2 SHLL.W Rd W 2 SHLL.W Rd W 2 SHLL.W #2,Rd W 2 SHLL.L #2,ERd L 2 SHLL.L #2,ERd L 2 SHLR.B #2,Rd B 2 SHLR.B #2,Rd B 2 SHLR.B #2,Rd B 2 SHLR.W Rd W 2 SHLR.W #2,Rd W 2 SHLR.B #2,Rd B 2 SHLR.B #2,Rd B 2 SHLR.B #2,Rd B 2 SHLR.W #2,Rd W 2 SHLR.W #2,Rd B 2 SHLR.W #2,Rd B 2 SHLR.W #2,Rd B 2 ROTXL.B #2,Rd B 2 ROTXL.B #2,Rd B 2 ROTXL.W #2,Rd W 2 ROTXR.B #2,Rd B 2 RO		SHAR.L ERd	L		2								MIOD EOD C	_	_	‡	‡	0	1	1
SHLL.B #2,Rd B 2 SHLL.W Rd W 2 SHLL.W #2,Rd W 2 SHLL.L #2,ERd L 2 SHLL.L #2,ERd L 2 SHLR.B Rd B 2 SHLR.W RLR B Rd B 2 SHLR.W #2,Rd B 2 SHLR.W #2,Rd W 2 SHLR.L #2,ERd L 2 ROTXL. B #2,Rd B 2 ROTXL.W #2,Rd B 2 ROTXL.W #2,Rd B 2 ROTXL.W #2,Rd W 2 ROTXL.W #2,Rd W 2 ROTXL.L #2,ERd L 2 ROTXL.L #2,ERd L 2 ROTXL. #2,ERd B 2 ROTXR.D #2,Rd B 2 ROTX.		SHAR.L #2,ERd	L		2									_	_	‡	‡	0	1	1
SHLLW Rd W 2 SHLLW #2,Rd W 2 SHLLW #2,Rd W 2 SHLLL ERd L 2 SHLL #2,ERd L 2 SHLR B Rd B 2 SHLR B #2,Rd B 2 SHLR W Rd W 2 SHLR W Rd W 2 SHLR W #2,Rd W 2 SHLR, ERd L 2 SHLR, ERd L 2 SHLR, B #2,Rd W 2 SHLR, B #2,Rd B 2 ROTXL, B #2,Rd B 2 ROTXL, B #2,Rd B 2 ROTXL, B #2,Rd B 2 ROTXL, B #2,Rd B 2 ROTXR, B #2,Rd B 2 ROTXR, B #2,Rd B 2 ROTXR, B #2,Rd B 2 ROTXR, B #2,Rd B 2	SHLL	SHLL.B Rd	В		2									_		‡	‡	0	1	1
SHLL.W #2,Rd W 2 SHLLL.ERd L 2 SHLLL.BERd L 2 SHLR.B Rd B 2 SHLR.B Rd B 2 SHLR.B #2,Rd B 2 SHLR.W Rd W 2 SHLR.W #2,Rd W 2 SHLR.L #2,ERd L 2 ROTXL.B #2,Rd B 2 ROTXL ROTXL B #2,Rd B 2 ROTXL ROTXL B #2,Rd B 2 ROTXL W #2,Rd W 2 ROTXL #2,ERd L 2 ROTXR B #2,Rd B 2 ROTXR W #2,Rd W 2 ROTXR W #2,Rd W 2 ROTL B Rd B 2		SHLL.B #2,Rd	В		2									-	_	‡	1	0	1	1
SHLLL ERd		SHLL.W Rd	W		2									-		‡	‡	0	1	1
SHILL ERG SHLLL #2,ERG SHLR BRG SHLR BRG SHLR BRG SHLR BRG SHLR BRG SHLR W #2,RG SHLR W #2,RG SHLR W #2,RG SHLR W #2,RG SHLR L #2,ERG SHLR BRG SHLR BRG SHLR BRG SHLR BRG SHLR BRG SHLR W #2,RG SHLR BRG SHAND		SHLL.W #2,Rd	W		2								C MSB ← LSB	-		‡	1	0	1	1
SHLR SHLR.B Rd B 2 SHLR.B #2,Rd B 2 SHLR.W Rd W 2 SHLR.W #2,Rd W 2 SHLR.L #2,ERd L 2 SHLR.L #2,ERd L 2 SHLR.L #2,ERd L 2 ROTXL.B Rd B 2 ROTXLB #2,Rd B 2 ROTXLL #2,ERd W 2 ROTXLL #2,ERd W 2 ROTXLL #2,ERd W 2 ROTXLL #2,ERd W 2 ROTXLL #2,ERd L 2 ROTXLL #2,ERd L 2 ROTXLL #2,ERd L 2 ROTXL, #2,ERd L 2 ROTXR, #32,Rd B 2 ROTXR, #32,Rd B 2 ROTXR, #32,Rd B 2 ROTXR, W2,Rd W 2 ROTXR, W2,Rd W 2 ROTX, W2,Rd W 2 ROTL, W2,Rd W 2 ROTL, W2,Rd W 2<		SHLL.L ERd	L		2								O WIGHT LOD		_	1	1	0	1	1
SHLR.B #2.Rd B 2 SHLR.W Rd W 2 SHLR.W #2.Rd W 2 SHLR.L ERd L 2 SHLR.L #2.ERd L 2 SHLR.L #2.ERd L 2 SHLR.L #2.ERd B 2 ROTXL.B #2.Rd B 2 ROTXL.W #2.Rd W 2 ROTXL.L #2.ERd L 2 ROTXL.L #2.ERd L 2 ROTXL.L #2.ERd B 2 ROTXR.B #2.Rd B 2 ROTXR.W #2.Rd B 2 ROTXR.W #2.Rd W 2 ROTXR.B Rd B 2 ROTXR.L #2.ERd L 2 ROTX.B Rd B 2 ROTX.B Rd B 2 ROTL.B #2.ERd W 2		SHLL.L #2,ERd	L		2									=	_	1	1	0	1	1
SHLR.W Rd W 2 SHLR.W #2,Rd W 2 SHLR.L #2,ERd L 2 SHLR.L #2,ERd L 2 ROTXL ROTXL.B Rd B 2 ROTXL.W #2,Rd B 2 ROTXL.W #2,Rd W 2 ROTXL.L #2,ERd L 2 ROTXL.L #2,ERd L 2 ROTXR.B #2,Rd B 2 ROTXR.B #2,Rd B 2 ROTXR.B #2,Rd B 2 ROTXR.B #2,Rd B 2 ROTXR.W #2,Rd W 2 ROTXR.W #2,Rd W 2 ROTXR.L #2,ERd L 2 ROTXR.L #2,ERd L 2 ROTXR.L #2,ERd L 2 ROTX.W #2,Rd W 2 ROTL.W Rd W 2 ROTL.W Rd W 2 ROTL.W Rd W 2 ROTR.B Rd B 2 ROTR.B Rd B 2 ROTR.B Rd B 2 <	SHLR	SHLR.B Rd	В		2									_	_	0	1	0	1	1
SHLR.W #2,Rd W 2 SHLR.L ERd L 2 SHLR.L #2,ERd L 2 ROTXL ROTXL.B Rd B 2 ROTXL.B #2,Rd B 2 ROTXL.W Rd W 2 ROTXL.W #2,Rd W 2 ROTXL.L #2,ERd L 2 ROTXL.L #2,ERd L 2 ROTXL.L #2,ERd L 2 ROTXL.B #2,Rd B 2 ROTXR.B #2,Rd B 2 ROTXR.W #2,Rd W 2 ROTXR.W #2,Rd W 2 ROTXR.B #2,Rd B 2 ROTXR.B #2,Rd B 2 ROTXR.B #2,Rd B 2 ROTX.B #2,Rd B 2 ROTL.B Rd B 2 ROTL.W #2,Rd W 2 ROTL.W #2,Rd W 2 ROTL.E Rd L 2 ROTL.B Rd B 2 ROTL.B Rd B 2 ROTR.B #2,Rd B 2		SHLR.B #2,Rd	В		2										_	0	1	0	1	1
SHLR.L ERd		SHLR.W Rd	W		2								0 →		_	0	1	0	1	1
SHLR.L ERd		SHLR.W #2,Rd	W		2								MSB — JSB C	=	_	0	1	0	1	1
ROTXL ROTXLB Rd B 2		SHLR.L ERd	L		2								WIGD LOD C		_	0	1	0	1	1
ROTXL.B #2,Rd		SHLR.L #2,ERd	L		2									_	_	0	1	0	1	1
ROTXL.W Rd W 2	ROTXL	ROTXL.B Rd	В		2									_	_	‡	1	0	1	1
ROTXL.W #2,Rd W 2 C MSB		ROTXL.B #2,Rd	В		2										_	1	1	0	1	1
ROTXLL ERd		ROTXL.W Rd	W		2									_	_	1	1	0	1	1
ROTXLL #2,ERd		ROTXL.W #2,Rd	W		2										_	‡	1	0	1	1
ROTXL #2,ERd			L		2								C MSB ← LSB		_	İ	İ	0	İ	1
ROTXR.B #2,Rd		ROTXL.L #2,ERd	L		2										_	İ	İ			1
ROTXR.B #2,Rd	ROTXR	ROTXR.B Rd	В		2										_	1	İ	0	1	1
ROTXR.W #2,Rd W 2	-		В		_										_	1	İ	0	Ħ	1
ROTXR.W #2,Rd W 2		ROTXR.W Rd	W		2										_	1	1	0	Ť	1
ROTXR.L ERd			w		2										_	t	İ		Ì	1
ROTXR.L #2,ERd													MSB ──→ LSB C		_	i	İ			
ROTL BRd B 2			_												=	-	İ		+ :	
ROTL.B #2,Rd B 2	ROTL	,	_		_										=				+ -	
ROTL.W Rd W 2 ROTL.W #2,Rd W 2 ROTL.L ERd L 2 ROTL.L #2,ERd L 2 ROTR.B Rd B 2 ROTR.B #2,Rd B 2 ROTR.W #2,Rd W 2 ROTR.W Rd W 2 ROTR.W Rd W 2 ROTR.W #2,Rd W 2 ROTR.W #2,Rd W 2			-													t	i		+ :	
ROTL.W #2,Rd W 2 C MSB - LSB - 0 1 1 1 1 1 1			-													1	Ť	-	+ :	
ROTL L ERd L 2																_	Ť	-	-	
ROTL. #2,ERd L 2 ROTR ROTR.B Rd B 2 ROTR.B #2,Rd B 2 ROTR.W Rd W 2 ROTR.W #2,Rd W 2 MSR — ISB C — I I 0 I 1					_								C MSB ← LSB			_	i		+ '	
ROTR ROTR.B Rd B 2			_													_				
ROTR.B #2,Rd B 2 ROTR.W Rd W 2 ROTR.W #2,Rd W 2 MSR — ISB C — I I 0 I 1	ROTE	· ·	-		_							Н				-				
ROTR.W Rd W 2	INO IN		_	_												_	-			
ROTR.W #2,Rd W 2 MSR \longrightarrow LSR C $$ \downarrow \downarrow 0 \downarrow 1			_														\rightarrow		+ -	
$MSB \longrightarrow ISB V$			_	_										H	3		+	+	+ :	
NOTICE NO. L Z				_									MSB ──→ LSB C	Θ	Ξ	-	+	-	+ '	
ROTR.L #2,ERd L 2			_		_									\exists	Ξ	+	+		+	

• Bit-manipulation instructions

			Add	dres	sing	Moc	le/In:		ction	Len	gth								
	Mnemonic	Operand Size	*xx	Rn	@ERn	@(d,ERn)	@-ERn/@ERn+	@aa	@(d,PC)	@ @ aa	1	Operation			diti				Number of States*1
BSET	BSET #xx:3,Rd	В		2								(#xx:3 of Rd8)←1			t	t	t	t	1
	BSET #xx:3,@ERd	В			4							(#xx:3 of @ERd)←1			1	1	1		4
	BSET #xx:3,@aa:8	В						4				(#xx:3 of @aa:8)←1		_	_	_	_		4
	BSET #xx:3,@aa:16	В						6				(#xx:3 of @aa:16)←1	=	=	1	1	1		5
	BSET #xx:3,@aa:32	В						8				(#xx:3 of @aa:32)←1	=	=	F	F	F		6
	BSET Rn,Rd	В		2								(Rn8 of Rd8)←1	=	_	_	_	_		1
	BSET Rn,@ERd	В			4							(Rn8 of @ERd)←1	_	_	_	_	_		4
	BSET Rn,@aa:8	В						4				(Rn8 of @aa:8)←1	-	_	_	_	_		4
	BSET Rn,@aa:16	В						6				(Rn8 of @aa:16)←1	Н	_	Н	Н	Н		5
	BSET Rn,@aa:32	В						8				(Rn8 of @aa:32)←1	-	_	-	-	-		6
BCLR	BCLR #xx:3,Rd	В		2								(#xx:3 of Rd8)←0	\equiv	Ē	Ē	Ē	Ē	Ŀ	1
	BCLR #xx:3,@ERd	В			4							(#xx:3 of @ERd)←0			E	E	E	L	4
	BCLR #xx:3,@aa:8	В						4				(#xx:3 of @aa:8)←0			E	E	E	L	4
	BCLR #xx:3,@aa:16	В						6				(#xx:3 of @aa:16)←0				E	E	1	5
	BCLR #xx:3,@aa:32	В						8				(#xx:3 of @aa:32)←0			╘	╘	╘	╙	6
	BCLR Rn,Rd	В		2								(Rn8 of Rd8)←0			╘	╘	╘	╙	1
	BCLR Rn,@ERd	В			4							(Rn8 of @ERd)←0		_	-	1	1	╙	4
	BCLR Rn,@aa:8	В						4				(Rn8 of @aa:8)←0	=		┢	E	╘	_	4
	BCLR Rn,@aa:16	В						6				(Rn8 of @aa:16)←0			╘	╘	╘	╙	5
	BCLR Rn,@aa:32	В						8				(Rn8 of @aa:32)←0			╘	╘	╘	╙	6
BNOT	BNOT #xx:3,Rd	В		2								(#xx:3 of Rd8)← [¬(#xx:3 of Rd8)]		_	⊨	⊨	⊨	┶	1
	BNOT #xx:3,@ERd	В			4							(#xx:3 of @ERd)← [¬(#xx:3 of @ERd)]			⊨	⊨	⊨	╙	4
	BNOT #xx:3,@aa:8	В						4				(#xx:3 of @aa:8)← [¬(#xx:3 of @aa:8)]			⊨	⊨	⊨	╙	4
	BNOT #xx:3,@aa:16	В						6				(#xx:3 of @aa:16)←[¬(#xx:3 of @aa:16)]			⊨	⊨	⊨	1	5
	BNOT #xx:3,@aa:32	В		_				8				(#xx:3 of @aa:32)← [¬(#xx:3 of @aa:32)]			F	F	F	1	6
	BNOT Rn,Rd	В		2								(Rn8 of Rd8)← [¬(Rn8 of Rd8)]			F	F	F	1	1
	BNOT Rn,@ERd	В		_	4	_		<u>.</u>				(Rn8 of @ERd)← [¬(Rn8 of @ERd)]	=	E	F	F	F	1	4
	BNOT Rn,@aa:8	В						4				(Rn8 of @aa:8)← [¬(Rn8 of @aa:8)]		E	⊨	⊨	⊨	1	4
	BNOT Rn,@aa:16	В						6				(Rn8 of @aa:16)← [¬(Rn8 of @aa:16)]		=	F	⊨	⊨	1	5
DT0T	BNOT Rn,@aa:32	В		_		-		8				(Rn8 of @aa:32)← [¬(Rn8 of @aa:32)]		F	F	1	⊨	1	6
BTST	BTST #xx:3,Rd	В		2	-	_						¬(#xx:3 of Rd8)→Z		F	F	H	F	-	1
	BTST #xx:3,@ERd	В			4			-				¬(#xx:3 of @ERd)→Z		E	F	H	F	╄	3
	BTST #xx:3,@aa:8	В						4				¬(#xx:3 of @aa:8)→Z	=	E	₽	H	F	╄	3
	BTST #xx:3,@aa:16	В						6				¬(#xx:3 of @aa:16)→Z	=	F	F	+	F	+	4
	BTST #xx:3,@aa:32 BTST Rn,Rd	В		_				8				¬(#xx:3 of @aa:32)→Z	=	F	F	+	F	-	5
	· ·	В		2	4							¬(Rn8 of Rd8)→Z		F	F	+	F	1	1
	BTST Rn,@ERd BTST Rn,@aa:8	В			+			4				¬(Rn8 of @ERd)→Z ¬(Rn8 of @aa:8)→Z	Н	Е	E	1	E	+	3
	BTST Rn,@aa:16	В						6				¬(Rn8 of @aa:16)→Z	Η	Е	E	1	E	+	4
	BTST Rn,@aa:32	В						8				¬(Rn8 of @aa:16)→Z ¬(Rn8 of @aa:32)→Z		E	E	1	E	+	5
BLD	BLD #xx:3,Rd	В		2				٦				(#xx:3 of Rd8)→C	H	E	E	Ľ	E	1	1
	BLD #xx:3,@ERd	В		<u> </u>	4							(#xx:3 of @ERd)→C		Ē	E	E	E	H	3
	BLD #xx:3,@aa:8	В			-			4				(#xx:3 of @aa:8)→C	\exists		E	Ē	Ē	1	3
	BLD #xx:3,@aa:16	В						6				(#xx:3 of @aa:16)→C			E	E	E	i	4
	BLD #xx:3,@aa:32	В						8				(#xx:3 of @aa:32)→C			E	E	E	H	5
BILD	BILD #xx:3,Rd	В		2				Ť				¬(#xx:3 of Rd8)→C	\exists		t	E	E	H	
	BILD #xx:3,@ERd	В		_	4							¬(#xx:3 of @ERd)→C	\exists		t	E	t	i	<u> </u>
	BILD #xx:3,@aa:8	В			Ť			4				¬(#xx:3 of @aa:8)→C	H		t	E	t	l i	3
	BILD #xx:3,@aa:16	В						6				¬(#xx:3 of @aa:16)→C	H		t	t	t	H	4
	BILD #xx:3,@aa:32	В						8				¬(#xx:3 of @aa:32)→C			t	t	t	H	5
BST	BST #xx:3,Rd	В		2				Ť				C→(#xx:3 of Rd8)	\equiv	=	t	t	t	Ť	1
	BST #xx:3,@ERd	В		Ē	4							C→(#xx:3 of @ERd)			1=		t	1	4
	BST #xx:3,@aa:8	В						4				C→(#xx:3 of @aa:8)	=	<u> </u>	E	E	E		4
	BST #xx:3,@aa:16	В						6				C→(#xx:3 of @aa:16)	\equiv	<u> </u>	E	E	F		5
	BST #xx:3,@aa:32	В						8				C→(#xx:3 of @aa:32)		_	1_	1_	1_		6

• Bit-manipulation instructions (cont)

