

# **16 BY 16 BIT COMPLEX MULTIPLIER**

(Supersedes version October 1995 verison, DS3707 - 3.0)

The PDSP16116A will multiply two complex (16 + 16) bit words every 50ns and can be configured to output the complete complex (32 + 32) bit result within a single cycle. The data format is fractional two's complement.

The PDSP16116/A contains four 16 x 16 Array Multipliers, two 32 bit Adder/Subtractors and all the control logic required to support Block Floating Point Arithmetic as used in FFT applications. In combination with a PDSP16318, the PDSP16116A forms a two chip 10MHz Complex Multiplier Accumulator with 20 bit accumulator registers and output shifters. The PDSP16116 in combination with two PDSP16318s and two PDSP1601s forms a complete 10MHz Radix 2 DIT FFT Butterfly solution which fully supports Block Floating Point Arithmetic. The PDSP16116/A has an extremely high throughput that is suited to recursive algorithms as all calculations are performed with a single pipeline delay (two cycle fall-through).

### **FEATURES**

- Complex Number (16 + 16) X (16 + 16) Multiplication
- Full 32 bit Result
- 20MHz Clock Rate
- Block Floating Point FFT Butterfly Support
- -1 times -1 Trap
- Two's Complement Fractional Arithmetic
- TTL Compatible I/O
- Complex Conjugation
- 2 Cycle Fall Through
- 144 pin PGA or QFP packages

### **APPLICATION**

- Fast Fourier Transforms
- Digital Filtering
- Radar and Sonar Processing
- Instrumentation
- Image Processing

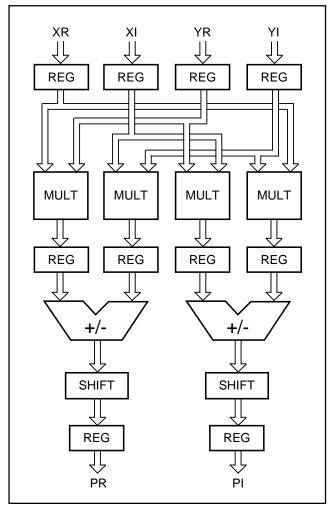


Fig.1 Simplified Block Diagram

#### ASSOCIATED PRODUCTS

PDSP16510

PDSP16318/A Complex Accumulator
PDSP16112/A (16 + 16) X (12 + 12) Complex Multiplier
PDSP16330/A Pythagoras Processor
PDSP16350 PDSP16256 Programmable FIR Filter

Single Chip FFT Processor

The PDSP16116 has a number of features tailored for System applications.

# -1 x -1 Trap

In multiply operations utilising Twos Complement Fractional notation, the -1  $\times$  -1 operation forms an invalid result as +1 is not representable in the fractional number range. The PDSP16116/A eliminates this problem by trapping the -1  $\times$  -1 operation and forcing the Multiplier result to become the most positive representable number.

# **Complex Conjugation**

Many algorithms utilising complex arithmetic require conjugation of complex data stream. This operation has

traditionally required an adiditional ALU to multiply the imaginary component by -1. The PDSP16116 eliminates the requirement for the extra ALU by offering on chip complex conjugation of either of the two incoming complex data words with no loss in throughput.

#### **Easy Interfacing**

As with all PDSP family members the PDSP16116 has registered I/O for data and control. Data inputs have independent clock enables and data outputs have independent three state output enables.