			Ad	dres	sing		de/In Byte		ction	Len	gth								
	Mnemonic	Operand Size	,		ERn	@(d,ERn)	@-ERn/@ERn+	aa	@(d,PC)	@aa		Operation	C	on	diti	on (Cod	de	Number of States*
		g	XX#	ೱ	@	(e)	e	00	@	0	1		ı	Н	N	z	٧	С	Advanced
BIST	BIST #xx:3,Rd	В		2								¬C→(#xx:3 of Rd8)		L	L	L	_	_	1
	BIST #xx:3,@ERd	В			4							¬C→(#xx:3 of @ERd)		┕	<u> </u>	上	_	<u> </u>	4
	BIST #xx:3,@aa:8	В						4				¬C→(#xx:3 of @aa:8)	-	-	<u> </u>	<u> -</u>	_	_	4
	BIST #xx:3,@aa:16	В						6				¬C→(#xx:3 of @aa:16)		-	-	L	—	_	5
	BIST #xx:3,@aa:32	В						8				¬C→(#xx:3 of @aa:32)	Ŀ	Ŀ	Ŀ	Ŀ	Ē	Ŀ	6
BAND	BAND #xx:3,Rd	В		2								C∧(#xx:3 of Rd8)→C	Œ	E	E	E	E	1	1
	BAND #xx:3,@ERd	В			4							C∧(#xx:3 of @ERd)→C	-	-	-	_	_	1	3
	BAND #xx:3,@aa:8	В						4				C∧(#xx:3 of @aa:8)→C	_	-	-	_	_	1	3
	BAND #xx:3,@aa:16	В						6				C∧(#xx:3 of @aa:16)→C		-	-	_	_	1	4
	BAND #xx:3,@aa:32	В						8				C∧(#xx:3 of @aa:32)→C		-	-	_	_	1	5
BIAND	BIAND #xx:3,Rd	В		2								C∧[¬(#xx:3 of Rd8)]→C	-	-	-	_	_	1	1
	BIAND #xx:3,@ERd	В			4							C∧[¬(#xx:3 of @ERd)]→C	_	-	-	<u> </u>	_	1	3
	BIAND #xx:3,@aa:8	В						4				C∧[¬(#xx:3 of @aa:8)]→C		1=	1=	_	_	1	3
	BIAND #xx:3,@aa:16	В						6				C∧[¬(#xx:3 of @aa:16)]→C	-	-	1_	<u> </u>	_	1	4
	BIAND #xx:3,@aa:32	В						8				C_\[¬(#xx:3 of @aa:32)]→C	-	-	_	_	_	1	5
BOR	BOR #xx:3,Rd	В		2								C√(#xx:3 of Rd8)→C	-	1=	1_	<u> </u>	_	1	1
	BOR #xx:3,@ERd	В			4							C√(#xx:3 of @ERd)→C	_	-	-	_	_	1	3
	BOR #xx:3.@aa:8	В						4				C√(#xx:3 of @aa:8)→C	_	1=	1_	<u> </u>	_	1	3
	BOR #xx:3,@aa:16	В						6				C√(#xx:3 of @aa:16)→C	_	-	1_	<u> </u>	_	1	4
	BOR #xx:3,@aa:32	В						8				C√(#xx:3 of @aa:32)→C	_	1=	1_			i	5
BIOR	BIOR #xx:3.Rd	В		2								C√[¬(#xx:3 of Rd8)]→C	_	1=	1_	<u> </u>		i	1
	BIOR #xx:3,@ERd	В			4							C√[¬(#xx:3 of @ERd)]→C	_	-	1_		_	İ	3
	BIOR #xx:3,@aa:8	В						4				C√[¬(#xx:3 of @aa:8)]→C	_	-	1_	_	_	1	3
	BIOR #xx:3.@aa:16	В						6				C√[¬(#xx:3 of @aa:16)]→C	_	1=	1_			i	4
	BIOR #xx:3.@aa:32	В						8				C√[¬(#xx:3 of @aa:32)]→C	_	1=	1_	<u> </u>		i	5
BXOR	BXOR #xx:3.Rd	В		2								C⊕(#xx:3 of Rd8)→C	_	1_	1_	<u> </u>	_	İ	1
	BXOR #xx:3,@ERd	В			4							C⊕(#xx:3 of @ERd)→C	_	E	L	E	<u> </u>	Ħ	3
	BXOR #xx:3,@aa:8	В						4				C⊕(#xx:3 of @aa:8)→C		E	1_	\vdash	<u> </u>	İ	3
	BXOR #xx:3,@aa:16	В						6				C⊕(#xx:3 of @aa:16)→C		1	1_		<u> </u>	İ	4
	BXOR #xx:3,@aa:32	В						8				C⊕(#xx:3 of @aa:32)→C	_	-	1_	_	E	1	5
BIXOR	BIXOR #xx:3,Rd	В		2								C⊕[¬(#xx:3 of Rd8)]→C	_	E	1_	\vdash	<u> </u>	İ	1
	BIXOR #xx:3,@ERd	В			4							C⊕[¬(#xx:3 of @ERd)]→C	_	t	1=	\vdash	⊏	İ	3
	BIXOR #xx:3,@aa:8	В			Ė			4				C⊕[¬(#xx:3 of @aa:8)]→C		1	⇇			Ħ	3
	BIXOR #xx:3,@aa:16	В						6				C⊕[¬(#xx:3 of @aa:16)]→C	_		1_	E	E	Ħ	4
	BIXOR #xx:3.@aa:32	В						8				C⊕[¬(#xx:3 of @aa:32)]→C	_		₶	\vdash		Ť	5

• Branch instructions

			Add	dres	sing		le/In: Byte:		ction	Len	gth									
	Mnemonic	Operand Size			Rn	@(d,ERn)	@-ERn/@ERn+		@(d,PC)	aa		Operation		C	onc	ditio	on	Co	de	Number of States ^{*1}
		o	××	조	@ERn	@(0	9	@aa	9	0	ı		Branch Condition	ı	н	N	z	٧	С	Advanced
Bcc	BRA d:8(BT d:8)								2			if condition is true then	Always		_	_	_	_	_	2
	BRA d:16(BT d:16)	_							4			PC←PC+d		_	_	_	_	_	_	3
	BRN d:8(BF d:8)								2			else next;	Never	=	_	_	_	_	_	2
	BRN d:16(BF d:16)								4					<u> -</u>	_	_	<u> </u>	_	_	3
	BHI d:8	_							2				C∨Z=0	_	_	_	_	_	_	2
	BHI d:16								4					_	_	_	_	_	_	3
	BLS d:8								2				C∨Z=1	느	느	=	느		느	2
	BLS d:16								4					\vdash	⊨	_	⊨	_	⊨	3
	BCC d:8(BHS d:8)	_							2				C=0	\vdash		_	E	_	-	2
	BCC d:16(BHS d:16)	_							4					_	_	_	_	_	_	3
	BCS d:8(BLO d:8)								2				C=1	_	_	_	_	_		2
	BCS d:16(BLO d:16)								4					_	_	_	_	_	_	3
İ	BNE d:8	_							2				Z=0	_	_	_	_	_	_	2
İ	BNE d:16	<u> </u>							4					<u> </u>	_	_	_	_	_	3
	BEQ d:8	_							2				Z=1	_	_	_	_	_	_	2
	BEQ d:16	_							4					_	_	_	_	_	_	3
	BVC d:8	_							2				V=0	_	_	_	_	_	_	2
	BVC d:16	_							4					_	_	_	_	_	_	3
ĺ	BVS d:8	<u> </u>							2				V=1	_	—	_	_	—	_	2
İ	BVS d:16	_							4					_	_	_	_	_	_	3
	BPL d:8	_							2				N=0	上	_	_	_	_	_	2
	BPL d:16	_							4					_	_	_	_	_	_	3
	BMI d:8	_							2				N=1	_	_	_	_	_	_	2
	BMI d:16	_ _							4					_	_	_	_	_	_	3
	BGE d:8								2				N⊕V=0	_	_	_	_	_	_	2
	BGE d:16								4					_	_	_	_	_	_	3
	BLT d:8	_							2				N⊕V=1	_	—	_	_	_	_	2
	BLT d:16	_							4					E			E		E	3
	BGT d:8	_							2				Z∨(N⊕V)=0		_	_		_	_	2
i	BGT d:16								4					<u> </u>	_	_	_	_		3
	BLE d:8								2				Z∨(N⊕V)=1		_	_	L	_	Ŀ	2
	BLE d:16								4					F					E	3
JMP	JMP @ERn				2							PC←ERn		E		E	Ē		E	2
	JMP @aa:24	_						4				PC←aa:24		E			E		E	3
	JMP @@aa:8	_								2		PC←@aa:8		E					E	5
BSR	BSR d:8								2			PC→@-SP,PC←PC+d:8	3	E			E		E	4
	BSR d:16	_							4			PC→@-SP,PC←PC+d:1	16		=		_		E	5
JSR	JSR @ERn	_			2							PC→@-SP,PC←ERn		E	E		E		E	4
Í	JSR @aa:24	_						4				PC→@-SP,PC←aa:24		<u> </u>	_	_	_	_	_	5
<u></u>	JSR @@aa:8					Ĺ				2		PC→@-SP,PC←@aa:8		E	E	_			E	6
RTS	RTS	_									2	PC←@SP+		<u> </u>	_	_		_	<u> </u>	5

• System control instructions

			Add	dres	sing		le/In: Byte:		ction	Len	gth								
	Mnemonic	Operand Size	×		ERn	@(d,ERn)	@-ERn/@ERn+	aa	@(d,PC)	@aa		Operation	С	one	diti	on	Co	de	Number of States*1
	T=	ō	#	돈	0	0	@	(9)	@	(9)	ı		1	Н	N	z	٧	С	Advanced
TRAPA	TRAPA #xx:2	-										PC→@-SP,CCR→@-SP, EXR→@-SP, <vector>→PC</vector>	1			-	Г	_	8 [9]
RTE	RTE	-										EXR←@SP+,CCR←@SP+, PC←@SP+	1	‡	1	‡	‡	‡	5 [9]
SLEEP	SLEEP	_										Transition to the power-down state	-	_	_	-	_	_	2
LDC	LDC #xx:8,CCR	В	2									#xx:8→CCR	1	1	1	1	1	1	1
	LDC #xx:8,EXR	В	4									#xx:8→EXR	1-	Ė	Ė	Ė	Ė	Ė	2
	LDC Rs,CCR	В		2								Rs8→CCR	1	1	1	1	1	1	1
	LDC Rs,EXR	В		2								Rs8→EXR	1-	Ė	Ė	Ė	Ė	Ė	1
	LDC @ERs,CCR	w			4							@ERs→CCR	1	1	1	1	1	1	3
	LDC @ERs,EXR	w			4							@ERs→EXR	Ť	Ė	Ė	Ė	Ė	Ė	3
	LDC @(d:16,ERs),CCR	w			Ė	6						@(d:16,ERs)→CCR	1	1	1	1	1	1	4
	LDC @(d:16,ERs),EXR	W				6						@(d:16,ERs)→EXR	ΤĖ	Ė	Ė	Ė	Ė	Ė	4
	LDC @(d:32,ERs),CCR	w				10						@(d:32,ERs)→CCR	l	t	t	I	1	I	6
	LDC @(d:32,ERs),EXR	w				10						@(d:32,ERs)→EXR	ť	Ė	Ľ	Ė	Ė	Ľ	6
	LDC @ERs+,CCR	w					4					@ERs→CCR,ERs32+2→ERs32	1	t	t	Ιt	t	Ιŧ	4
	LDC @ERs+,EXR	w					4					@ERs→EXR,ERs32+2→ERs32	Ľ	Ľ	Ľ	Ľ	Ľ	Ľ	4
	LDC @aa:16,CCR	w					_	6				@aa:16→CCR	I	l t	1	l t	t	l t	4
	LDC @aa:16,EXR	w						6				@aa:16→EXR	Ľ	Ľ	Ľ	Ľ	Ľ	Ľ	4
	LDC @aa:32.CCR	W						8				@aa:32→CCR	t	1	+	1	t	1	5
	LDC @aa:32,EXR	w						8				@aa:32→EXR	Ľ	Ľ	Ľ	Ľ	Ľ	Ľ	5
STC	STC CCR,Rd	В	\vdash	2				۳				CCR→Rd8	E	⊨	H	\vdash	F	F	1
310	STC EXR.Rd	В		2								EXR→Rd8	F	F	F	F	F	F	1
	STC CCR,@ERd	W		-	4							CCR→@ERd	₽	F	F	F	F	F	3
		W	-		4							EXR→@ERd	₽	F	F	\vdash	F	F	3
	STC EXR,@ERd	W	\vdash	-	4	6				_			\vdash	F	F	F	F	F	4
	STC CCR,@(d:16,ERd)	+	-		-	6				_		CCR→@(d:16,ERd)	\vdash	F	F	\vdash	F	F	· ·
	STC EXR,@(d:16,ERd)	W	-		-	<u> </u>				_		EXR→@(d:16,ERd)	+	F	F	F	F	F	4
	STC CCR,@(d:32,ERd)	W	\vdash		\vdash	10						CCR→@(d:32,ERd)	+	F	F	F	F	F	6
	STC EXR,@(d:32,ERd)	W	\vdash	-	-	10	-			_		EXR→@(d:32,ERd)	+	F	F	\vdash	F	F	6
	STC CCR,@-ERd	W	-	-	-		4			_		ERd32-2→ERd32,CCR→@ERd	+	F	F	F	F	F	4
	STC EXR,@-ERd	W	-	-	-		4	_		_		ERd32-2→ERd32,EXR→@ERd	+	F	F	F	F	F	4
	STC CCR,@aa:16	W	-		-			6				CCR→@aa:16	+	F	F	F	F	F	4
	STC EXR,@aa:16	W	├	_	-	_		6		_		EXR→@aa:16	+	F	H	F	E	H	4
	STC CCR,@aa:32	W	-	-	-			8		_		CCR→@aa:32	+	F	F	F	F	F	5
	STC EXR,@aa:32	W	<u> </u>	_	_			8				EXR→@aa:32	丰	Ļ	1	上	Ļ	F	5
ANDC	ANDC #xx:8,CCR	В	2	_	-					_		CCR∧#xx:8→CCR	‡	1	‡	1	1	H	1
	ANDC #xx:8,EXR	В	4		_					_		EXR∧#xx:8→EXR	1	F	F	1	Ļ	F	2
ORC	ORC #xx:8,CCR	В	2	<u> </u>	_	_						CCR√#xx:8→CCR	1	1	H	ļļ	ļļ	ļļ	1
	ORC #xx:8,EXR	В	4		_							EXR√#xx:8→EXR	上	Ļ	F	上	Ļ	Ļ	2
XORC	XORC #xx:8,CCR	В	2		_							CCR⊕#xx:8→CCR	1	1	‡	1	1	‡	1
	XORC #xx:8,EXR	В	4	_								EXR⊕#xx:8→EXR	\perp	\vdash		\vdash	E	\vdash	2
NOP	NOP	_									2	PC←PC+2				E	E		1

• Block transfer instructions

Mnemonic		Operand Size	Addressing Mode/Instruction Length (Bytes)								gth								
			×		@ERn	@(d,ERn)	@-ERn/@ERn+	@aa	@(q,PC)	@aa		Operation	Condition Code					Number of States*1	
			xx#	Ru	0	ø		0		(9)	- 1		ı	н	N	ı z	v	С	Advanced
EEPMOV	EEPMOV.B	_									4	if R4L≠0 Repeat @ER5→@ER6 R5+1→R5 R6+1→R6 R4L-1→R4L Until R4L=0 else next;							4+2n* ²
	EEPMOV.W	_									4	if R4≠0 Repeat @ER5→@ER6 R5+1→R5 R6+1→R6 R4-1→R4 Until R4=0 else next;		_					4+2n* ²

Notes: 1. The number of states is the number of states required for execution when the instruction code and operands are located in on-chip memory.

- 2. n is the initial value of R4L or R4.
- [1] 7 states when the number of restored/saved registers is 2, 9 states when 3, and 11 states when 4.
- [2] Cannot be used with the H8S/2623 Series.
- [3] Set to 1 when there is a carry from or borrow to bit 11; otherwise cleared to 0.
- [4] Set to 1 when there is a carry from or borrow to bit 27; otherwise cleared to 0.
- [5] If the result is zero, the previous value of the flag is retained; otherwise the flag is cleared to 0.
- [6] Set to 1 if the divisor is negative: otherwise cleared to 0.
- [7] Set to 1 if the divisor is zero; otherwise cleared to 0.
- [8] Set to 1 if the quotient is negative; otherwise cleared to 0.
- [9] When EXR is valid, the number of states is increased by 1.

Number of States Required for Execution

The number of states shown in the instruction set table is the number of states required for execution when the op code and operand data are located in a one-cycle area on which word access is possible, such as on-chip memory. When the op code or operand data is accessed from an on-chip supporting module or an external address, the number of states increases as shown in the table below.

• Number of States Required for Execution (Cycle)

Access Conditions

		•			Extern	al Device	
			n-Chip ting Module	8-B	it Bus	16-E	Bit Bus
Execution State (Cycle)	On-Chip Memory	8-Bit Bus	16-Bit Bus	2-State Access	3-State Access	2-State Access	3-State Access
Instruction fetch	1	4	2	4	6 + 2m	2	3 + m
Branch instruction read	_						
Stack operation	_						
Byte data access	_	2		2	3 + m	=	
Word data access	_	4		4	6 + 2m	=	
Internal operation	1	1	1	1	1	1	1

Legend

m: Number of wait states inserted in external device access

Condition Code Notation

Symbol	Meaning
‡	Changes according to operation result.
*	Indeterminate (value not guaranteed).
0	Always cleared to 0.
1	Always set to 1.
_	Not affected by operation result.

Operation Notation

Rd	General register (destination)*					
Rs	General register (source)*					
Rn	General register*					
ERn	General register (32-bit register)					
MAC	Multiply-and-accumulate register (32-bit register)					
(EAd)	Destination operand					
(EAs)	Source operand					
EXR	Extend register					
CCR	Condition-code register					
N	N (negative) flag in CCR					
Z	Z (zero) flag in CCR					
V	V (overflow) flag in CCR					
С	C (carry) flag in CCR					
PC	Program counter					
SP	Stack pointer					
#IMM	Immediate data					
disp	Displacement					
+	Addition					
_	Subtraction					
×	Multiplication					
÷	Division					
^	Logical AND					
<u> </u>	Logical OR					
⊕	Logical exclusive OR					
\rightarrow	Transfer					
7	NOT (logical complement)					
() <>	Operand contents					
:8/:16/:24/:32	8-, 16-, 24-, or 32-bit length					

Note: * General registers include 8-bit registers (R0H to R7H, R0L to R7L), 16-bit registers (R0 to R7, E0 to E7), and 32-bit registers (ER0 to ER7).

2.6 Basic Bus Timing

The CPU operates on the basis of the system clock (Ø). One Ø clock cycle is called a state. The bus cycle consists of one, two, or three states. Different access methods are used for on-chip memory, on-chip supporting modules, and the external address space.

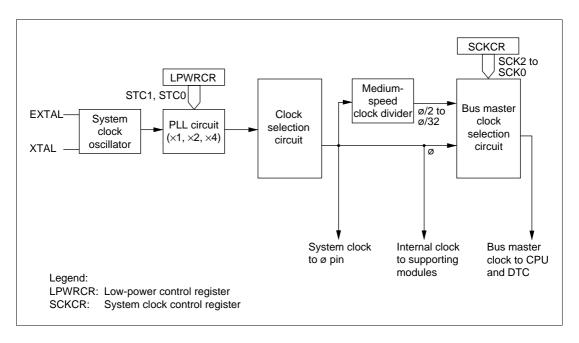
System Clock

The system clock (ø) is generated by the oscillator and PLL circuit.

The following two methods can be used to supply a clock to the oscillator. In either case, the input clock should not exceed 20 MHz.

- 1. Connecting a crystal resonator to the EXTAL and XTAL pins
- 2. Inputting an external clock to the EXTAL pin

The frequency of the clock from the oscillator can be multiplied by a factor of 1, 2, or 4 by means of the PLL circuit. Ensure that, after multiplication, the clock frequency does not exceed the maximum operating frequency of the chip.



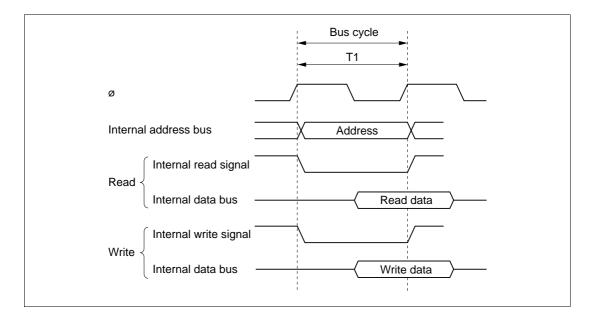
CPU Read/Write Cycles

The CPU operates on the basis of the system clock (ø). One ø clock cycle is called a state, and a bus cycle consists of one, two, or three states. Different access methods are used for on-chip memory, on-chip supporting modules, and external address space. Access to the external address space can be controlled by the bus controller.

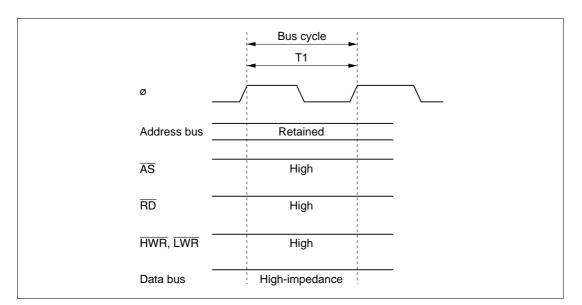
On-Chip Memory

On-chip memory is accessed in one state. The data bus is 16 bits wide, permitting both byte and word access.

• On-Chip Memory Access Cycle (One-State Access)



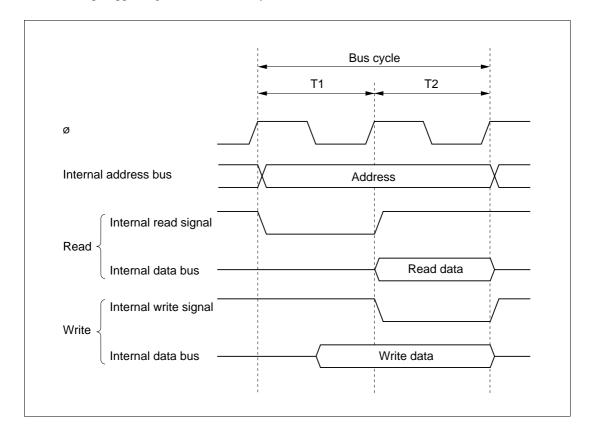
• Pin States during On-Chip Memory Access



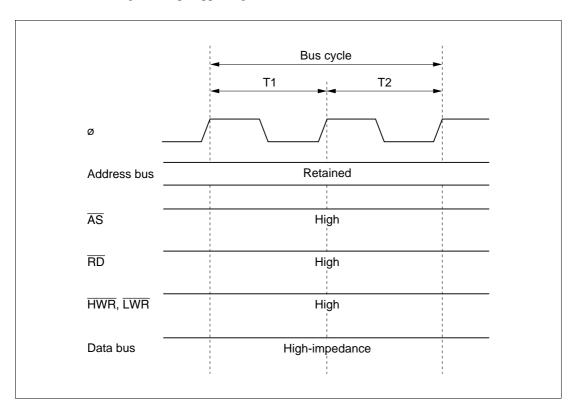
On-Chip Supporting Modules

The on-chip supporting modules are accessed in two states. The data bus is either 8 bits or 16 bits wide, depending on the internal I/O register.

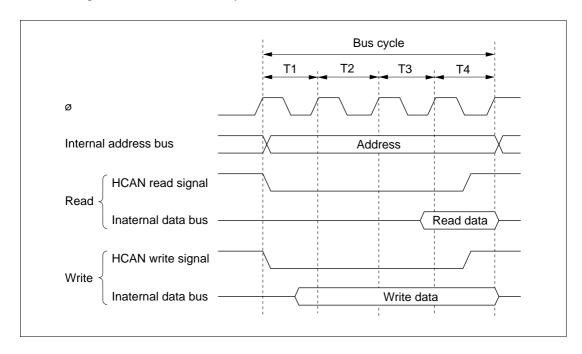
• On-Chip Supporting Module Access Cycle (Two-State Access)



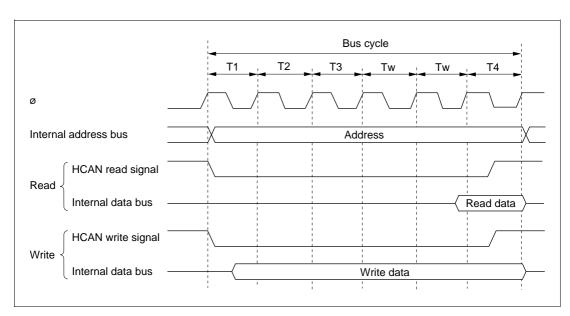
• Pin States during On-Chip Supporting Module Access



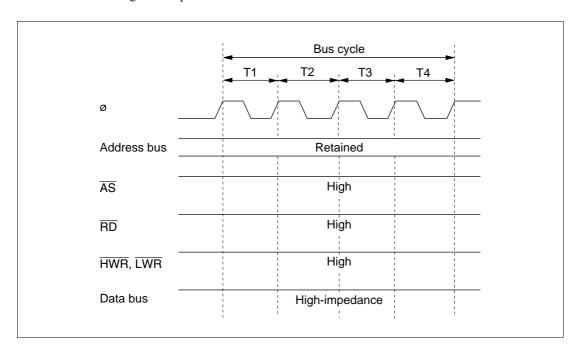
• On-Chip HCAN Module Access Cycle (No Wait States)



• On-Chip HCAN Module Access Cycle (Wait States Inserted)



• Pin States during On-Chip HCAN Module Access



External Address Space

The external address space is accessed with an 8-bit or 16-bit data bus width in a two-state or three-state bus cycle. In three-state access, wait states can be inserted. For further details, refer to section 3.1, Bus Controller (BSC).

2.7 Processing States

The H8S/2600 CPU has five processing states: the reset state, program execution state, exception-handling state, bus-released state, and power-down state.

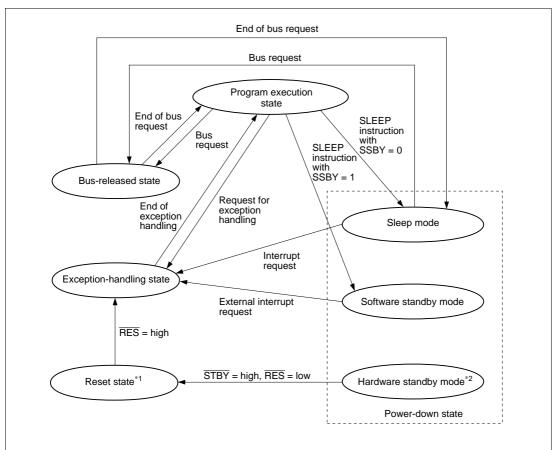
Reset State: State in which the CPU and all on-chip supporting modules are initialized and halted.

Program Execution State: State in which the CPU executes the program sequentially.

Exception-Handling State: Transient state in which the normal processing flow is altered and exception handling executed as the result of a reset, interrupt, or trap instruction exception handling source.

Bus-Released State: State in which the external bus is released in response to a bus request signal from a bus master other than the CPU.

Power-Down State: State in which CPU operation is stopped, and power consumption is kept low (sleep mode, software standby mode, hardware standby mode). The power-down state also includes medium-speed mode and module stop mode.



Notes: 1. From any state except hardware standby mode, a transition to the power-on reset state occurs whenever RES goes low. A transition can also be made to the reset state when the watchdog timer overflows.

2. From any state, a transition to hardware standby mode occurs whenever STBY goes low.

2.8 Exception Handling

H8S/2600 CPU exception handling is activated by a reset, a trap instruction, or an interrupt. A priority system is provided for exception handling, and simultaneously generated exceptions are handled in order of priority.

Exception Types and Priority

Priority	Exception Type	Start of Exception Handling
High	Reset	Starts immediately after a low-to-high transition at the $\overline{\text{RES}}$ pin, or when the watchdog timer overflows. When the $\overline{\text{RES}}$ pin is low, the power-on reset state is in effect.
	Trace*1	Starts when execution of the current instruction or exception handling ends, if the trace (T) bit is set to 1.
	Interrupt	Starts when execution of the current instruction or exception handling ends, if an interrupt request has been issued.*2
Low	Trap instruction*3 (TRAPA)	Started by execution of a trap instruction (TRAPA).

Notes: 1. The trace function is enabled only in interrupt control mode 2. Trace exception handling is not executed after execution of an RTE instruction.

- 2. Interrupt detection is not performed at the end of execution of an ANDC, ORC, XORC, or LDC instruction, or at the end of reset exception handling.
- 3. Trap instruction exception handling is always accepted in the program execution state.

Exception Handling Operation

Exception handling is initiated by various exception handling sources. Trap instruction exception handling is always accepted in the program execution state. Trap instructions and interrupts are handled as follows:

- 1. The program counter (PC), condition-code register (CCR), and extend register (EXR) are pushed onto the stack.
- 2. The interrupt mask bits are updated. The T bit is cleared to 0.
- 3. A vector address corresponding to the activation source is generated, and program execution starts from the address indicated in that vector address.

For a reset exception, steps 2 and 3 above are carried out.

Exception Vector Table

			Vector Address*1	
Exception Source		Vector Number	Advanced Mode	
Power-on reset		0	H'0000 to H'0003	
Manual reset*3		1	H'0004 to H'0007	
Reserved for system	use	2	H'0008 to H'000B	
		3	H'000C to H'000F	
		4	H'0010 to H'0013	
Trace		5	H'0014 to H'0017	
Reserved for system	use	6	H'0018 to H'001B	
External interrupt	NMI	7	H'001C to H'001F	
Trap instruction (4 so	ources)	8	H'0020 to H'0023	
		9	H'0024 to H'0027	
		10	H'0028 to H'002B	
		11	H'002C to H'002F	
Reserved for system	use	12	H'0030 to H'0033	
		13	H'0034 to H'0037	
		14	H'0038 to H'003B	
		15	H'003C to H'003F	
External interrupt	IRQ0	16	H'0040 to H'0043	
	IRQ1	17	H'0044 to H'0047	
	IRQ2	18	H'0048 to H'004B	
	IRQ3	19	H'004C to H'004F	
	IRQ4	20	H'0050 to H'0053	
	IRQ5	21	H'0054 to H'0057	
Reserved for system	use	22	H'0058 to H'005B	
		23	H'005C to H'005F	
Internal interrupt*2		24	H'0060 to H'0063	
		to 127	to H'01FC to H'01FF	

Notes: 1. Lower 16 bits of the address.

- 2. For details of internal interrupt vectors, see section 2.9, Interrupts.
- 3. Cannot be used with this series.

2.9 Interrupts

Interrupts are controlled by the interrupt controller. There are a total of 51 interrupt sources, comprising seven external interrupts from external pins (NMI, $\overline{IRQ0}$ to $\overline{IRQ5}$), and 44 internal interrupts from on-chip supporting modules (including options). A separate vector number is assigned to each interrupt.

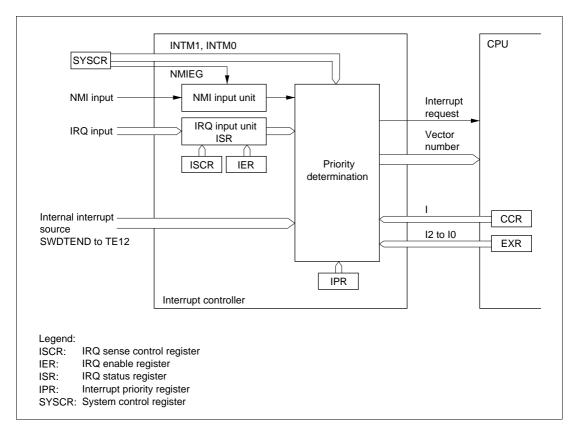
Interrupt Control

Either of two interrupt control modes can be set by means of the INTM1 and INTM0 bits in the system control register (SYSCR).

The interrupt controller controls interrupts on the basis of the control mode set by the INTM1 and INTM0 bits, the interrupt priorities set by the interrupt priority register (IPR), and the masking conditions set by the I bit in CCR and bits I2 to I0 in EXR.

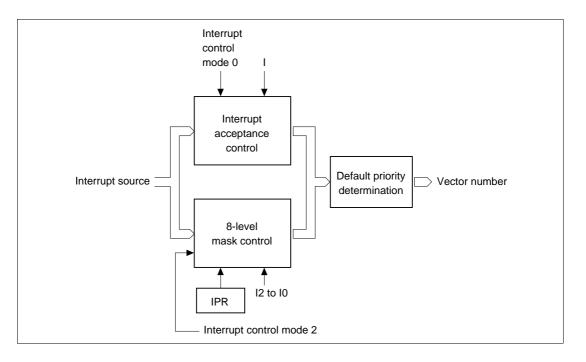
NMI is the highest-priority interrupt, and is always accepted.

Block Diagram of Interrupt Controller



Interrupt	SYSCR		Priority Setting	Interrupt	
Control Mode	INTM1	INTM0	Registers	Mask Bits	Description
0	0	0	_	1	Interrupt mask control is performed by the I bit.
_	-	1	_	_	Setting prohibited
2	1	0	IPR	l2 to l0	8-level interrupt mask control is performed by bits I2 to I0. 8 priority levels can be set with IPR.
_	_	1	_	_	Setting prohibited

• Block Diagram of Interrupt Control Operation



Interrupt Control Mode 0: Enabling and disabling of IRQ interrupts and on-chip supporting module interrupts can be set by means of the I bit in CCR. Interrupts are enabled when the I bit is cleared to 0, and disabled when set to 1.

Interrupt Control Mode 2: Eight-level masking can be implemented for IRQ interrupts and onchip supporting module interrupts by comparing the interrupt mask level bits (I2 to I0) in EXR and the IPR priority level.

Interrupt Sources, Vector Addresses, and Interrupt Priorities

	Origin of		Vector Address*		
Interrupt Source	Interrupt Source	Vector Number	Advanced Mode	- IPR	Priority
NMI	External	7	H'001C		High
IRQ0	pin	16	H'0040	IPRA6 to 4	_
IRQ1		17	H'0044	IPRA2 to 0	_
IRQ2		18	H'0048	IPRB6 to 4	_
IRQ3		19	H'004C		
IRQ4		20	H'0050	IPRB2 to 0	
IRQ5		21	H'0054		
Reserved	_	22	H'0058	IPRC6 to 4	_
		23	H'005C		
SWDTEND (software activated data transfer end)	DTC	24	H'0060	IPRC2 to 0	
WOVI0 (interval timer)	Watchdog timer 0	25	H'0064	IPRD6 to 4	_
Reserved	_	26	H'0068	IPRD2 to 0	_
PC break	PC break controller	27	H'006C	IPRE6 to 4	_
ADI (A/D conversion end)	A/D	28	H'0070	IPRE2 to 0	_
Reserved	_	29	H'0074		_
Reserved	_	30	H'0078		_
		31	H'007C		
TGI0A (TGR0A input capture/ compare match)	TPU channel 0	32	H'0080	IPRF6 to 4	_
TGI0B (TGR0B input capture/ compare match)		33	H'0084		
TGI0C (TGR0C input capture/ compare match)		34	H'0088		
TGI0D (TGR0D input capture/ compare match)		35	H'008C		
TCIOV (overflow 0)		36	H'0090		
Reserved	_	37	H'0094		
		38	H'0098		
		39	H'009C		Low

Note: * Lower 16 bits of the start address.