Signal	Туре	Description	Normal mode Configuration
XR15:0 XI15:0 YR15:0 YR15:0 PR15:0 PR15:0 PI15:0 CLK CEX CEY CONX CONY ROUND MBFP SOBFP EOPSS AR15:13 AI15:13 WTA1:0 WTB1:0 WTOUT1:0 SFTA1:0 SFTR2:0 GWR4:0 OSEL1:0 OER, OEI VDD GND	INPUT INPUT INPUT OUTPUT OUTPUT INPUT OUTPUT OUTPUT OUTPUT OUTPUT INPUT INPUT INPUT INPUT INPUT INPUT INPUT INPUT OUTPUT OUTPUT OUTPUT OUTPUT OUTPUT OUTPUT INPUT INPUT INPUT POWER POWER	16 bit input for real x data 16 bit input for imag x data 16 bit input for reaal y data 16 bit input for imag y data 16 bit output for imag y data 16 bit output for img p data 16 bit output for img p data Clock, new data is loaded on rising edge of CLK Clock, enable X-port input register Clock, enable Y-port input register Conjugate X data Conjugate Y data Rounds the real & imag results Mode select (BFP/Normal) Start of BFP operations ** End of pass ** 3 MSB's from real part of A-word ** Word tag from A-word Word tag from B-word / shift control * Word tag output ** Shift control for A-word / overflow flag * Shift control for accumulator resul ** Global weighting register contents ** Selects the desired output configuration Output enables +5V Supply All supply pins 0V Supply must be connected	Tie Low Tie Low Tie Low Tie Low Tie Low

<sup>\*</sup> Indicates pin performs different functions in BFP / Normal modes.

<sup>\*\*</sup> Indicates pin is used only in BFP mode

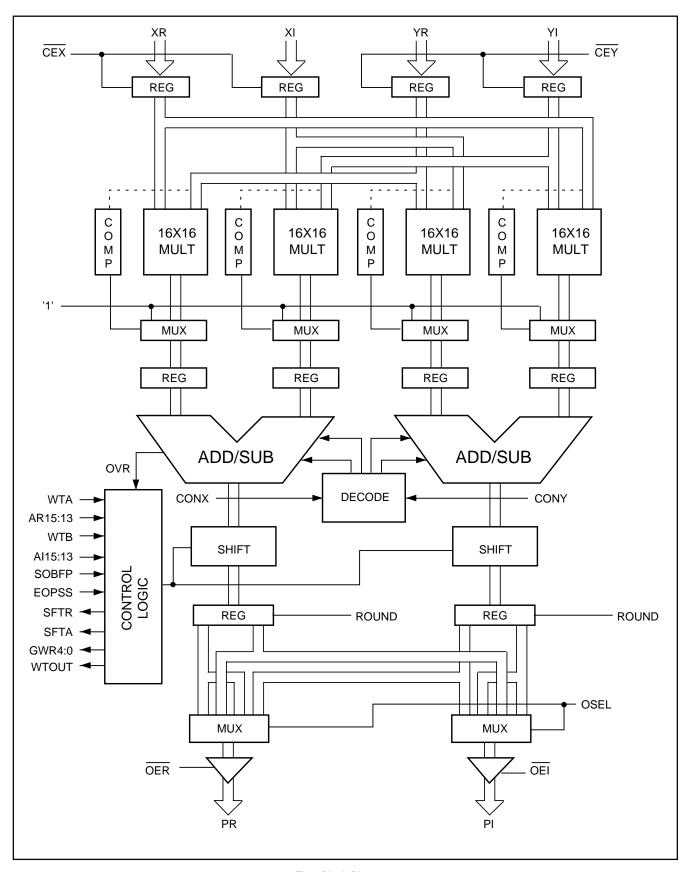


Fig.2 Block Diagram

R	VDD	XI3	XI6	XI8	XI11	XI14	XI15	CEX	XR13	XR12	XR8	XR6	XR5	XR1	YR13
Р	(YI1)	GND	XI2	XI4	XI7	XI10	XI12	CEY	XR11	XR10	XR7	XR4	XR2	YR14	YR12
N	YI5	YI2	(XIO)	XI1	XI5	(XI9)	XI13	XR15	XR14	XR9	XR3	XR0	YR15	VDD	YR9
M	YI6	YI4	YIO										GND	YR11	YR7
L	YI8	YI7	YI3										YR10	YR8	YR4
κ	YI12	YI10	YI9										YR6	YR5	YR1
J	YI13	YI11	YI14										YR2	YR3	YR0
Н	AR15	AR14	YI15										SOBFP	EOPSS	VDD
G	AR13	Al13	Al14										WTB1	WTA0	WTB0
F	Al15	CONX	CONY										CLK	MBFP	WTA1
E	ROUND	SFTR2	(SFTR0)										GWRO	OSELO	OSEL1)
D	(OEI)	WTOUTO	PI14	Loca	ition pin								GWR3	SFTA1	OER
С	SFTR1)	PI15	PI13	PI12	PI9	VDD	PIO	PRO	VDD	PR7	PR11	VDD	GWR4	GWR1)	SFTA0
В	WTOUT1)	GND	PI10	PI8	GND	Pl4	PI3	PR2	GND	PR6	PR9	PR12	GND	PR15	GWR2
Α	VDD	PI11	PI7	PI6	PI5	Pl2	PI1	PR1	PR3	PR4	PR5	PR8	PR10	PR13	PR14
•	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15

Fig.3 Pin Allocation Diagram (Bottom View) 144 pin PGA - AC144

GG	SIG	GG	SIG	GG	SIG	GG	SIG
1	PI14	37	XI1	73	GND	109	VDD
2	PI15	38	XI2	74	VDD	110	GND
3	WTOUT1	39	XI3	75	YR12	111	PR13
4	WTOUT0	40	XI4	76	YR11	112	PR12
5	SFTR0	41	XI5	77	YR10	113	PR11
6	SFTR1	42	XI6	78	YR9	114	PR10
7	SFTR2	43	XI7	79	YR8	115	PR9
8	OEI	44	XI8	80	YR7	116	PR8
9	CONX	45	XI9	81	YR6	117	PR7
10	CONY	46	XI10	82	YR5	118	PR6
11	ROUND	47	XI11	83	YR4	119	PR5
12	Al13	48	XI12	84	YR3	120	GND
13	Al14	49	XI13	85	YR2	121	VDD
14	Al15	50	XI14	86	YR1	122	PR4
15	AR13	51	XI15	87	YR0	123	PR3
16	AR14	52	CEY	88	<u>EOPSS</u>	124	PR2
17	AR15	53	CEX	89	VDD	125	PR1
18	YI15	54	XR15	90	SOBFP	126	PR0
19	YI14	55	XR14	91	WTB1	127	PI0
20	YI13	56	XR13	92	WTB0	128	PI1
21	YI12	57	XR12	93	WTA1	129	PI2
22	YI11	58	XR11	94	WTA0	130	PI3
23	YI10	59	XR10	95	MBFP	131	PI4
24	YI9	60	XR9	96	CLK	132	VDD
25	YI8	61	XR8	97	OSEL1	133	PI5
26	YI7	62	XR7	98	OSEL0	134	GND
27	YI6	63	XR6	99	<u>OER</u>	135	PI6
28	YI5	64	XR5	100	SFTA0	136	PI7
29	YI4	65	XR4	101	SFTA1	137	PI8
30	YI3	66	XR3	102	GWR0	138	PI9
31	YI2	67	XR2	103	GWR1	139	PI10
32	YI1	68	XR1	104	GWR2	140	PI11
33	YI0	69	XR0	105	GWR3	141	PI12
34	XI0	70	YR15	106	GWR4	142	PI13
35	GND	71	YR14	107	PR15	143	GND
36	VDD	72	YR13	108	PR14	144	VDD

All GND and VDD pins must be used.

Fig.3A Pin Allocation Diagram - 144 pin ceramic QFP - GG144

#### NORMAL MODE OPERATION

When the MBFP mode select input is held low the 'Normal' mode of operation is selected. This mode supports all Complex Multiply operations that do not require Block Floating Point arithmetic.

#### **Multiplier Stage**

Complex two's complement fractional data is loaded into the X and Y input registers via the X and Y Ports on the rising edge of CLK. The Real and Imaginary components of the fractional data are each assumed to have the following format

BIT NUMBER	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
WEIGHTING	S	2 <sup>-1</sup>	2 <sup>-2</sup>	2-3	2-4	2 <sup>-5</sup>	2 <sup>-6</sup>	2.7	2 <sup>-8</sup>	2-9	2 <sup>-10</sup>	2 <sup>-11</sup>	2 <sup>-12</sup>	2 <sup>-13</sup>	2 <sup>-14</sup>	2 <sup>-15</sup>

Where  $S = \text{sign bit which has an effective weighting } -2^{0}$ 

The value of the 16 bit two's complement word is

Value =  $(-1xS)+(bit14x2^{-1})+(bit13x2^{-2})+(bit12x2^{-3})$ . . .