Interrupt Source		Origin of		Vector Address*	_	
compare match) channel 1 TGI1B (TGR1B input capture/ compare match) 41 H'00A4 TCI1V (overflow 1) 42 H'00A8 TCI1U (underflow 1) 43 H'00AC TGI2A (TGR2A input capture/ compare match) TPU 44 H'00B0 IPRG6 to 4 TGI2B (TGR2B input capture/ compare match) 45 H'00B4 H'00B4 TCI2V (overflow 2) 46 H'00B8 H'00B0 TCI2U (underflow 2) 47 H'00BC IPRG2 to 0 TGI3A (TGR3A input capture/ compare match) TPU 48 H'00C4 IPRG2 to 0 TGI3C (TGR3C input capture/ compare match) 50 H'00C8 H'00C8 IPRG2 to 0 TGI3D (TGR3D input capture/ compare match) 51 H'00C8 IPRG2 to 0 IPRG2 to 0 TGI3D (TGR3D input capture/ compare match) 51 H'00C8 IPRG2 to 0 IPRG2 to 0 TGI3A (TGR3A input capture/ compare match) 52 H'00D0 IPRH6 to 4 IPRH6 to 4 IPRH6 to 4 IPRH6 to 4 IPRH6 to 4 IPRH6 to 4 IPRH6 to 4 IPRH6 to 4 IPRH6 to 4	Interrupt Source	Interrupt Source	Vector Number	Advanced Mode	IPR	Priority
TGI1B (TGR1B input capture/ compare match) TCI1V (overflow 1) TCI1U (underflow 1) TGI2A (TGR2A input capture/ channel 2 TGI2B (TGR2B input capture/ channel 2 TGI2B (TGR2B input capture/ channel 2 TGI2B (TGR3B input capture/ compare match) TCI2V (overflow 2) TGI2U (underflow 2) TGI3A (TGR3A input capture/ channel 3 TGI3B (TGR3B input capture/ compare match) TGI3C (TGR3C input capture/ channel 3 TGI3C (TGR3C input capture/ compare match) TGI3C (TGR3C input capture/ compare match) TGI3C (TGR3C input capture/ compare match) TGI3C (TGR3C input capture/ compare match) TGI3C (TGR3C input capture/ compare match) TGI3V (overflow 3) Reserved TPU 56 H'00D0 TGI3A (TGR4A input capture/ compare match) TGI4A (TGR4A input capture/ channel 4 TGI4B (TGR4B input capture/ compare match) TGI4C (overflow 4) TGI4C (overflow 4) TGI4C (overflow 4) TGI4C (TGR5A input capture/ compare match) TGI4C (overflow 4) TGIAC (TGR5A input capture/ compare match)			40	H'00A0	IPRF2 to 0	High
Compare match TCI1V (overflow 1)		channel 1				A
TCI1U (underflow 1)			41	H'00A4		
TGI2A (TGR2A input capture/ compare match) TPU channel 2 44 H'00B0 IPRG6 to 4 TGI2B (TGR2B input capture/ compare match) 45 H'00B4 TCI2V (overflow 2) 46 H'00BC TCI2U (underflow 2) 47 H'00BC TGI3A (TGR3A input capture/ compare match) TPU channel 3 48 H'00C0 IPRG2 to 0 TGI3B (TGR3B input capture/ compare match) 50 H'00C8 H'00C8 H'00C8 TGI3D (TGR3C input capture/ compare match) 51 H'00CC H'00D0 H'00D0 H'00D0 IPRH6 to 4 H'00D0 IPRH6 to 4 TGI3A (TGR4A input capture/ compare match) TPU channel 4 56 H'00E0 IPRH6 to 4 H'00E4 IPRH6 to 4 TGI4A (TGR4A input capture/ compare match) 57 H'00E4 H'00E0 IPRH6 to 4 TGI4A (TGR5A input capture/ compare match) 59 H'00EC TGI5A (TGR5A input capture/ compare match) 60 H'00F0 IPRH2 to 0 TGI5B (TGR5B input capture/ compare match) 61 H'00F4 H'00F4 H'00F4 H'00F4 H'00F4 H'00F4 H'00F4 H'00F4 TGI5D (Voerflow 5)	TCI1V (overflow 1)		42	H'00A8		
compare match) channel 2 TGI2B (TGR2B input capture/ compare match) 45 H'00B4 TCI2V (overflow 2) 46 H'00B8 TCI2U (underflow 2) 47 H'00BC TGI3A (TGR3A input capture/ compare match) TPU 48 H'00C0 IPRG2 to 0 TGI3B (TGR3B input capture/ compare match) 50 H'00C8 H'00C8 TGI3D (TGR3C input capture/ compare match) 51 H'00CC TGI3V (overflow 3) 52 H'00D0 Reserved — 53 H'00D4 54 H'00D8 55 H'00DC TGI4A (TGR4A input capture/ compare match) TPU 56 H'00E0 IPRH6 to 4 TGI4B (TGR4B input capture/ compare match) 57 H'00E4 H'00E4 TCI4V (overflow 4) 59 H'00EC TGI5A (TGR5A input capture/ compare match) 59 H'00F0 IPRH2 to 0 TGI5B (TGR5B input capture/ compare match) Channel 5 61 H'00F4 H'00F4	TCI1U (underflow 1)		43	H'00AC		
Compare match TCI2V (overflow 2)			44	H'00B0	IPRG6 to 4	
TCI2U (underflow 2)	· · · · · ·		45	H'00B4		
TGI3A (TGR3A input capture/ compare match) TPU channel 3 48 H'00C0 IPRG2 to 0 TGI3B (TGR3B input capture/ compare match) 49 H'00C4 H'00C8 TGI3C (TGR3C input capture/ compare match) 50 H'00C8 H'00CC TGI3D (TGR3D input capture/ compare match) 51 H'00CC H'00D0 Reserved — 53 H'00D4 54 H'00D8 55 H'00DC TGI4A (TGR4A input capture/ compare match) TPU channel 4 56 H'00E0 IPRH6 to 4 TGI4B (TGR4B input capture/ compare match) 58 H'00E4 H'00E4 TCI4V (overflow 4) 59 H'00EC IPRH2 to 0 TGI5A (TGR5A input capture/ compare match) TPU channel 5 60 H'00F0 IPRH2 to 0 TGI5B (TGR5B input capture/ compare match) 61 H'00F4 H'00F4 H'00F4	TCI2V (overflow 2)		46	H'00B8		
compare match) channel 3 TGI3B (TGR3B input capture/compare match) 49 H'00C4 TGI3C (TGR3C input capture/compare match) 50 H'00C8 TGI3D (TGR3D input capture/compare match) 51 H'00CC TCI3V (overflow 3) 52 H'00D0 Reserved — 53 H'00D4 54 H'00D8 55 H'00DC TGI4A (TGR4A input capture/compare match) TPU 56 H'00E0 IPRH6 to 4 TGI4B (TGR4B input capture/compare match) 57 H'00E4 57 H'00E8 TCI4V (overflow 4) 59 H'00EC H'00F0 IPRH2 to 0 TGI5A (TGR5A input capture/compare match) Channel 5 61 H'00F4 H'00F4 TGI5B (TGR5B input capture/compare match) 61 H'00F4 H'00F4 H'00F4 TGI5V (overflow 5) 62 H'00F8 H'00F8 H'00F4	TCI2U (underflow 2)		47	H'00BC		
compare match) 50 H'00C8 TGI3C (TGR3C input capture/ compare match) 51 H'00CC TGI3D (TGR3D input capture/ compare match) 51 H'00D0 TCI3V (overflow 3) 52 H'00D4 Reserved — 53 H'00D4 54 H'00D8 55 H'00DC TGI4A (TGR4A input capture/ compare match) TPU channel 4 56 H'00E0 IPRH6 to 4 TGI4B (TGR4B input capture/ compare match) 57 H'00E4 H'00E4 TCI4V (overflow 4) 59 H'00EC TGI5A (TGR5A input capture/ compare match) TPU compare match) 60 H'00F0 IPRH2 to 0 TGI5B (TGR5B input capture/ compare match) 61 H'00F4 H'00F4 H'00F4 TCI5V (overflow 5) 62 H'00F8 H'00F8 H'00F4 H'00F8			48	H'00C0	IPRG2 to 0	
compare match) TGI3D (TGR3D input capture/ compare match) 51 H'00CC TCI3V (overflow 3) 52 H'00D0 Reserved — 53 H'00D4 54 H'00D8 55 H'00DC TGI4A (TGR4A input capture/ compare match) TPU channel 4 56 H'00E0 IPRH6 to 4 TGI4B (TGR4B input capture/ compare match) 57 H'00E4 H'00E8 TCI4V (overflow 4) 59 H'00EC TGI5A (TGR5A input capture/ compare match) TPU channel 5 60 H'00F0 IPRH2 to 0 TGI5B (TGR5B input capture/ compare match) 61 H'00F4 H'00F4 TCI5V (overflow 5) 62 H'00F8			49	H'00C4		
TCI3V (overflow 3) 52			50	H'00C8		
TGI4A (TGR4A input capture/ compare match)			51	H'00CC		
54	TCI3V (overflow 3)		52	H'00D0		
TGI4A (TGR4A input capture/ compare match) TGI4B (TGR4B input capture/ compare match) TCI4V (overflow 4) TCI4U (underflow 4) TGI5A (TGR5A input capture/ compare match) TGI5B (TGR5B input capture/ compare match) TGI5B (TGR5B input capture/ compare match) TGI5V (overflow 5) 55 H'00E0 IPRH6 to 4 FR H'00E4 FR H'00E8 H'00EC IPRH2 to 0 FR H'00F6 FR H'00F6 IPRH2 to 0 FR H'00F6 FR H'00F8	Reserved	_	53	H'00D4		_
TGI4A (TGR4A input capture/ channel 4 TPU 56 H'00E0 IPRH6 to 4 compare match) TGI4B (TGR4B input capture/ compare match) TCI4V (overflow 4) 58 H'00E8 TCI4U (underflow 4) 59 H'00EC TGI5A (TGR5A input capture/ compare match) TGI5B (TGR5B input capture/ channel 5 TGI5B (TGR5B input capture/ compare match) TCI5V (overflow 5) 62 H'00F8			54	H'00D8		
compare match) TGI4B (TGR4B input capture/ compare match) TCI4V (overflow 4) TCI4U (underflow 4) TGI5A (TGR5A input capture/ compare match) TGI5B (TGR5B input capture/ channel 5 TGI5B (TGR5B input capture/ compare match) TCI5V (overflow 5) Channel 4 H'00E4 TOUE TOU			55	H'00DC		
compare match) 58 H'00E8 TCI4V (overflow 4) 59 H'00EC TGI5A (TGR5A input capture/ compare match) TPU 60 H'00F0 IPRH2 to 0 TGI5B (TGR5B input capture/ compare match) 61 H'00F4 TCI5V (overflow 5) 62 H'00F8			56	H'00E0	IPRH6 to 4	
TCI4U (underflow 4) TGI5A (TGR5A input capture/ compare match) TGI5B (TGR5B input capture/ compare match) TGI5V (overflow 5) TGI5U (underflow 4) TPU 60 H'00F0 IPRH2 to 0 Channel 5 61 H'00F4 CH 100F8	· · · · · ·		57	H'00E4		
TGI5A (TGR5A input capture/ TPU 60 H'00F0 IPRH2 to 0 compare match) channel 5 TGI5B (TGR5B input capture/ 61 H'00F4 compare match) TCI5V (overflow 5) 62 H'00F8	TCI4V (overflow 4)		58	H'00E8		
compare match) channel 5 TGI5B (TGR5B input capture/ 61 H'00F4 compare match) TCI5V (overflow 5) 62 H'00F8	TCI4U (underflow 4)		59	H'00EC		
compare match) TCI5V (overflow 5) 62 H'00F8			60	H'00F0	IPRH2 to 0	
			61	H'00F4		
TCI5U (underflow 5) 63 H'00FC Low	TCI5V (overflow 5)		62	H'00F8		
	TCI5U (underflow 5)		63	H'00FC		Low

Note: * Lower 16 bits of the start address.

	Origin of		Vector Address*		
Interrupt Source	Interrupt Source	Vector Number	Advanced Mode	- IPR	Priority
Reserved	_	64	H'0100	IPRI6 to 4	High
		65	H'0104		A
		66	H'0108		
		67	H'010C		
		68	H'0110	IPRI2 to 0	
		69	H'0114		
		70	H'0118		
		71	H'011C		
		72	H'0120	IPRJ6 to 4	
		73	H'0124		
		74	H'0128		
		75	H'012C		
		76	H'0130		
		77	H'0134		
		78	H'0138		
		79	H'013C		
ERI0 (receive error 0)	SCI	80	H'0140	IPRJ2 to 0	
RXI0 (reception completed 0)	channel 0	81	H'0144		
TXI0 (transmit data empty 0)		82	H'0148		
TEI0 (transmission end 0)		83	H'014C		
ERI1 (receive error 1)	SCI	84	H'0150	IPRK6 to 4	
RXI1 (reception completed 1)	channel 1	85	H'0154		
TXI1 (transmit data empty 1)		86	H'0158		
TEI1 (transmission end 1)		87	H'015C		
ERI2 (receive error 2)	SCI	88	H'0160	IPRK2 to 0	
RXI2 (reception completed 2)	channel 2	89	H'0164		
TXI2 (transmit data empty 2)		90	H'0168		
TEI2 (transmission end 2)		91	H'016C		
ERS0	HCAN	104	H'01A0	IPRM6 to 4	
OVR0		105	H'01A4		
RM0		106	H'01A8		
SLE0		107	H'01AC		
RM0		108	H'01B0	IPRM2 to 0	Low

Note: * Lower 16 bits of the start address.

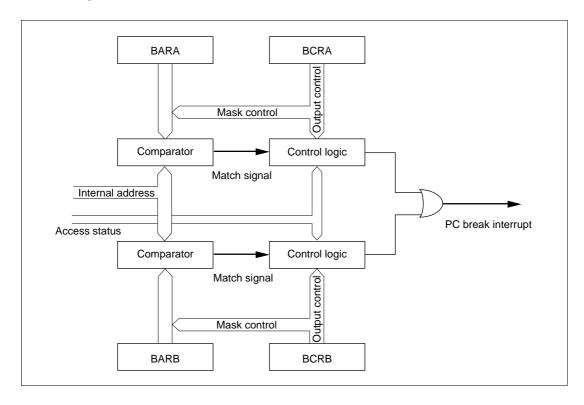
2.10 PC Break Controller (PBC)

This series has a two-channel on-chip PC break controller (PBC) providing functions that simplify program debugging. Using these functions, it is easy to create a sophisticated self-monitoring debugger, enabling programs to be debugged with the chip alone, without using a large-scale incircuit emulator.

Features

- Two break channels (A and B)
- The following can be set as break compare conditions:
 - 24 address bits
 - Bit masking possible
 - Bus cycle
 - Instruction fetch
 - Data access: data read, data write, data read/write
 - Bus master
 - Either CPU or CPU/DTC can be selected
- The timing of PC break exception handling after the occurrence of a break condition is as follows:
 - Immediately before execution of the instruction fetched at the set address (instruction fetch)
 - Immediately after execution of the instruction that accesses data at the set address (data access)
- Module stop mode can be set
 - As the initial setting, PBC operation is halted. Register access is enabled by exiting module stop mode.

Block Diagram of PBC



2.11 Operating Modes

In this series, there are four operating modes. The operating mode is determined by the settings of the mode pins (MD2 to MD0).

Mode 4: The CPU can access a 16-Mbyte address space in advanced mode. The on-chip ROM is disabled. Ports 1, A, B, and C function as an address bus, ports D and E function as a data bus, and part of port F carries bus control signals.

The initial bus mode after a reset is 16 bits, with 16-bit access to all areas. However, if 8-bit access is designated by the bus controller for all areas, the bus mode switches to 8 bits.

Mode 5: The CPU can access a 16-Mbyte address space in advanced mode. The on-chip ROM is disabled. Ports 1, A, B, and C function as an address bus, port D functions as a data bus, and part of port F carries bus control signals.

The initial bus mode after a reset is 8 bits, with 8-bit access to all areas. However, if any area is designated for 16-bit access by the bus controller, the bus mode switches to 16 bits and port E becomes a data bus.

Mode 6: The CPU can access a 16-Mbyte address space in advanced mode. The on-chip ROM is enabled. Ports 1, A, B, and C function as input ports immediately after a reset. These pins can be set to output addresses by setting the corresponding data direction register (DDR) bits to 1. Port D functions as a data bus, and part of port F carries bus control signals.

The initial bus mode after a reset is 8 bits, with 8-bit access to all areas. However, if any area is designated for 16-bit access by the bus controller, the bus mode switches to 16 bits and port E becomes a data bus.

Mode 7: The CPU can access a 16-Mbyte address space in advanced mode. The on-chip ROM is enabled, but external addresses cannot be accessed.

All I/O ports are available for use as I/O ports.

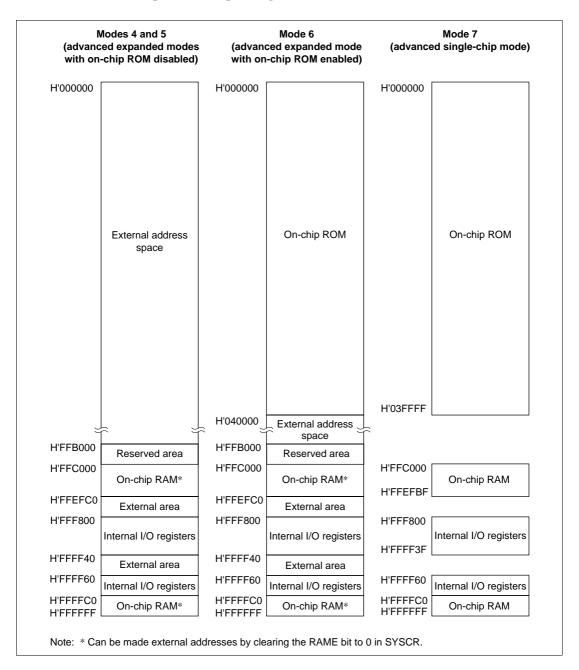
MCU Operating Modes

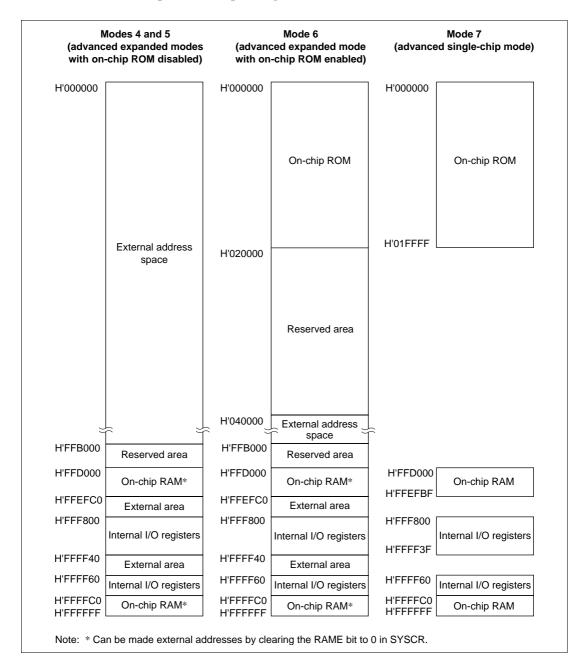
MCU				CPU				nal Data Bus
Operating Mode	MD2	MD1	MD0	Operating Mode	Description	On-Chip ROM	Initial Width	Max. Width
4	1	0	0	Advanced	Expanded mode with	Disabled	16 bits	16 bits
5			1		on-chip ROM disabled		8 bits	16 bits
6		1	0		Expanded mode with on-chip ROM enabled	Enabled	8 bits	16 bits
7	_		1	_	Single-chip mode		_	_

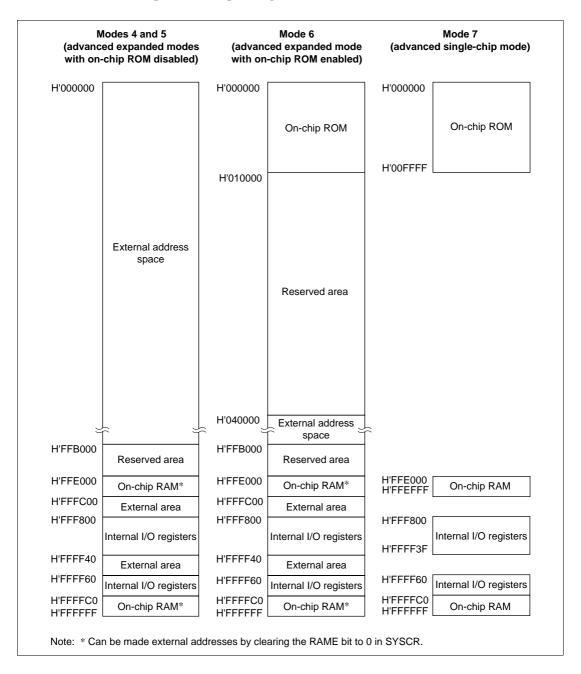
2.12 Address Maps

The address space is 16 Mbytes in modes 4 to 7 (advanced modes).

H8S/2623 Address Maps in Each Operating Mode







Section 3 Supporting Modules

3.1 Bus Controller (BSC)

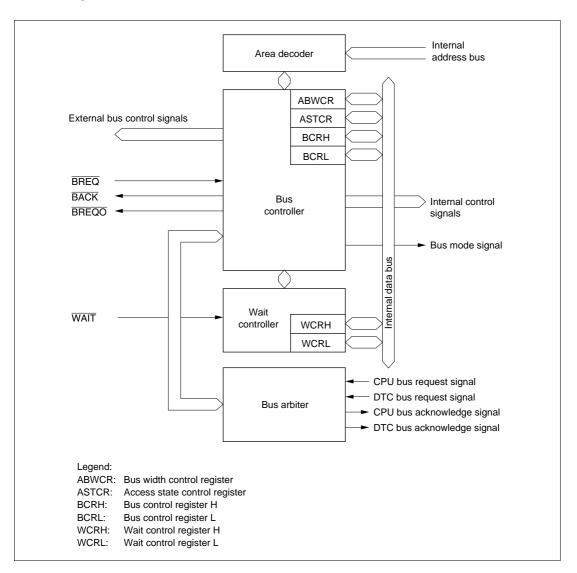
The H8S/2623 Series has a built-in bus controller (BSC) that manages the external address space divided into eight areas. The bus specifications, such as bus width and number of access states, can be set independently for each area, enabling multiple memories to be connected easily.

The bus controller also has a bus arbitration function, and controls the operation of the internal bus masters: the CPU and data transfer controller (DTC).

Features

- Manages external address space in area units
 - Manages the external space as 8 areas of 2 Mbytes
 - Bus specifications can be set independently for each area
 - Burst ROM interface can be set
- · Basic bus interface
 - 8-bit access or 16-bit access can be selected for each area
 - 2-state access or 3-state access can be selected for each area
 - Program wait states can be inserted for each area
- Burst ROM interface
 - Burst ROM interface can be set for area 0
 - Choice of 1- or 2-state burst access
- Idle cycle insertion
 - An idle cycle can be inserted in case of an external read cycle between different areas
 - An idle cycle can be inserted in case of an external write cycle immediately after an external read cycle
- Write buffer function
 - External write cycle and internal access can be executed in parallel
- Bus arbitration function
 - Includes a bus arbiter that arbitrates bus mastership between the CPU and DTC
- Other features
 - External bus release function

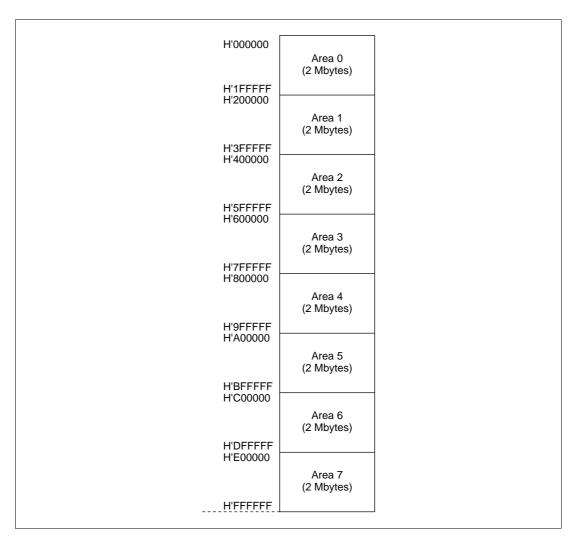
Block Diagram of Bus Controller



3.1.1 Area Partitioning

The bus controller partitions the 16-Mbyte address space into eight areas, 0 to 7, in 2-Mbyte units, and performs bus control for external space in area units.