The X & Y port registers are individually enabled by the CEX & CEY signals respectively. If the registers are required to be permanently enabled, then these signals may be tied to ground. On each clock cycle the contents of the input registers are passed to the four multipliers to start a new Complex Multiply operation. Each Complex Multiply operation requires four partial products (Xr x Yr), (Xr x Yi), (Xi x Yr), (Xi x Yi), all of which are calculated in parallel by the four 16 x 16 Multipliers. Only one clock cycle is required to complete the multiply stage before the Multiplier results are loaded into the Multiplier output registers for passing on to the Adder/Subtractors in the next cycle. Each multiplier produces a 31 bit result with the duplicate sign bit eliminated. The format of the output data from the Multipliers is

BIT NUMBER	30	29	28	27	26	25	24	 7	6	5	4	3	2	1	0
WEIGHTING	S	2 <sup>-1</sup>	2 <sup>-2</sup>	2-3	2 <sup>-4</sup>	2 <sup>-5</sup>	2 <sup>-6</sup>	 2 <sup>-23</sup>	2 <sup>-24</sup>	2 <sup>-25</sup>	2 <sup>-26</sup>	2 <sup>-27</sup>	2 <sup>-28</sup>	2 <sup>-29</sup>	2 <sup>-30</sup>

The effective weighting of the sign bit is -20

#### **Result Correction**

Due to the nature of the fraction twos complement representation it is possible to represent -1 exactly but not 1. With conventional multipliers this causes a problem when -1 is multiplied by -1 as the multiplier produces an incorrect result. The PDSP16116 includes a trap to ensure that the most positive number (value =  $1.2^{-30}$ ), (hex = 7FFFFFFF) is subsituted for the incorrect result. The multiplier result is therefore always a (correct) fractional value.

# **Complex Conjugation**

Either the X or Y input data may be complex conjugated by asserting the CONX or CONY signals respectively. Asserting either of these signals has the effect of inverting (multiplying by -1) the imaginary component of the respective input. Table 3 shows the effect of CONX and CONY on the X and Y inputs.

FUNCTION	OPERATION	CONX	CONY
X x Y X x Conj Y Conj X x Y Invalid	(XR+XI)x(YR+YI)	low	low
	(XR+XI)x(YR-YI)	high	low
	(XR-XI)x(YR+YI)	low	high
	Invalid	high	high

Table 3 Conjugate Functions

# Adder / Subtractor Stage

The 31 bit Real and Imaginary results from the Multipliers are passed to two 32 bit Adder/Subtractors. The Adder calculates the imaginary result ( $(Xr \times Yi) + (Xi \times Yr)$ ) and the Subtractor calculates the Real result ( $(Xr \times Yr) = (Xi \times Yi)$ ). Each Adder/Subtractor produces a 32 bit result with the following format.

BIT NUMBER	31	30	29	28	27	26	 8	7	6	5	4	3	2	1	0
WEIGHTING	s	2 <sup>0</sup>	2 <sup>-1</sup>	2-2	2.3	2 <sup>-4</sup>	 2 <sup>-22</sup>	2 <sup>-23</sup>	2 <sup>-24</sup>	2 <sup>-25</sup>	2 <sup>-26</sup>	2 <sup>-27</sup>	2 <sup>-28</sup>	2 <sup>-29</sup>	2 <sup>-30</sup>

The effective weighting of the sign bit is -21

# Rounding

The ROUND control when asserted rounds the most significant 16 bits of the full 32 bit result from the shifter. If the ROUND signal is active (High), then bit 16 is set to a one, rounding the most significant 16 bits of the shifted result. (The least significant 16 bits are unaffected). Inserting a one ensures that the rounding error is never greater than 1LSB, and that no DC bias is introduced as a result of the rounding processes.