Overview of Area Partitioning



Bus Specifications

The external address space bus specifications consist of three elements: bus width, number of access states, and number of program wait states. The bus width and number of access states for on-chip memory and internal I/O registers are fixed, and are not affected by the bus controller.

Bus specifications can be set as shown below by means of the bus controller control registers.

• Bus Specifications for Each Area (Basic Bus Interface)

ABWCR	ASTCR	WCR	H, WCRL	Bus Specifications (Basic Bus Interface)			
ABWn	ASTn	Wn1	Wn0	Bus Width	Access States	Program Wait States	
0	0	_	_	16	2	0	
	1	0	0	•	3	0	
			1	-		1	
		1	0	•		2	
			1	-		3	
1	0	_	_	8	2	0	
	1	0	0	•	3	0	
			1	-		1	
		1	0	•		2	
			1	-		3	

Memory Interfaces

This series' memory interfaces comprise a basic bus interface that allows direct connection of ROM, SRAM, and so on, and a burst ROM interface that allows direct connection of burst ROM. The interface can be designated independently for each area.

An area for which the basic bus interface is designated is ordinary space, and on area for which the burst ROM interface is designated is burst ROM space.

3.1.2 Basic Bus Interface

The basic bus interface allows direct connection of ROM, SRAM, etc.

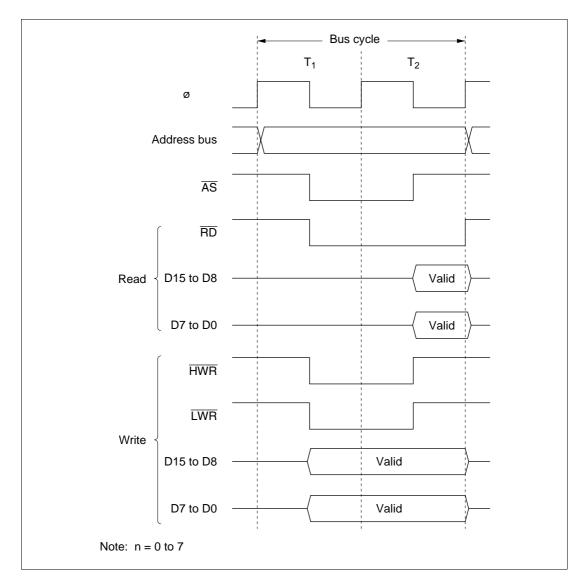
The bus specifications can be selected by means of ABWCR, ASTCR, WCRH, and WCRL.

This interface can be designated for areas 0 to 7.

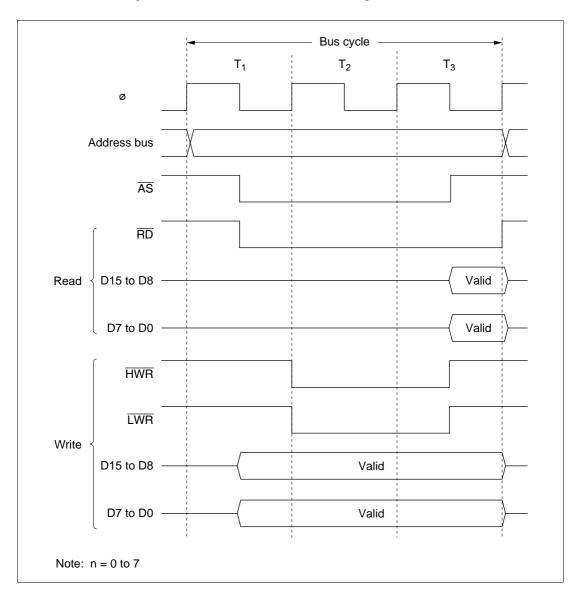
In 3-state access space, 0 to 3 program wait states or a pin wait by means of the \overline{WAIT} pin can be inserted.

After a reset, all areas are designated as basic bus interface, 3-state access space (the bus width is determined by the MCU operating mode).

• Basic Bus Timing (Word Access to 16-Bit 2-State Access Space)



• Basic Bus Timing (Word Access to 16-Bit 3-State Access Space)

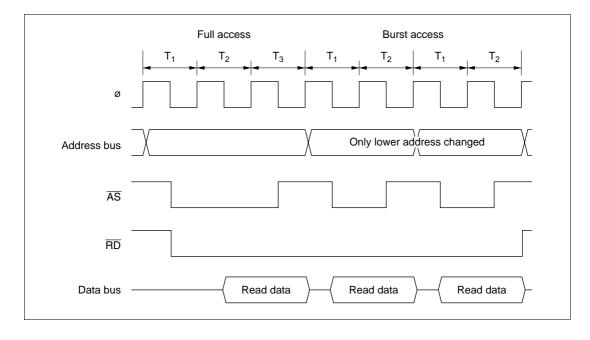


3.1.3 Burst ROM Interface

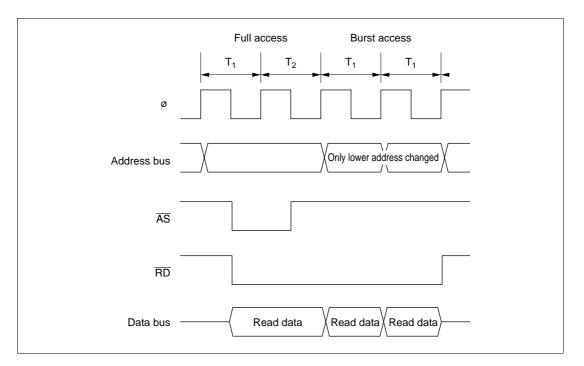
External space area 0 can be designated as burst ROM space, and burst ROM space interfacing can be performed. The burst ROM space interface enables 16-bit configuration ROM with burst access capability to be accessed at high speed.

Consecutive burst accesses of a maximum 4 words or 8 words can be performed for CPU instruction fetches only. One or two states can be selected for burst access.

• Example of Burst ROM Access Timing (When AST0 = BRSTS1 = 1)



• Example of Burst ROM Access Timing (When AST0 = BRSTS1 = 0)



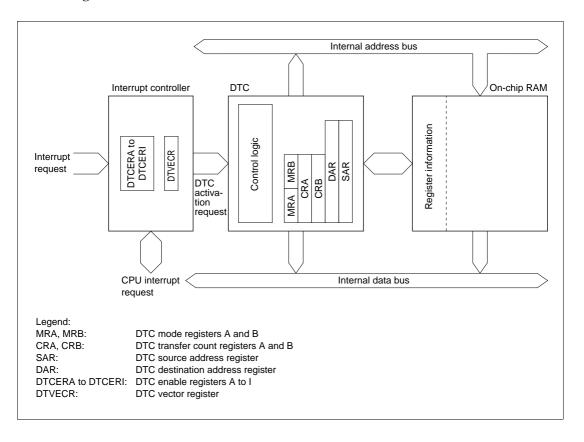
3.2 Data Transfer Controller (DTC)

The data transfer controller (DTC) is activated by an interrupt or software, and can transfer data without imposing any load on the CPU.

Features

- Transfer possible over any number of channels
 - Transfer information is stored in memory
 - One activation source can trigger a number of data transfers (chain transfer)
- · Variety of transfer modes
 - Normal, repeat, and block transfer modes available
 - Incrementing, decrementing, or fixing of source and destination addresses can be selected
- Direct specification of 16-Mbyte address space possible
 - 24-bit specification of transfer source and destination addresses
- Transfer can be set in byte or word units
- A CPU interrupt can be requested for the interrupt that activated the DTC
 - An interrupt request can be issued to the CPU after one data transfer ends
 - An interrupt request can be issued to the CPU after all specified data transfers have ended
- Can be activated by software
- Module stop mode can be set
 - The initial setting enables DTC registers to be accessed. DTC operation is halted by setting module stop mode.

Block Diagram of DTC

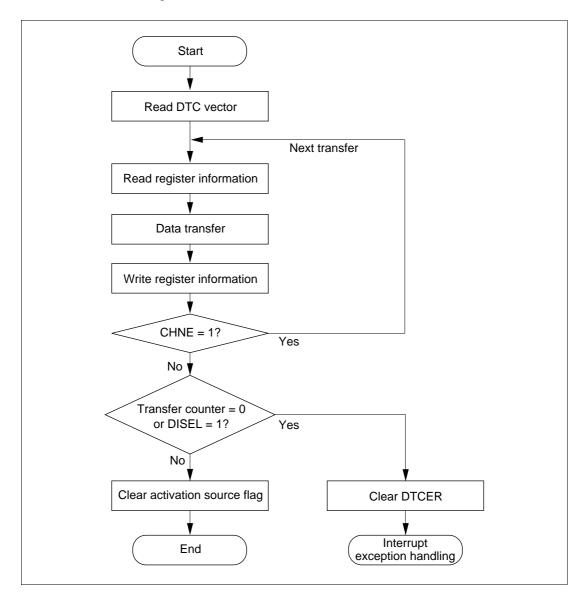


Data Transfer Operation

When a DTC activation source occurs, the DTC reads register information previously stored in memory, and transfers data on the basis of that register information. After the data transfer, it writes updated register information back to memory.

Pre-storage of register information in memory makes it possible to transfer data over any required number of channels. The DTC can also execute a number of transfers with a single activation source (chain transfer).

• Flowchart of DTC Operation

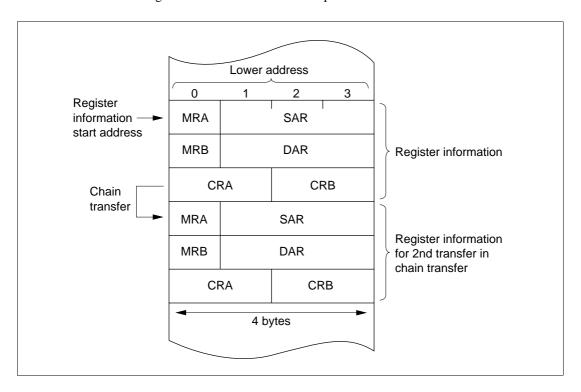


DTC Activation Sources

The DTC is activated by an interrupt source or by a vector number write to the DTC vector register (DTVECR) by software. An interrupt request can be designated as a CPU interrupt source or a DTC activation source.

When an interrupt has been designated as a DTC activation source, existing CPU mask level and interrupt controller priorities have no effect. If there is more than one activation source at the same time, the DTC operates in accordance with the default priorities.

- Interrupt Sources and DTC Vector Address
 - The DTC vector address indicates the start address of the register information in memory. The MRA, SAR, MRB, DAR, CRA, and CRB registers are located in that order from the start address of the register information. Locate the register information in the on-chip RAM (addresses H'FFEBC0 to H'FFEFBF).
- Location of DTC Register Information in Address Space



• Interrupt Sources, DTC Vector Addresses, and Corresponding DTCEs

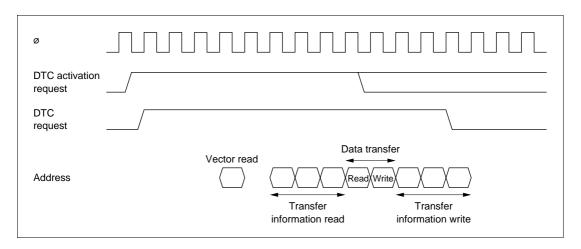
Interrupt Source	Origin of Interrupt Source	Vector Number	Vector Address	DTCE*	Priority
Write to DTVECR	Software	DTVECR	H'0400+ (DTVECR [6:0]<<1)	_	High A
IRQ0	External pin	16	H'0420	DTCEA7	_
IRQ1		17	H'0422	DTCEA6	_
IRQ2		18	H'0424	DTCEA5	_
IRQ3		19	H'0426	DTCEA4	_
IRQ4		20	H'0428	DTCEA3	
IRQ5		21	H'042A	DTCEA2	
Reserved for system use	_	22	H'042C	DTCEA1	_
		23	H'042E	DTCEA0	_
ADI (A/D conversion end)	A/D	28	H'0438	DTCEB6	
TGI0A (GR0A compare match/input capture)	TPU channel 0	32	H'0440	DTCEB5	
TGI0B (GR0B compare match/input capture)		33	H'0442	DTCEB4	_
TGI0C (GR0C compare match/ input capture)		34	H'0444	DTCEB3	
TGI0D (GR0D compare match/ input capture)		35	H'0446	DTCEB2	_
TGI1A (GR1A compare match/input capture)	TPU channel 1	40	H'0450	DTCEB1	_
TGI1B (GR1B compare match/input capture)		41	H'0452	DTCEB0	_
TGI2A (GR2A compare match/input capture)	TPU channel 2	44	H'0458	DTCEC7	
TGI2B (GR2B compare match/ input capture)		45	H'045A	DTCEC6	Low

Interrupt Source	Origin of Interrupt Source	Vector Number	Vector Address	DTCE*	Priority
TGI3A (GR3A compare match/ input capture)	TPU channel 3	48	H'0460	DTCEC5	High A
TGI3B (GR3B compare match/ input capture)		49	H'0462	DTCEC4	
TGI3C (GR3C compare match/ input capture)		50	H'0464	DTCEC3	
TGI3D (GR3D compare match/input capture)		51	H'0466	DTCEC2	
TGI4A (GR4A compare match/ input capture)	TPU channel 4	56	H'0470	DTCEC1	
TGI4B (GR4B compare match/ input capture)		57	H'0472	DTCEC0	
TGI5A (GR5A compare match/ input capture)	TPU channel 5	60	H'0478	DTCED5	
TGI5B (GR5B compare match/ input capture)		61	H'047A	DTCED4	
Reserved for system use	_	64	H'0480	DTCED3	_
		65	H'0482	DTCED2	
		68	H'0488	DTCED1	
		69	H'048A	DTCED0	_
		72	H'0120	DTCEE7	_
		73	H'0124	DTCEE6	_
		74	H'0128	DTCEE5	_
		75	H'012C	DTCEE4	_
RXI0 (reception complete 0)	SCI	81	H'04A2	DTCEE3	_
TXI0 (transmit data empty 0)	channel 0	82	H'04A4	DTCEE2	
RXI1 (reception complete 1)	SCI	85	H'04AA	DTCEE1	
TXI1 (transmit data empty 1)	channel 1	86	H'04AC	DTCEE0	_
RXI2 (reception complete 2)	SCI	89	H'04B2	DTCEF7	
TXI2 (transmit data empty 2)	channel 2	90	H'04B4	DTCEF6	
RM0	HCAN	108	H'04D4	DTCEG5	Low

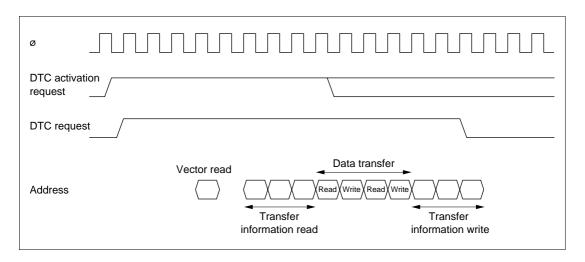
Note: * DTCE bits with no corresponding interrupt are reserved, and should be written with 0.

DTC Operation Timing

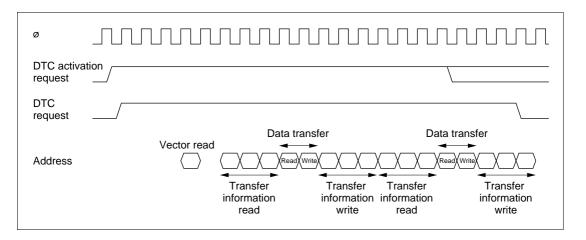
• Example in Normal Mode or Repeat Mode



• Example of Block Transfer Mode (Block Size = 2)



• Example of Chain Transfer



Number of DTC Execution States

Mode	Vector Read I	Register Information Read/Write J	Data Read K	Data Write L	Internal Operations M
Normal	1	6	1	1	3
Repeat	1	6	1	1	3
Block transfer	1	6	N	N	3

N: Block size (initial setting of CRAH and CRAL)

• Number of States Required in Each Execution State

Access To	Access To:					Internal I/O Registers External Devices				
Bus width			32	16	8	16	8	8	16	16
Access sta	tes		1	1	2	2	2	3	2	3
Execution state	Vector read	SI	_	1	_	_	4	6+2m	2	3+m
	Register information read/write	SJ	1	_	_	_	_	_	_	_
	Byte data read	S _K	1	1	2	2	2	3+m	2	3+m
	Word data read	S _K	1	1	4	2	4	6+2m	2	3+m
	Byte data write	S_L	1	1	2	2	2	3+m	2	3+m
	Word data write	S _L	1	1	4	2	4	6+2m	2	3+m
	Internal operation	S_{M}	1	1	1	1	1	1	1	1

Note: The number of execution states is calculated from the formula below.

Number of execution states = I · S_I + Σ (J · S_J + K · S_K + L · S_L) + M · S_M

 Σ indicates the sum of all transfers activated by one activation event (the number in which the CHNE bit is set to 1, plus 1).

Transfer Modes

There are three DTC transfer modes—normal mode, repeat mode, and block transfer mode.

In the DTC, the 24-bit source address register (SAR) designates the transfer source address and the 24-bit destination address register (DAR) designates the transfer destination address. After each transfer, SAR and DAR are independently incremented, decremented, or left unchanged.

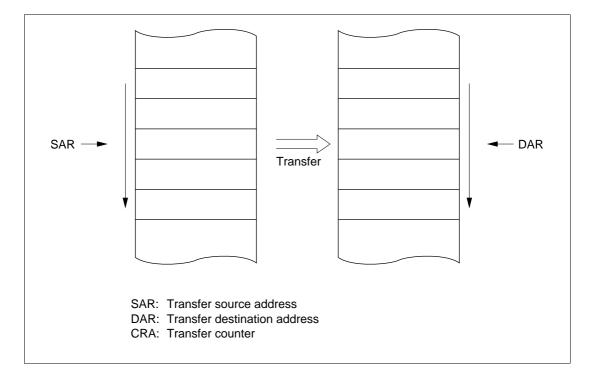
• Overview of DTC Functions

either the source or destination

		Address Registers			
Transfer Mode	Activation Source	Transfer Source	Transfer Destination		
 Normal mode 1-byte or 1-word transfer executed for one transfer request Memory addresses are incremented or decremented by 1 or 2 Up to 65,536 transfers possible 	IRQTPU TGISCI TXI or RXIA/D converter ADISoftware	24 bits	24 bits		
 Repeat mode 1-byte or 1-word transfer executed for one transfer request Memory addresses are incremented or decremented by 1 or 2 After the specified number of transfers (1 to 256), the initial state is restored and operation continues 					
 Block transfer mode One transfer request transfers a block of the specified size Block size is from 1 to 256 bytes or words Up to 65,536 transfers possible A block area can be designated at 					

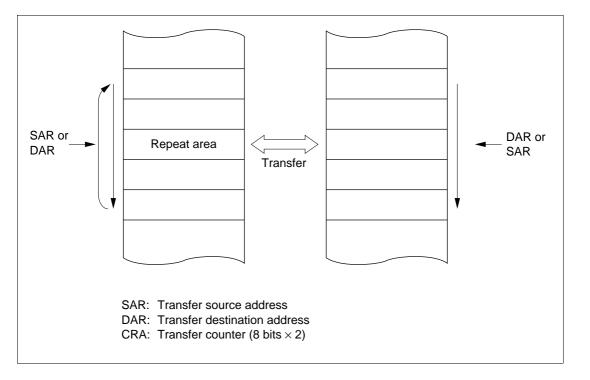
Operation in Normal Mode: In normal mode, one byte or one word of data is transferred in one operation. From 1 to 65,536 transfers can be specified. When the specified number of transfers have ended, a CPU interrupt can be requested.

• Normal Mode Memory Map



Operation in Repeat Mode: In repeat mode, one byte or one word of data is transferred in one operation. From 1 to 256 transfers can be specified. When the specified number of transfers have ended, the initial state specified in the transfer counter and repeat area is restored and transfer is repeated.