The format of the Rounded result is;

BIT NUMBER	31	30	29	28	27	 18	17	16	15	14	13		2	1	0
WEIGHTING	S	20	2 <sup>-1</sup>	2-2	2.3	 2 <sup>-12</sup>	2 <sup>-13</sup>	2 <sup>-14</sup>	2 <sup>-15</sup>	2 <sup>16</sup>	2 <sup>-17</sup>		2 <sup>-28</sup>	2 <sup>-29</sup>	2 <sup>-30</sup>
_				D 1/41				_				<u> </u>			

The effective weighting of the sign is -21

#### **Shifter**

Each of the two Adder/Subtractors are followed by Shifters controlled via the WTB control input. These shifters can each apply two different shifts, however the same shift is applied to both real and imaginary components. The two shift options are:

i) WTB1:0 = 11 Shift complex product one place to the left.

The effective weighting of the sign bit is -20

 ii) WTB1:0 = 00, 01 or 10 Shift complex product three place to the right. The effective weighting of the sign bit is -2<sup>4</sup>.

#### Overflow

If the left shift option is selected and the Adder/Subtractor contain a 32 bit word, then an invalid result will be passed to the output. An invalid output arising from this combination of events will be flagged by the SFTA0 flag output. The SFTA0 Flag will go high if either the real or imaginary reslut is invalid.

#### **Output Select**

The output from the Shifters is passed to the Output Select Mux, which is controlled via the OSEL inputs. These inputs are not registered and hence allow the output combination to be changed within each cycle. The full complex 64 bit result from the multiplier may therefore be output within a single cycle. The OSEL control selects four different output combinations as as summarised in Table 4.

OSEL1	OSEL0	PR	PI
0	0	MSR	MSI
0	1	LSR	LSI
1	0	MSR	LSR
1	1	MSI	LSI

Table 3 Output Selection

(Where MSR and LSR are the most and least siginificant 16 bit words of the Real Shifter output, MSI and LSI are the most and least significant 16 bit words of the imaginary Shifter output).

The output select options allow two different modes for extracting the full 32 bit result from the PDSP16116. The first mode treats the two 16 bit outputs as real and imaginary ports allowing the real and imaginary results to be output in two halves on the real and imaginary output ports. The second mode treats the two 16 bit outputs as one 32 bit output and allows the real and imaginary results to be output as 32 bit words.

#### PIN DESCRIPTIONS

#### XR, XI, YR, YI

Data inputs 16 bits: Data is loaded into the input registers from these ports on the rising edge of CLK. The data format is Twos Complement Fractional, where the MSB (sign bit) is bit 15. In normal mode the weighting of the MSB is -20 ie -1.

#### PR, PI

Data outputs 16 bits: Data is clocked into the output registers and passed to the PR and PI outputs on the rising edge of CLK. The data format is Twos Complement Fractional. The field of the internal result selected for output via PR and PI is controlled by signals OSEL1:0 (see Table 4).

#### **CLK**

Common Clock to all internal register.

#### CEX, CEY

Clock enables for X and Y input ports: When low these inputs enable the CLK signal to the X or Y input registers allowing new data to be clocked into the Multiplier.

#### CONX, CONY

If either of these inputs are high on the rising edge of CLK, then the data in the associated input has its imaginary component inverted (multiplied by -1), see Table 3. CONX and CONY affect data input on the same clock rising edge.

#### ROUND

The ROUND control is used to round the most significant 16 bits of the output register. The rounding operation takes place one cycle after the ROUND input is taken high. The ROUND input is not latched and is intended to be tied high or low depending upon the application.

#### **MBFP**

Mode select: When high, Block Floating Point (BFP) mode is selected. This allows the device to maintain the dynamic range of the data using a series of word tags. This is especially useful in FFT appllications. When low, the chip operates in normal mode for more general applications. This pin is intended to be tied high or low, depending on application.

#### **SOBFP** (BFP MODE ONLY)

Start of BFP: This input should be held low for the first cycle of the first pass of the BFP calculations (see Fig.7). It serves to reset the internal registers associated with BFP control. When operating in normal mode this input should be tied low.