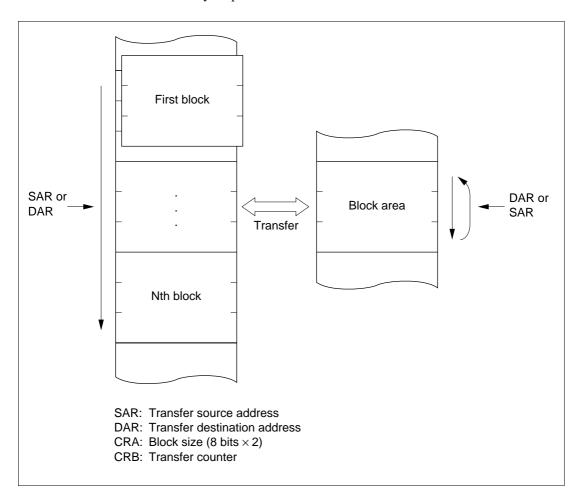
Repeat Mode Memory Map



Operation in Block Transfer Mode: In block transfer mode, one operation transfers one block of data. Either the transfer source or the transfer destination is specified as a block area. The block size is 1 to 256. When the transfer of one block ends, the initial settings of the block size counter and the address register specified in the block area are restored. The other address register is continually incremented, decremented, or left unchanged.

From 1 to 65,536 transfers can be specified. When the specified number of transfers have ended, a CPU interrupt can be requested.

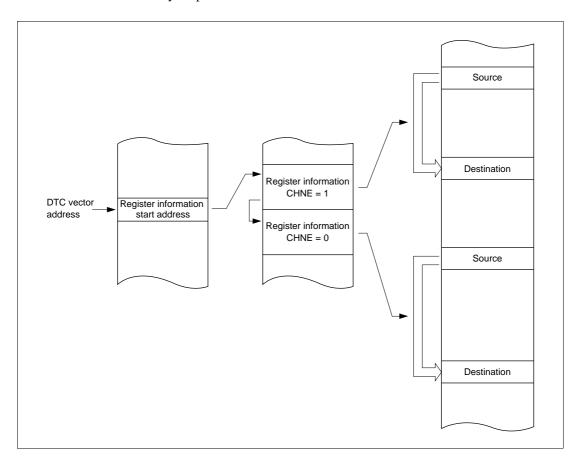
• Block Transfer Mode Memory Map



Chain Transfer

Setting the CHNE bit to 1 enables a number of data transfers to be performed consecutively in response to a single transfer request.

• Chain Transfer Memory Map



3.3 I/O Ports

This series has ten input/output ports (ports 1, A, B, C, D, E, and F), and two input ports (ports 4 and 9).

Each port includes a data direction register (DDR) that controls input/output, a data register (DR) that stores output data, and a port register (PORT) used to read the pin states.

Ports A to E incorporate a MOS input pull-up, and in addition to DDR and DR, have a MOS input pull-up control register (PCR) that turns the MOS input pull-up on or off.

H8S/2623 Series Port Functions in Each Operating Mode

Port	Description	Pins	Mode 4	Mode 5	Mode 6	Mode 7
Port 1	8-bit I/O port Schmitt- triggered input (P16, P14)	P17/PO15/TIOCB2/ TCLKD P16/PO14/TIOCA2/ IRQ1 P15/PO13/TIOCB1/ TCLKC P14/PO12/TIOCA1/ IRQ0 P13/PO11/TIOCD0/ TCLKB/A23 P12/PO10/TIOCC0/ TCLKA/A22 P11/PO9/TIOCB0/ A21 P10/PO8/TIOCA0/ A20	(TCLKA, TCLKI TIOCBO, TIOCO TIOCA2, TIOCE PO8), interrupt address outputs	so functioning as B, TCLKC, TCLKC, TCLKC, TCLKC0, TIOCD0, TI	(D, TIOCAO, DCA1, TIOCB1, pins (PO15 to	8-bit I/O port also functioning as TPU I/O pins (TCLKA, TCLKB, TCLKC, TCLKD, TIOCA0, TIOCB0, TIOCA1, TIOCB1, TIOCA2, TIOCB2), PPG output pins (PO15 to PO8), and interrupt input pins (IRQ0, IRQ1)
Port 4	8-bit input port	P47/AN7 P46/AN6 P45/AN5 P44/AN4 P43/AN3 P42/AN2 P41/AN1 P40/AN0	8-bit input port (AN7 to AN0)	also functioning	as A/D converte	r analog inputs

Port	Description	Pins	Mode 4	Mode 5	Mode 6	Mode 7				
Port 9	• 8-bit input port	P97/AN15 P96/AN14 P95/AN13 P94/AN12 P93/AN11 P92/AN10 P91/AN9 P90/AN8	8-bit input port a (AN15 to AN8)							
	 6-bit I/O port Built-in MOS input pull-up Open-drain output capability 	PA5 PA4 PA3/A19/SCK2 PA2/A18/RxD2 PA1/A17/TxD2 PA0/A16	6-bit I/O port als 2) I/O pins (TxD outputs (A19 to	6-bit I/O port also functioning as SCI (channel 2) I/O pins (TxD2, RxD2, SCK2)						
Port B	8-bit I/O port Built-in MOS input pull-up Open-drain output capability	PB7/A15/TIOCB5 PB6/A14/TIOCA5 PB5/A13/TIOCB4 PB4/A12/TIOCA4 PB3/A11/TIOCD3 PB2/A10/TIOCC3 PB1/A9/TIOCB3 PB0/A8/TIOCA3	8-bit I/O port als (TIOCB5, TIOC TIOCC3, TIOCE outputs (A15 to	8-bit I/O port also functioning as TPU I/O pins (TIOCB5, TIOCA5, TIOCB4, TIOCA4, TIOCD3, TIOCC3, TIOCB3, TIOCC3)						
Port C	8-bit I/O port Built-in MOS input pull-up Open-drain output capability	PC7/A7 PC6/A6 PC5/A5/SCK1/IRQ5 PC4/A4/RxD1 PC3/A3/TxD1 PC2/A2/SCK0/IRQ4 PC1/A1/RxD0 PC0/A0/TxD0	Address outputs	s (A7 to A0)	8-bit I/O port also function- ing as SCI (channel 0 and 1) I/O pins (TxD0, RxD0, SCK0, TxD1, RxD1, SCK1), interrupt input pins (IRQ4, IRQ5), and address outputs (A7 to A0)	8-bit I/O port also functioning as SCI (channel 0 and 1) I/O pins (TxD0, RxD0, SCK0, TxD1, RxD1, SCK1) and interrupt input pins (IRQ4, IRQ5)				
Port D	8-bit I/O portBuilt-in MOS input pull-up	PD7/D15 to PD0/D8	Data bus input/o	output		I/O port				
Port E	8-bit I/O portBuilt-in MOS input pull-up	PE7/D7 to PE0/D0	In 8-bit-bus mod		put/output	I/O port				

Port	Description	Pins	Mode 4	Mode 5	Mode 6	Mode 7			
Port F	8-bit I/O port	PF7/ø	When DDR = 0 When DDR = 1	: input port (after reset): Ø	output	When DDR = 0 (after reset): input port			
				When DDR = 1: Ø output					
		PF6/AS	PF6/AS RD, HWR, LWR, AS outputs						
		PF5/RD	ADTRG, IRQ3	ADTRG, IRQ3					
		PF4/HWR		input					
		PF3/LWR/ADTRG/ IRQ3							
		PF2/WAIT/BREQO	When WAITE = 0 and BREQOE = 0 (after reset): I/O port When WAITE = 1 and BREQOE = 0: WAIT input						
			When WAITE = input						
		PF1/BACK	When BRLE = 0 (after reset): I/O port						
		$PF0/\overline{BREQ}/\overline{IRQ2}$ When $BRLE = 1$: \overline{BREQ} input, \overline{BACK} output							
			IRQ2 output						

3.4 16-Bit Timer Pulse Unit (TPU)

The 16-bit timer pulse unit (TPU) comprises six 16-bit timer channels.

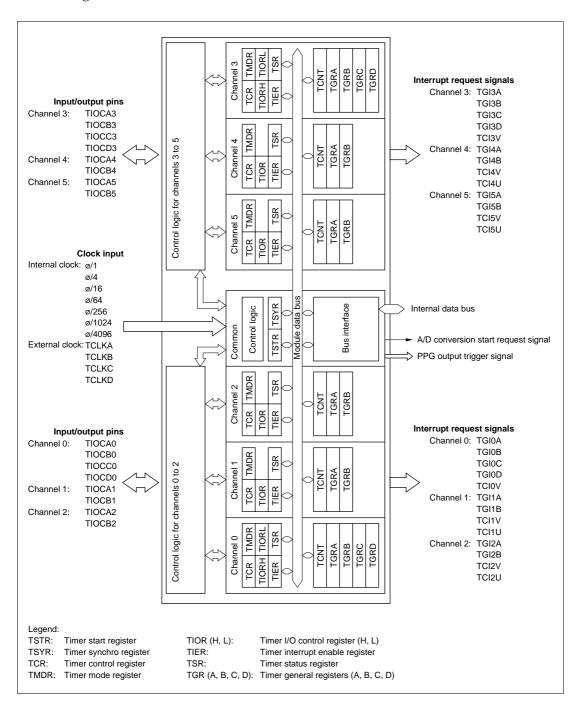
The TPU can perform PWM output, pulse width measurement, and two-phase encoder processing, and can activate the data transfer controller (DTC). It can also generate a programmable pulse generator (PPG) output trigger and A/D converter start trigger.

Features

- Maximum 16-pulse input/output
 - A total of 16 timer general registers (TGRs) are provided (four each for channels 0 and 3, and two each for channels 1, 2, 4, and 5), each of which can be set independently as an output compare/input capture register
- Selection of 8 counter input clocks for each channel
 - Internal clocks: ø, ø/4, ø/16, ø/64, ø/256, ø/1024, ø/4096
 - External clocks: TCLKA, TCLKB, TCLKC, TCLKD
- The following operations can be set for each channel:
 - Waveform output at compare match: Selection of 0, 1, or toggle output
 - Input capture function: Selection of rising edge, falling edge, or both edge detection
 - Counter clear operation: Counter clearing possible by compare match or input capture
 - Synchronous operation:
 - Multiple timer counters (TCNT) can be written to simultaneously
 - Simultaneous clearing by compare match and input capture possible
 - Register simultaneous input/output possible by counter synchronous operation
 - PWM mode:
 - Any PWM output duty can be set
 - Maximum of 15-phase PWM output possible by combination with synchronous operation
- Buffer operation settable for channels 0 and 3
 - Input capture register double-buffering possible
 - Automatic rewriting of output compare register possible
- Phase counting mode settable independently for each of channels 1, 2, 4, and 5
 - Two-phase encoder pulse up/down-count possible
- Cascaded connection operation
 - Channel 2 (channel 5) input clock operates as 32-bit counter by setting channel 1 (channel 4) overflow/underflow
- Fast access via internal 16-bit bus
 - Fast access is possible via a 16-bit bus interface

- 26 interrupt sources
 - For channels 0 and 3, four compare match/input capture dual-function interrupts and one overflow interrupt can be requested independently
 - For channels 1, 2, 4, and 5, two compare match/input capture dual-function interrupts, one overflow interrupt, and one underflow interrupt can be requested independently
- Automatic transfer of register data
 - Block transfer, 1-word data transfer, and 1-byte data transfer possible by data transfer controller (DTC) activation
- Programmable pulse generator (PPG) output trigger can be generated
 - Channel 0 to 3 compare match/input capture signals can be used as PPG output trigger
- A/D converter conversion start trigger can be generated
 - Channel 0 to 5 compare match A/input capture A signals can be used as A/D converter conversion start trigger
- Module stop mode can be set
 - As the initial setting, TPU operation is halted. Register access is enabled by exiting module stop mode.

Block Diagram of TPU



Interrupt Sources and Data Transfer Controller (DTC) Activation

• List of TPU Interrupts

Channel	Interrupt Source	Description	DTC Activation	Priority		
0	TGI0A	TGR0A input capture/compare match	Possible	High		
	TGI0B	TGR0B input capture/compare match	Possible	_		
	TGI0C	TGR0C input capture/compare match	Possible	_		
	TGI0D	TGR0D input capture/compare match	Possible	_		
	TCI0V	TCNT0 overflow	Not possible	_		
1	TGI1A	TGR1A input capture/compare match	Possible	_		
	TGI1B	TGR1B input capture/compare match	Possible	_		
	TCI1V	TCNT1 overflow	Not possible	_		
	TCI1U	TCNT1 underflow	Not possible	_		
2	TGI2A	TGR2A input capture/compare match	Possible	_		
	TGI2B	TGR2B input capture/compare match	Possible	_		
	TCI2V	TCNT2 overflow	Not possible	_		
	TCI2U	TCNT2 underflow	Not possible	_		
3	TGI3A	TGR3A input capture/compare match	Possible	_		
	TGI3B	TGR3B input capture/compare match	Possible	_		
	TGI3C	TGR3C input capture/compare match	Possible	_		
	TGI3D	TGR3D input capture/compare match	Possible	_		
	TCI3V	TCNT3 overflow	Not possible	_		
4	TGI4A	TGR4A input capture/compare match	Possible	_		
	TGI4B	TGR4B input capture/compare match	Possible	_		
	TCI4V	TCNT4 overflow	Not possible	_		
	TCI4U	TCNT4 underflow	Not possible	_		
5	TGI5A	TGR5A input capture/compare match	Possible	_		
	TGI5B	TGR5B input capture/compare match	Possible	_		
	TCI5V	TCNT5 overflow	Not possible	_		
	TCI5U	TCNT5 underflow	Not possible	Low		

Note: This table shows the initial state immediately after a reset. The relative channel priorities can be changed by the interrupt controller.

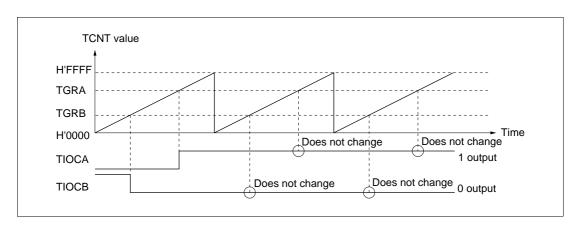
Operation

Each channel has a TCNT and TGR register. TCNT performs up-counting, and is also capable of free-running operation, cyclic counting, and external event counting. Each TGR can be used as an input capture register or output compare register.

- Buffer Operation
 - When TGR is an output compare register
 When a compare match occurs, the value in the buffer register for the relevant channel is transferred to TGR.
 - When TGR is an input capture register
 When input capture occurs, the value in TCNT is transferred to TGR and the value previously held in TGR is transferred to the buffer register.

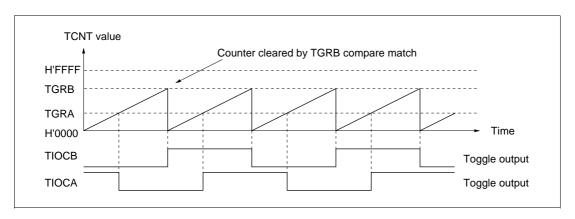
Waveform Output by Compare Match

- 0, 1, or toggle output can be selected.
- Example of 0 Output/1 Output Operation
 In this example, TCNT has been designated as a free-running counter, and settings have been made so that 0 is output by compare match A, and 1 is output by compare match B.



Example of Toggle Output

In this example, settings have been made so that TCNT counter clearing is performed by compare match B, and output is toggled by both by compare match A and compare match B.

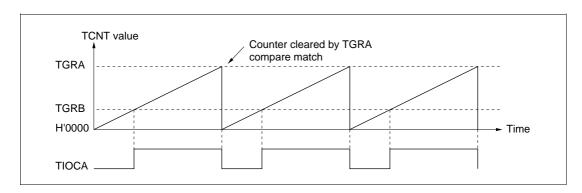


PWM Modes

In PWM mode, PWM waveforms are output from the output pins. There are two PWM modes—PWM mode 1 with a maximum of 8-phase pulse output, and PWM mode 2 with a maximum of 15-phase pulse output.

PWM Mode 1: PWM output is generated by pairing TGRA with TGRB and TGRC with TGRD. In PWM mode 1, a maximum 8-phase PWM output is possible.

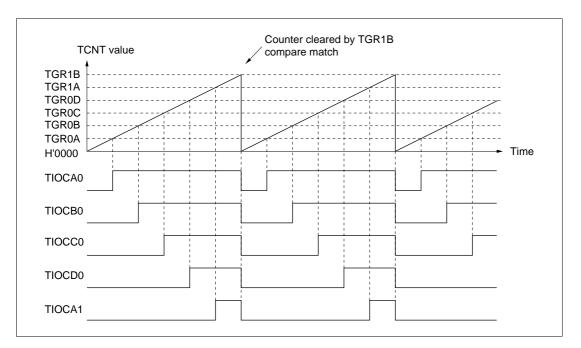
• Example of Operation in PWM Mode 1
In this example, TGRA compare match is set as the TCNT clearing source, 0 is set for the TGRA initial output value and output value, and 1 is set as the TGRB output value. In this case, the value set in TGRA is the cycle, and the value set in TGRB is the duty.



PWM Mode 2: PWM output is generated using one TGR register as the cycle register and the others as duty registers. In PWM mode 2, a maximum 15-phase PWM output is possible by combined use with synchronous operation.

• Example of Operation in PWM Mode 2

In this example, synchronous operation is designated for channels 0 and 1, TGR1B compare match is set as the TCNT clearing source, and 0 is set for the initial output value and 1 for the output value of the other TGR registers, to output a 5-phase PWM waveform. In this case, the value set in TGR1B is the cycle, and the value set in the other TGR registers is the duty.

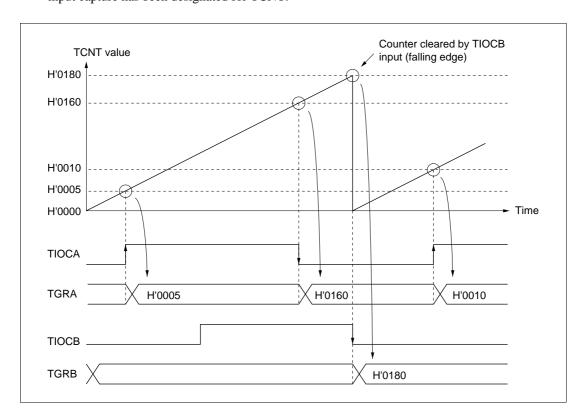


Input Capture Operation

The TCNT value can be transferred to TGR on detection of the TIOC pin input edge.

Rising edge, falling edge, or both edges can be selected as the input edge.

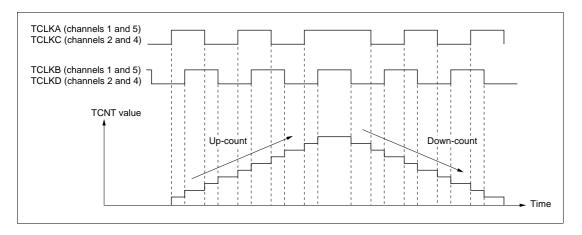
Example of Input Capture Operation
 In this example, both rising and falling edges have been selected as the TIOCA pin input edge, falling edge has been selected as the TIOCB pin input edge, and counter clearing by TGRB input capture has been designated for TCNT.



Phase Counting Mode

In phase counting mode, the phase difference between two external clock inputs is detected and TCNT operates as an up/down-counter. There are four modes (phase counting modes 1 to 4) with different setting conditions. These modes can be set for channels 1, 2, 4, and 5.

• Example of Operation in Phase Counting Mode 1



• Up/Down-Count Conditions in Phase Counting Mode

TCLKA (Channels 1 and 5) TCLKC (Channels 2 and 4)	TCLKB (Channels 1 and 5) TCLKD (Channels 2 and 4)	Operation
High level		Up-count
Low level	7_	
	Low level	
7_	High level	
High level	7_	Down-count
Low level		
	High level	
7_	Low level	

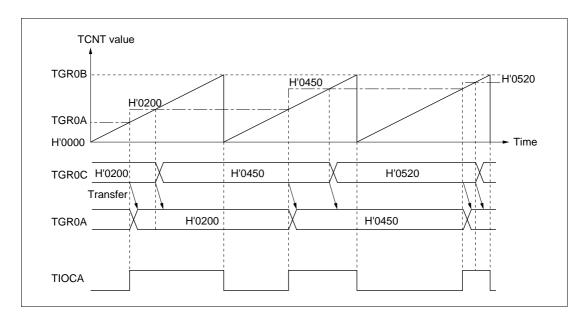
Legend

∴ Falling edge → : Falling edge

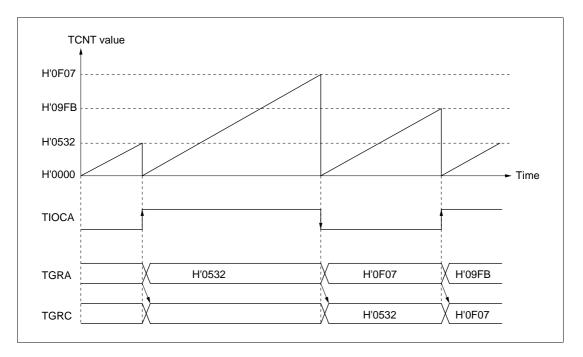
Buffer Operation

Buffer operation, provided for channels 0 and 3, enables TGRC and TGRD to be used as buffer registers.

• Example of Buffer Operation (1) (When TGR is an Output Compare Register)
In this example, PWM mode 1 has been designated for channel 0, and buffer operation has been designated for TGRA and TGRC. The settings used are TCNT clearing by a compare match B, 1 output at compare match A, and 0 output at compare match B. When a compare match A occurs, the output is changed and the value in buffer register TGRC is simultaneously transferred to timer general register TGRA.



Example of Buffer Operation (2) (When TGR is an Input Capture Register)
In this example, TGRA has been designated as an input capture register, and buffer operation has been designated for TGRA and TGRC. Counter clearing by TGRA input capture has been set for TCNT, and detection of both rising and falling edges has been selected for the TIOCA pin. When the TCNT value is stored in TGRA upon occurrence of input capture A, the value previously stored in TGRA is simultaneously transferred to TGRC.



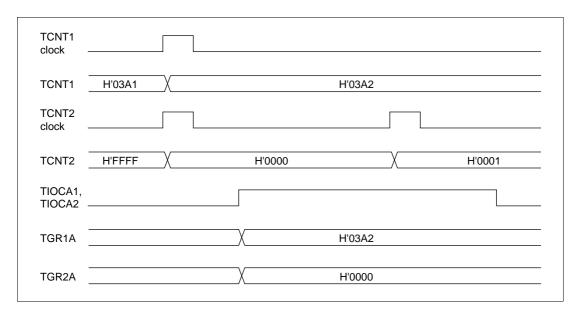
Cascading

In cascaded operation, two 16-bit counters for different channels are used together as a 32-bit counter. Channels 1 and 2, and channels 4 and 5, can be cascaded.

• Example of Cascaded Operation

In this example, counting upon TCNT2 overflow/underflow has been set for TCNT1, TGR1A and TGR2A have been designated as input capture registers, and TIOC pin rising edge detection has been selected. When a rising edge is input to the TIOCA1 and TIOCA2 pins simultaneously, the upper 16 bits of the 32-bit data are transferred to TGR1A, and the lower 16 bits to TGR2A.

Example of Cascaded Operation (32-Bit Input Capture Operation)



Synchronous Operation

When synchronous operation is designated for a channel, TCNT for that channel performs synchronous presetting and clearing. That is, when TCNT for a channel designated for synchronous operation is rewritten, the TCNT counters for the other channels are also rewritten at the same time. When any clearing condition occurs, the TCNT counters for the other channels are also cleared simultaneously.

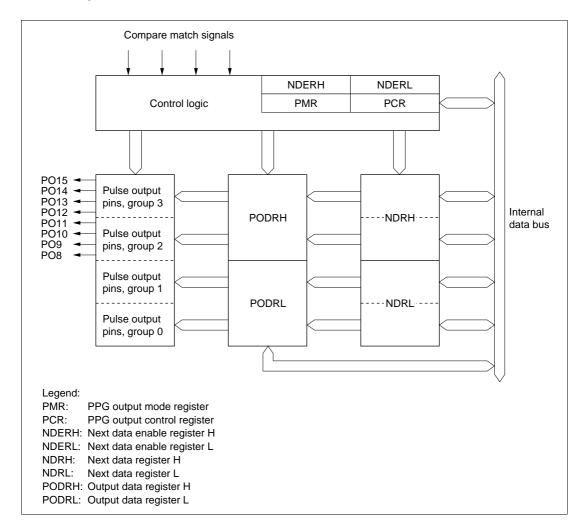
3.5 Programmable Pulse Generator (PPG)

The programmable pulse generator (PPG) can handle up to 8 outputs simultaneously, using a signal from the 16-bit timer-pulse unit (TPU) as its input.

Features

- 8-bit output data
 - Maximum 8-bit data can be output, and pulse output can be enabled on a bit-by-bit basis
- Two output groups
 - Output trigger signals can be selected in 4-bit groups to provide a maximum of two 4-bit outputs
- Selectable output trigger signals
 - Output trigger signals can be selected for each group from the compare match signals of four TPU channels
- Non-overlap operation
 - A non-overlap interval can be set between pulse outputs
- Can operate together with the data transfer controller (DTC)
 - The compare match signals selected as output trigger signals can activate the DTC for sequential output of data without CPU intervention
- Inverted output can be specified
 - Inverted data can be output for each group

Block Diagram of PPG



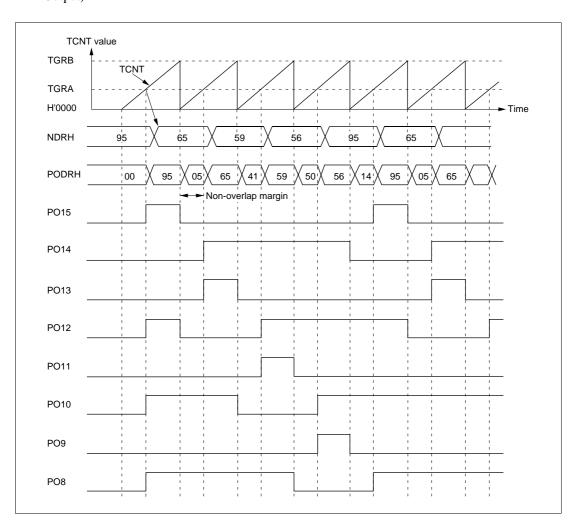
Example of Four-Phase Complementary Non-Overlapping Output

In this example, pulse output is used for four-phase complementary non-overlapping pulse output.

When a TGRB compare match occurs, outputs change from 1 to 0. When a TGRA compare match occurs, outputs change from 0 to 1. Set the non-overlap margin in the TPU TGRA for which the output trigger is selected, and set the cycle in TGRB.

If the DTC is set for activation by a TGIA interrupt, pulse output can be performed without imposing a load on the CPU.

 Non-Overlapping Pulse Output Example (Four-Phase Complementary Non-Overlapping Output)



3.6 Watchdog Timer (WDT)

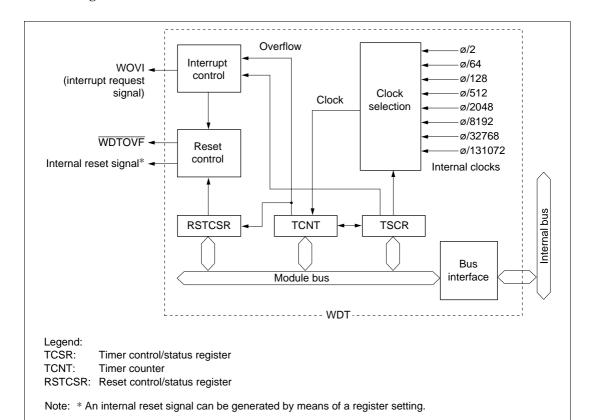
This series has a one-channel on-chip watchdog timer (WDT) for monitoring system operation. When this watchdog function is not needed, the WDT can be used as an interval timer.

When the subclock is selected as the input clock, the WDT can be used as a realtime clock timer.

Features

- Switchable between watchdog timer mode and interval timer mode
- WDTOVF output when in watchdog timer mode
 - If the counter overflows, the WDT outputs the WDTOVF signal externally. It is possible to select whether the chip is internally reset or an NMI interrupt is generated at the same time.
- Interrupt generation when in interval timer mode
 - If the counter overflows, the WDT generates an interval timer interrupt.
- Choice of eight counter input clocks

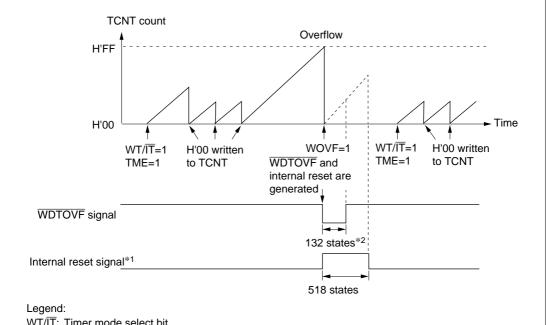
Block Diagram of WDT



Watchdog Timer Operation

The examples below show this timer used as a watchdog timer. The timer counter (TCNT) starts counting up using the specified clock.

• WDT Watchdog Timer Operation



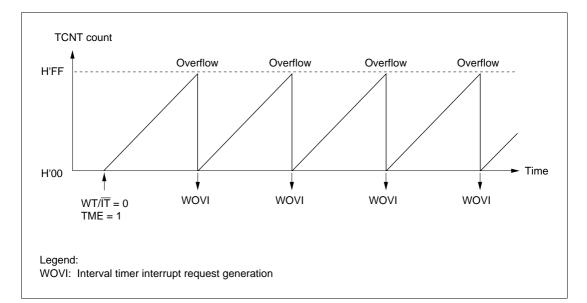
WT/IT: Timer mode select bit TME: Timer enable bit

Notes: 1. The internal reset signal is generated only if the RSTE bit is set to 1.

2. 130 states when the RSTE bit is cleared to 0.

Interval Timer Operation

An example of the use of the WDT as an interval timer is shown here. The timer counter (TCNT) starts counting up on the specified clock, and an interval timer interrupt (WOVI) occurs each time TCNT overflows. This function can be used to generate interrupt requests at regular intervals.



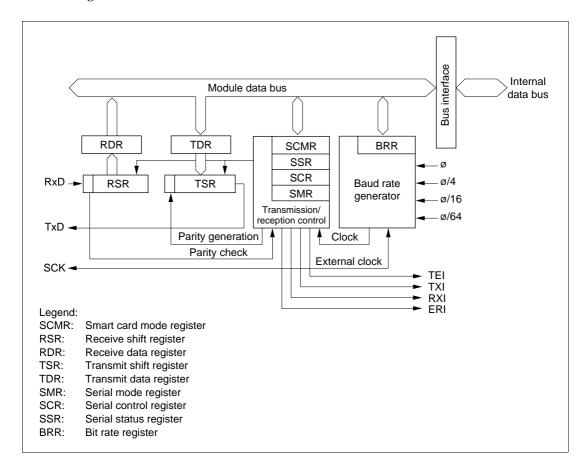
3.7 Serial Communication Interface (SCI)

This series has a serial communication interface (SCI) with three independent channels. All three channels have the same functions, and can handle both asynchronous and synchronous serial communication. A function is also provided for serial communication between processors (multiprocessor communication function).

Features

- Choice of asynchronous or synchronous serial communication mode
- Full-duplex communication capability
- Data register double-buffering enables continuous transmission/reception
- Internal dedicated baud rate generator allows any bit rate to be selected
- Selection of internal baud rate generator or external clock input (SCK pin) as serial clock source
- Detection of three receive errors
 - Overrun errors, framing errors, and parity errors can be detected
- Break detection
- Four interrupt sources
 - Four interrupt sources—transmit data empty, transmission end, receive data full, and receive error—that can issue requests independently
 - The transmit data empty interrupt and receive data full interrupt can activate the data transfer controller (DTC) to execute data transfer
- Built-in multiprocessor communication function
- Selection of LSB-first or MSB-first transfer
 - This choice can be made regardless of the communication mode (with the exception of 7-bit data transfer in asynchronous mode)
- Module stop mode can be set
 - As the initial setting, SCI operation is halted. Register access is enabled by exiting module stop mode.

Block Diagram of SCI



SCI Interrupt Sources

Channel	Interrupt Source	Description	DTC Activation	Priority*
0	ERI	Interrupt due to receive error (ORER, FER, or PER)	Not possible	High Å
	RXI	Interrupt due to receive data full (RDRF)	Possible	_
	TXI	Interrupt due to transmit data empty (TDRE)	Possible	_
	TEI	Interrupt due to transmission end (TEND)	Not possible	_
1	ERI	Interrupt due to receive error (ORER, FER, or PER)	Not possible	
	RXI	Interrupt due to receive data full (RDRF)	Possible	_
	TXI	Interrupt due to transmit data empty (TDRE)	Possible	_
	TEI	Interrupt due to transmission end (TEND)	Not possible	_
2	ERI	Interrupt due to receive error (ORER, FER, or PER)	Not possible	_
	RXI	Interrupt due to receive data full (RDRF)	Possible	_
	TXI	Interrupt due to transmit data empty (TDRE)	Possible	_
	TEI	Interrupt due to transmission end (TEND)	Not possible	Low

Note: *This table shows the initial state immediately after a reset. Relative priorities among channels can be changed by means of the interrupt controller.

SCI Asynchronous Communication

Asynchronous mode is a serial communication mode in which synchronization is achieved on a character by character basis, using a start bit and one or two stop bits.

- Twelve serial data transfer formats
 - Data length: 7 or 8 bitsStop bit length: 1 or 2 bits
 - Parity: Even/odd/none
 - Multiprocessor bit: 1 or 0
- Selection of internal baud rate generator or external clock from SCK pin as clock source
- Transmit/receive clock can be output from SCK pin
- Break detection
 - A break can be detected by reading the RxD pin level directly in case of a framing error
- Multiprocessor communication capability

Serial Transfer Formats and Frame Lengths in Asynchronous Mode

	SMR	Setting	gs		Serial Transmit/Receive Format and Frame Length										
CHR	PE	MP	STOP	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	S				8-bit	data				STOP	-	
0	0	0	1	S		8-bit data							STOP STOP		
0	1	0	0	s		8-bit data							Р	STOP	-
0	1	0	1	S		8-bit data							Р	STOP	STOP
1	0	0	0	s			7	-bit da	ata			STOF	- >		
1	0	0	1	s			7	-bit da	ata			STOP	STOP	-	
1	1	0	0	s			7	-bit da	ata			Р	P STOP		
1	1	0	1	s			7	-bit da	ata			Р	STOP	STOP	-
0	_	1	0	S		8-bit data							МРВ	STOP	-
0	_	1	1	S		8-bit data						МРВ	STOP	STOP	
1	_	1	0	S		7-bit data MPB						MPB	STOP	-	
1	_	1	1	S		7-bit data MPB						STOP	STOP	=	

Legend:

S: Start bit STOP: Stop bit P: Parity bit

MPB: Multiprocessor bit

Multiprocessor Communication Function

A multiprocessor format, in which a multiprocessor bit is added to the transfer data, can be used for serial communication, enabling data transfer to be performed among a number of processors.

The transmitting station first sends the ID of the receiving station with which it wants to perform serial communication as data with a "1" MPB (multiprocessor bit) added. It then sends transmit data as data with a "0" MPB added.

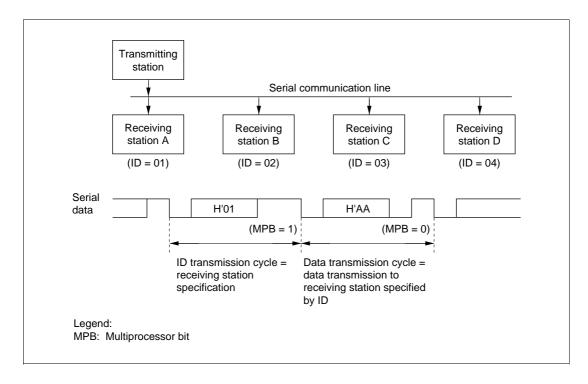
Receiving stations skip data until data with a "1" MPB is received. Each receiving station then compares that data with its own ID. The station whose ID matches then continues with reception, and accepts data. Stations whose ID does not match continue to skip the data until data with a "1" MPB is sent again.

SCI Synchronous Communication

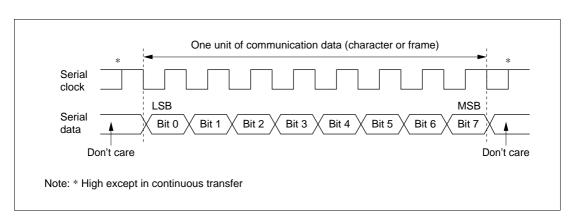
In synchronous mode, data is transmitted or received in synchronization with clock pulses, making it suitable for continuous, high-speed serial communication.

- Data length: 8 bits per character
- Overrun error detection
- Selection of internal baud rate generator or external clock from SCK pin as transmit/receive clock source
- Selection of LSB-first or MSB-first transfer
- Communication is possible with chips provided with a synchronous mode, such as the H8 Series, HD64180, and HD6301

Example of Inter-Processor Communication Using Multiprocessor Format (Transmission of Data H'AA to Receiving Station A)



Data Format in Synchronous Communication



BRR Settings for Various Bit Rates (Synchronous Mode)

Operating Frequency ø (MHz)

Bit Rate (bits/s)	ø = 2		ø = 4			ø = 8		ø = 10		ø = 16		ø = 20	
	n	N	n	N	n	N	n	N	n	N	n	N	
110	3	70	_	_									
250	2	124	2	249	3	124	_	_	3	249			
500	1	249	2	124	2	249	_	_	3	124	_	_	
1 k	1	124	1	249	2	124	_	_	2	249	_	_	
2.5 k	0	199	1	99	1	199	1	249	2	99	2	124	
5 k	0	99	0	199	1	99	1	124	1	199	1	249	
10 k	0	49	0	99	0	199	0	249	1	99	1	124	
25 k	0	19	0	39	0	79	0	99	0	159	0	199	
50 k	0	9	0	19	0	39	0	49	0	79	0	99	
100 k	0	4	0	9	0	19	0	24	0	39	0	49	
250 k	0	1	0	3	0	7	0	9	0	15	0	19	
500 k	0	0*	0	1	0	3	0	4	0	7	0	9	
1 M			0	0*	0	1			0	3	0	4	
2.5 M							0	0*			0	1	
5 M											0	0*	

Note: As far as possible, the setting should be made so that the error is no more than 1%.

Legend:

Blank: Cannot be set.

—: Can be set, but there will be a degree of error.

*: Continuous transfer is not possible.

The BRR setting is found from the following formulas.

Asynchronous mode:

$$N = \frac{\phi}{64 \times 2^{2n-1} \times B} \times 10^6 - 1$$

Synchronous mode:

$$N = \frac{\phi}{8 \times 2^{2n-1} \times B} \times 10^6 - 1$$

Where N: BRR setting for band rate generator $(0 \le N \le 255)$

ø: Operating frequency (MHz)

B: Bit rate (bits/s)

n: Baud rate generator input clock (n = 0 to 3)
(See the table below for the relation between n and the clock.)