#### **EOPSS (BFP MODE ONLY)**

End of pass: This input should be held low for the last cycle of each pass and for the lay time between passes. It instructs the control logic to update the value of the global weighting register and prepare the BFP circuitry for the next pass. When operating in normal mode this input should be tied low.

#### AR15:13 (BFP MODE ONLY)

Three Msbs of the real part of the A-word: These are used in the FFT butterfly application to deteremine the magnitude of the real part of the A-word and, hence, to determine if there will be any chage of word growth in the PDSP16318 Complex Accumulator. When operating in normal mode, these inputs are not used and may be tied low.

#### AI15:13 (BFP MODE ONLY)

Three Msbs of the imaginary part of the A-word : used in the same fashion as AR.

#### SFTR2:0 (BFP MODE ONLY)

Accumulator result shift control. These pins should be linked directly to the S2:0 pins on the PDSP16318 Complex Accumulator. They control the accumulator's barrel shifter (see Table 5). The purpose of this shift is to minimise sign extension in the multiplier or accumulator ALU's. When operating in normal mode, these output are superfluous.

SFTR2:0	FUNCTION
000	Reserved
001	Reserved
010	Reserved
011	Shift right by one
100	No shift
101	Shift left by one
110	Shift left by two
111	Reserved

Table 5 Auccumulator Shifts (BFP mode)

#### **GWR4:0** (BFP MODE ONLY)

Contents of the global weighting register: This stores the weighting of the largest word present with respect to the weighting of the original input words. Hence, if the contents of the GWR are 00010, this indicates that the largest word currently being processed has its binary point two bits to the right of the original data at the start of the BFP calculations. The contents of this register are updated at the end of each pass, according to the largest value of WTOUT occuring during that pass. (i.e. If WTOUT = 11, then GWR will be increased by 2). The GWR is presented in two's complement format. These outputs are superfluous in normal mode.

#### WTOUT1:0 (BFP MODE ONLY)

Word tag output. This tag records the weighting of the output words from the current cycle relative to the current global weighting register (see Table 6). It should be stored along with the A' and B' words as it will form the input word tags, WTA and WTB, for each complex word during the next pass. These outputs are superfluous in normal mode.

WTOUT1:0	Weighting of the output relative to the current global weighting register
0 0	One less
0 1	The same
1 0	One more
11	Two more

Table 6 Word Tag Weightings

#### WTA1:0 (BFP MODE ONLY)

Word tag from the A-word. This word records the weighting of the A-word relative to the global weighting register on the previous pass. Although the A-word inself is not processed in the PDSP16116, this information is required by the control logic for the radix-2 butterfly FFT application. These inputs should be tied low in normal mode.

#### WTB1:0 (BFP & NORMAL MODES)

In BFP mode, this is the word tag from the B-word. This is operated in the same manner as WTA but for the B-word. The value of the word tags are used to ensure that the binary weighting of the A word and the product of the complex multiplier are the same at the inputs to the complex accumulator. Depending on which word is the larger, the weighting adjustment is performed using either the internal shifter or an external shifter controlled by SFTA. The word tags are also used to maintain the weighting of the final result to within plus two and minus one binary points relative to the new GWR. (On the first pass all word tags will be ignored).

In normal mode, these inputs perform a different function. They directly control the internal shifter at the output port as shown in Table 7.

WTB1:0	FUNCTION
11 00 01 10	shift complex product one place to the left shift complex product three places to the right

Table 7 Normal Mode Shift Control

#### SFTA1:0 (BFP & NORMAL MODES)

In BFP mode, these signals act as as the A-word shift control. They allow shifting from one to four places to the right, see Table 8. Depending on the relative weightings of the A-words and the complex product, the A-word may have to be shifted to the right to ensure compatible weightings at the inputs to the PDSP16318 complex accumulator. (The two words must have the same weighting if they are to be added).