		SMR Settings				
n	Clock	CKS1	CKS0			
0	Ø	0	0			
1	ø/4	0	1			
2	ø/16	1	0			
3	ø/64	1	1			

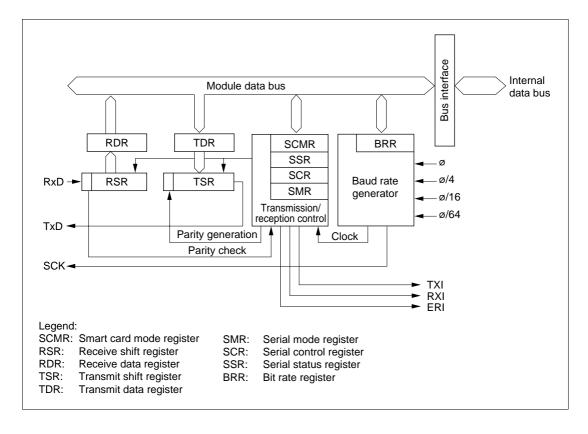
3.8 Smart Card Interface

The SCI supports a smart card interface as an IC card interface serial communication function conforming to ISO/IEC7816-3 (Identification Card).

Features

- · Asynchronous mode
 - Data length: 8 bits
 - Parity bit generation and checking
 - Transmission of error signal (parity error) in receive mode
 - Error signal detection and automatic data retransmission in transmit mode
 - Direct convention and inverse convention both supported
- Internal baud rate generator allows any bit rate to be selected
- Three interrupt sources
 - Three interrupt sources—transmit data empty, receive data full, and transmit/receive error—that can issue requests independently
 - The transmit data empty interrupt and receive data full interrupt can activate the data transfer controller (DTC) to execute data transfer

Block Diagram of Smart Card Interface

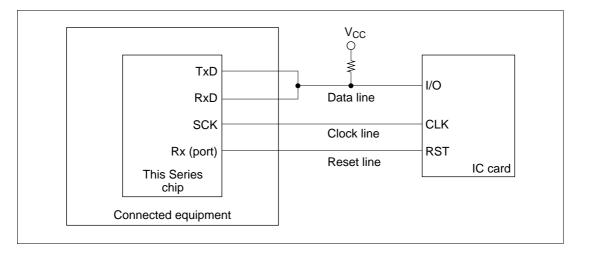


Operation

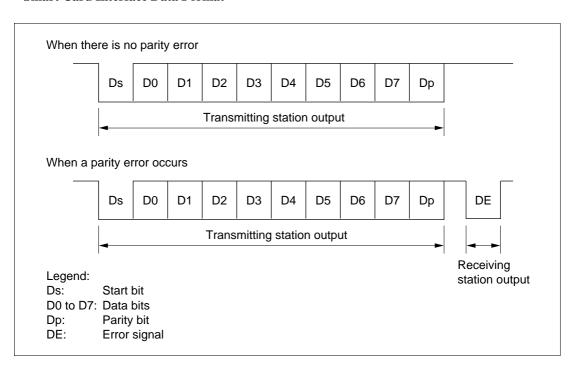
Overview

- Only asynchronous communication is supported, with one frame consisting of 8-bit data plus a parity bit.
- In transmission, a guard time of at least 2 etu (Elementary Time Unit: the time for transfer of one bit) is left between the end of the parity bit and the start of the next frame.
- If a parity error is detected during reception, a low error signal level is output for a 1 etu period 10.5 etu after the start bit (except in block transfer mode).
- If the error signal is sampled during transmission, the same data is transmitted automatically after the elapse of 2 etu or longer (except in block transfer mode).

Schematic Diagram of Smart Card Interface Pin Connections



Smart Card Interface Data Format



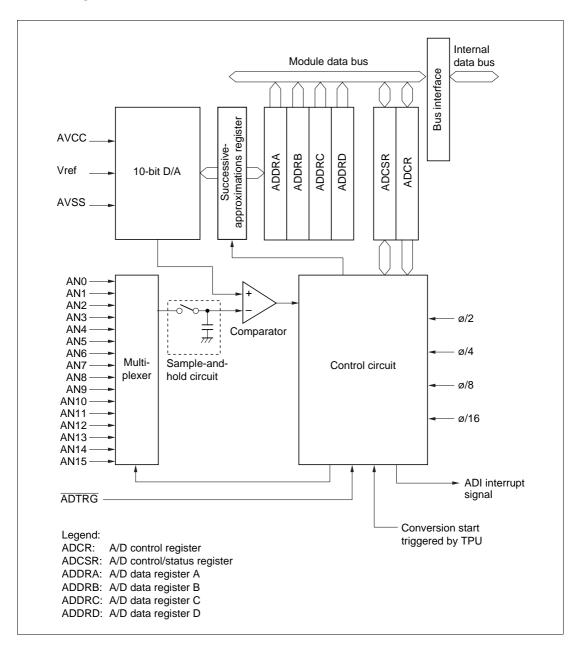
3.9 A/D Converter

This series incorporates an on-chip A/D converter with 10-bit precision. Analog signals can be input on up to 16 channels.

Features

- 10-bit resolution
- 16 input channels
- Settable analog conversion voltage range
 - Conversion of analog input from 0 V to V_{ref}, with the reference voltage pin (Vref) as the analog reference voltage
- High-speed conversion
 - Minimum conversion time: 13.3 µs per channel (at 20 MHz operation)
- Selection of single mode or scan mode
 - Single mode: A/D conversion on one channel
 - Scan mode: continuous A/D conversion on one to four channels
- Three kinds of conversion start
 - Selection of software or timer conversion start trigger (TPU), or ADTRG pin
- Four data registers
 - Conversion results held in a 16-bit data register for each channel
- Sample-and-hold function
- A/D conversion end interrupt generation
 - A/D conversion end interrupt (ADI) request can be generated at the end of A/D conversion
- Module stop mode can be set
 - As the initial setting, A/D converter operation is halted. Register access is enabled by exiting module stop mode.

Block Diagram of A/D Converter



Input Channel Setting

16-channel analog input is performed by means of the scan mode bit (SCAN) and channel select bits (CH3 to CH0) in ADCSR.

		CH1	CH0	Description				
СНЗ	CH2			Single Mode (SCAN = 0)	Scan Mode (SCAN = 1)			
0	0	0	0	AN0 (initial value)	AN0			
			1	AN1	AN0, AN1			
		1	0	AN2	AN0 to AN2			
			1	AN3	AN0 to AN3			
	1	0	0	AN4	AN4			
			1	AN5	AN4, AN5			
		1	0	AN6	AN4 to AN6			
			1	AN7	AN4 to AN7			
1	0	0	0	AN8	AN8			
			1	AN9	AN8, AN9			
		1	0	AN10	AN8 to AN10			
			1	AN11	AN8 to AN11			
	1	0	0	AN12	AN12			
			1	AN13	AN12, AN13			
		1	0	AN14	AN12 to AN14			
			1	AN15	AN12 to AN15			

Operation

The successive approximation method is used for A/D conversion, with a 10-bit resolution. There are two operating modes—single or scan.

• Single Mode

conversion start condition.

Single mode is selected when A/D conversion is to be performed on a single channel only. A/D conversion is started when the ADST bit is set to 1, according to the specified A/D

On completion of conversion, the ADF flag is set to 1. If the ADIE bit is set to 1 at this time, an ADI interrupt request is generated.

• Scan Mode

Scan mode is selected when A/D conversion is to be performed repeatedly on a number of channels.

Once the ADST bit is set to 1 according to the specified A/D conversion start condition, A/D conversion is performed repeatedly on the selected channels until the ADST bit is cleared to 0 by software.

An ADI interrupt request can be generated on completion of the first conversion operation for all the selected channels.

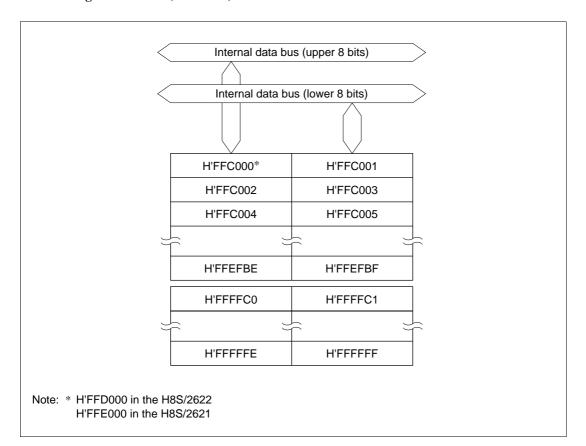
3.10 RAM

The H8S/2623 has 12 kbytes of on-chip high-speed static RAM, the H8S/2622 has 8 kbytes, and the H8S/2621 has 4 kbytes.

The on-chip RAM is connected to the bus master by a 16-bit data bus, enabling both byte data and word data to be accessed in one state. This makes it possible to perform fast word data transfer.

The on-chip RAM can be enabled or disabled by means of the RAM enable bit (RAME) in the system control register (SYSCR).

Block Diagram of RAM (H8S/2623)



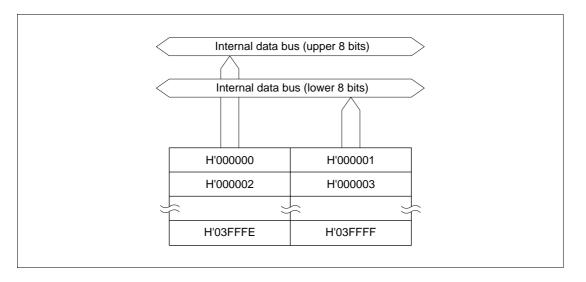
3.11 ROM

The H8S/2623 has 256 kbytes of on-chip flash memory or mask ROM. The H8S/2622 has 128 kbytes of on-chip mask ROM, and the H8S/2621 has 64 kbytes.

The ROM is connected to the bus master by a 16-bit data bus, enabling both byte data and word data to be accessed in one state. This allows rapid instruction fetches and improves processing speed.

In addition to erasing and programming with a dedicated PROM programmer, the flash memory version can also be erased and programmed on-board.

ROM Block Diagram (H8S/2623)

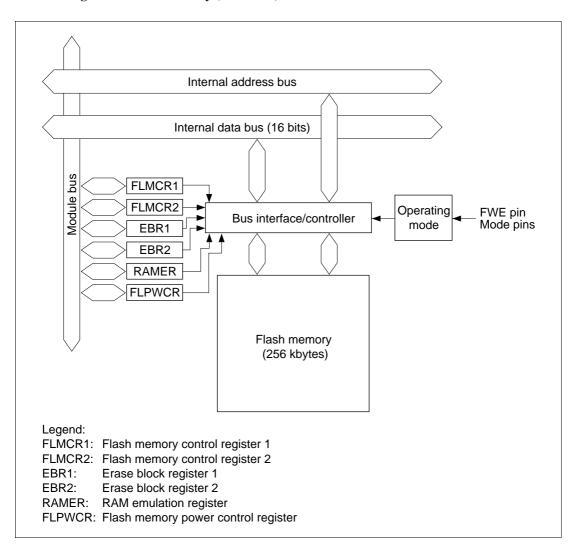


On-chip ROM is enabled or disabled by means of mode pins MD1 and MD0.

Features of Flash Memory

- Four operating modes
 - Program mode
 - Erase mode
 - Program-verify mode
 - Erase-verify mode
- Programming/erase methods
 - 128 bytes programmed simultaneously
 - 4-kbyte, 32-kbyte, or 64-kbyte units can be set for block erasing
- Programming/erase times
 - Programming time: TBD (typ.) per 128-byte programming operation TBD (typ.) per byte
 - Erase time: TBD (typ.) per block
- Reprogramming up to 100 times
- On-board programming modes
 - Boot mode
 - User program mode
- · Automatic bit rate adjustment
 - With data transfer in boot mode, the bit rate of the chip can be automatically adjusted to match the transfer bit rate of the host
- Flash memory emulation by RAM
 - Part of the RAM area can be overlapped onto flash memory, to emulate flash memory programming in real time
- Protect modes
 - There are two protect modes, software and hardware, which allow protected status to be designated for flash memory program/erase/verify operations
- Writer mode
 - On-board programming mode and writer mode, using a PROM programmer, are available as flash memory program/erase modes

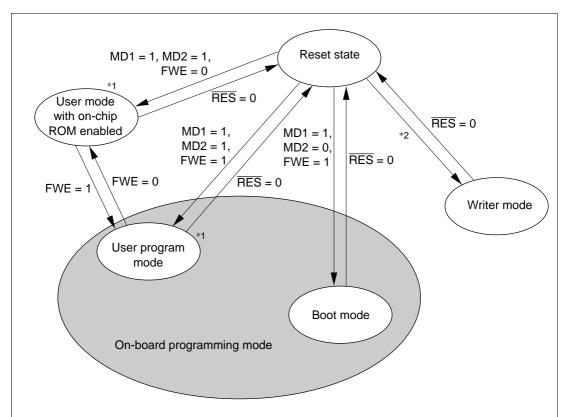
Block Diagram of Flash Memory (H8S/2623)



Mode Transitions

When the mode pins and FWE pins are set in the reset state and a reset-start is executed, the MCU enters one of the operating modes as shown in the figure below. In user mode, flash memory can be read but not programmed or erased.

Flash memory can be programmed and erased in boot mode, user program mode, and writer mode.



Notes: Only make a transition between user mode and user program mode when the CPU is not accessing the flash memory.

- 1. RAM emulation possible
- 2. MD0 = 0, MD1 = 0, MD2 = 0, P14 = 0, P16 = 0, PF0 = 1

Section 4 Power-Down Modes

In addition to the normal program execution state, this series has power-down modes in which operation of the CPU and oscillator is halted and power consumption is reduced. Low-power operation can be achieved by individually controlling the CPU, on-chip supporting modules, and so on.

This series' operating modes are as follows:

- 1. High-speed mode
- 2. Medium-speed mode
- 3. Sleep mode
- 4. Module stop mode
- 5. Software standby mode
- 6. Hardware standby mode

Of these, 2 to 6 are power-down modes. Sleep mode is a CPU mode, medium-speed mode is a CPU and bus master mode, and module stop mode is an on-chip supporting module mode (including bus masters other than the CPU). A combination of certain of these modes can be set.

Medium-Speed Mode: When bits SCK2, SCK1, and SCK0 in the system clock control register (SCKCR) are set to 1 in high-speed mode, medium-speed mode is entered as soon as the current bus cycle ends. In medium-speed mode, the CPU operates on the operating clock (\emptyset /2, \emptyset /4, \emptyset /8, \emptyset /16, or \emptyset /32) specified by bits SCK2 to SCK0. However, on-chip supporting functions other than the bus masters operate on the high-speed clock (\emptyset).

Sleep Mode: If a SLEEP instruction is executed when the SSBY bit in SBYCR is cleared to 0, the CPU enters sleep mode. In sleep mode, CPU operation stops but the contents of the CPU's internal registers are retained. Other supporting functions do not stop.

Module Stop Mode: Module stop mode can be used to start and stop individual on-chip supporting modules.

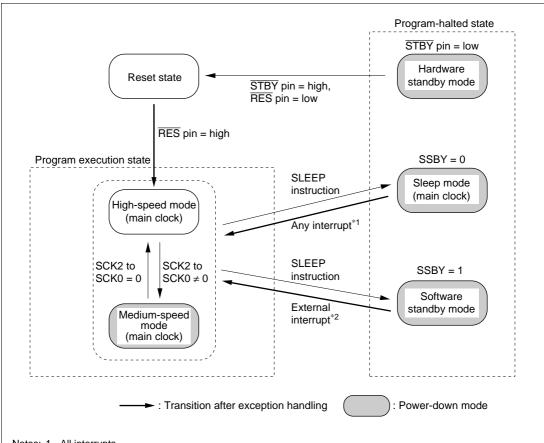
When the MSTP bit corresponding to a particular supporting function in the module stop control register (MSTPCR) is set to 1, operation of the specified module stops at the end of the bus cycle and a transition is made to module stop mode. The CPU continues operating independently.

Software Standby Mode: If a SLEEP instruction is executed when the SSBY bit in SBYCR is set to 1, software standby mode is entered. In this mode, the CPU, on-chip supporting functions, and oscillator all stop. However, the states of on-chip supporting functions other than the SCI and A/D, and the states of I/O ports, are retained.

Hardware Standby Mode: When the STBY pin is driven low, a transition is made to hardware standby mode from any state.

In hardware standby mode, all functions enter the reset state and stop operation, resulting in extremely low power consumption. As long as the prescribed voltage is supplied, on-chip RAM data is retained. I/O ports are set to the high-impedance state.

Mode Transition Diagram



- Notes: 1. All interrupts
 - 2. NMI, IRQ0 to IRQ5
 - When a transition is made between modes by means of an interrupt, transition cannot be made on interrupt source generation alone. Ensure that interrupt handling is performed after accepting the interrupt request.
 - From any state except hardware standby mode, a transition to the reset state occurs whenever RES goes
 - From any state, a transition to hardware standby mode occurs whenever STBY goes low.

Chip Internal States in Each Mode

Function System clock oscillator		High- Speed	Medium- Speed	Sleep	Module Stop	Software Standby	Hardware Standby		
		Functioning	Functioning	Functioning	Functioning	Halted	Halted		
CPU	Instruction registers	Functioning	Medium- speed operation	Halted (retained)	High/ medium- speed operation	Halted (retained)	Halted (undefined)		
External	NMI	Functioning	Functioning	Functioning	Functioning	Functioning	Halted		
interrupts	IRQ0 to IRQ5								
Supporting functions	WDT	Functioning	Functioning	Functioning	Halted (retained)	Halted (retained)	Halted (reset)		
	DTC	Functioning	Medium- speed operation	Functioning	Halted (retained)	Halted (retained)	Halted (reset)		
	TPU	Functioning	Functioning	Functioning	Halted (retained)	Halted (retained)	Halted (reset)		
	РСВ	_							
	PPG	_							
	SCI	Functioning	Functioning	Functioning	Halted (reset) Halted (reset) I		Halted (reset)		
	A/D	_							
	RAM	Functioning	Functioning	Functioning (DTC)	Functioning	Retained	Retained		
	I/O	Functioning	Functioning	Functioning	Functioning	Retained	High impedance		

Notes: "Halted (retained)" means that internal register values are retained. The internal state is "operation suspended."

Power-Down Mode Transition Conditions

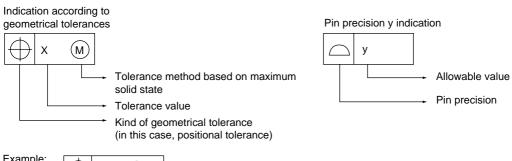
State before		Bit States whe	n _ State after Transition	State after Recovery		
Transition	SSBY	LSON	by SLEEP Instruction	by Interrupt		
High-speed/	0	0	Sleep	High-speed/medium-speed		
medium-speed	1	0	Software standby	High-speed/medium-speed		

[&]quot;Halted (reset)" means that internal register values and internal states are initialized. In module stop mode, only modules for which a stop setting has been made are halted (reset or retained).

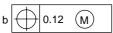
Appendix

Packages

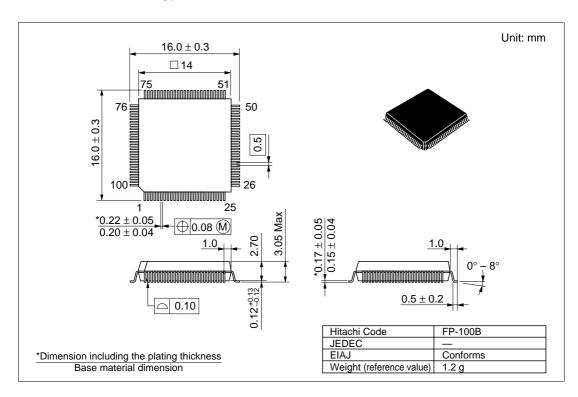
Package Dimension Diagrams (Unit: mm)



Example:



This indicates that the allowable pin displacement from the true central position is 0.12 mm when pin width b is the maximum dimension. If b is smaller than the maximum dimension, the tolerance can be extended accordingly.



H8S/2623 Series Overview

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