In normal mode, SFTA0 performs a different a different function. If WTB1:0 is set to implement a left shift, then overflow will occur if the data is fully 32 bits wide. This pin is used to flag such an overflow. SFTA1 is not used in normal mode.

WTB1:0	FUNCTION
0 0	Shift A-word 1 places to the right
0 1	Shift A-word 2 places to the right
1 0	Shift A-word 3 places to the right
1 1	Shift A-word 4 places to the right

Table 8 External A-word shift control

#### OSEL1:0

The outputs from the device are selected by the OSEL0 & OSEL1 instruction bits. These controls allow selection of the output combination during the current cycle. (They are not registered). These are four possible output configurations that allow either complex outputs of the most or least significant bytes, or real or imaginary outputs of the full 32 bit word (see Table 4). OSEL0 and OSEL1 should both be tied low when in BFP mode.

#### **BFP MODE FFT APPLICATION**

The PDSP16116 may be used as the main arithmetic unit of the butterfly processor which will allow the following FFT benchmarks:

1024 point complex radix-2 transform in 517us 512 point complex radix-2 transform in 235us

256 point complex radix-2 transform in 106us

In addition, with pin MBFP tied high, the BFP circuitry within the PDSP16116 can be used to adaptively rescale data throughout the course of the FFT so as to give high-resolution results.

The BFP system on the PDSP16116 can be used with any variation of the Radix-2 Decimation-In-Time FFT - e.g. the

Constant Geometry algorithm, the In-Place algorithm etc. An N-point Radix-2 DIT FFT is split into log (N) passes. Each pass consists of N/2 'butterflies', each performing the operation:

$$A' = A + B.W$$
  
 $B' = A - B.W$ 

Where W is the complex coefficient and A & B are the complex data

Fig.4 illustrates how a single PDSP16116 may be combined with two PDSP1601's and two PDSP16318's to form a complete BFP butterfly processor. The PDSP16318's are used to perform the complex addition and subtraction of the butterfly operation, while the PDSP1601's are used to match the data path of the A-word to the pipelining and shifting operations within the PDSP16116.

For more information on the theory and construction of this butterfly processor, refer to application note AN59.

#### **BFP MODE OPERATION**

The BFP mode on the PDSP16116 is intended for use in the FFT application described above. i.e. it is intended to prevent data degredation during the course of an FFT calculation. The operation of the PDSP16116 based BFP butterfly processor (see Fig.4) is described below.

#### The Block Floating Point System

A block floating point system is essentially an ordinary integer arithmetic system with some clever logic bolted on. The object of the extra logic is to lend the system some of the enormous dynamic range afforded by a true floating point system without suffering the corresponding loss in performance.

The initial data used by the FFT should all have the same binary arithmetic weighting. i.e. the binary point should occupy the same position in every data word, as is normal in integer arithmetic. However, during the course of the FFT, a variety of weightings are used in the data words to increase the dynamic range available. This situation is similar to that within a true floating point system, though the range of numbers representable is more limited. In the BFP system used in the PDSP16116, there are, within any one pass of the FFT, four possible positions of the binary point wihin the integer words. To record the position of its binary point, each word has a 2-bit word tag associated with it. By way of example, in a particular pass we may have the following four positions of binary point available, each denoted by a certain value of word tag:

XX.XXXXXXXXXXX	word tag = 00
XXX.XXXXXXXXXX	word tag = $01$
XXXX.XXXXXXXXX	word tag = 10
XXXXX.XXXXXXXX	word tag = 11

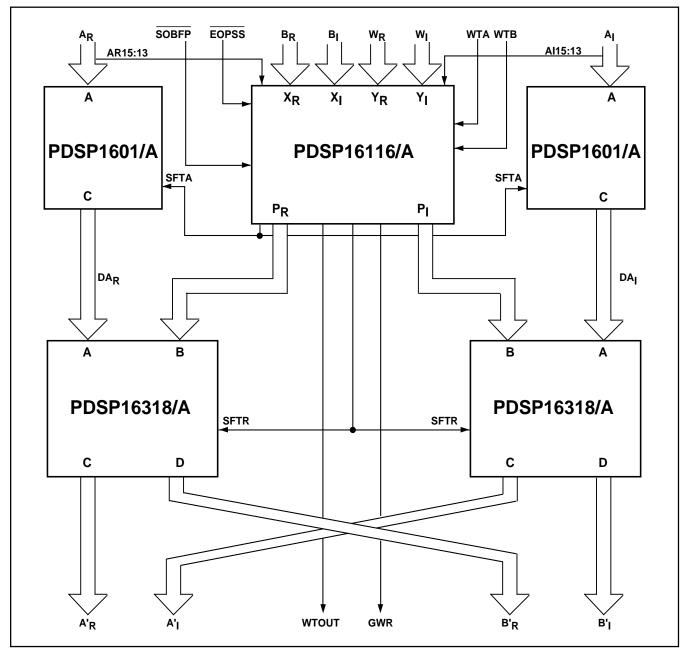


Fig.4 FFT Butterfly Processor

At the end of each constituent pass of the FFT, the positions of the binary point supported may change to reflect the trend of data increase or decreases in magnitude. Hence, in the pass following that of the above example, the four positions of binary point supported may be change to:

This variation in the range of binary points supported from pass to pass (i.e. the movement of the binary point relative to its position in the original data) is recorded in the GWR.

Thus we can determine the position of the binary point relative to its initial position by modifying the value of GWR by WTOUT for a given word as shown in Table 6.

As an example, if GWR=01001 and WTOUT=10 then the binary point has moved 10 places to the right of its original position.

#### The butterfly operation

The butterfly operation is the arithmetic operation which is repeated many times to produce an FFT. The PDSP16116A based butterfly processor performs this operation in a low power high accuracy chip set.

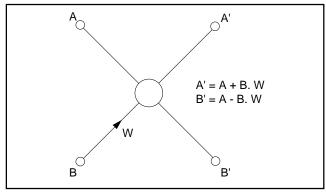


Fig.5 Butterfly Operation

A new butterfly operation is commenced each cycle, requiring a new set of data for , B, W, WTA and WTB. Five cycles later, the corresponding results A' and B' are produced along with their associated WTOUT. In between, the signals SFTA and SFTR are produced and acted upon by the shifters in the PDSP1601/A and PDSP16318/A. The timing of the data and control signals is shown in Fig.6.

The results (A' and B') of each butterfly calculation in a pass must be stored away to be used later as the input data (A and B) in the next pass. Each result must be stored together with its associated word tag, WTOUT. Although WTOUT is common to both A' and B', it must be stored separately with each word as the words are used on different cycles during the next pass. At the inputs, the word tag associated with the A word is known as WTA and the word tag associated with the B word is known as WTB. Hence, the WTOUTs from one pass will become the WTAs and WTBs for the following pass. It should be noted that the first pass is unique in that word tags need not be input into the butterfly as all data initially has the same weighting. Hence, during the first pass alone, the inputs WTA and WTB are ignored.

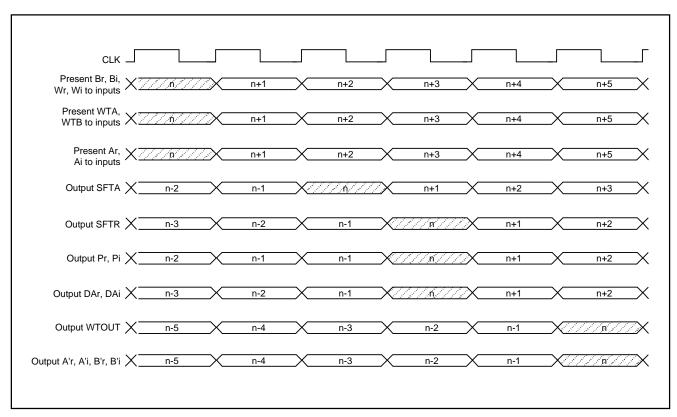


Fig.6 Butterfly Data and Control Signals

# Control of the FFT

To enable the block floating point hardware to keep track of the data, the following signals are provided:

SOBFP - start of the FFT EOPSS - end of current pass

These inform the PDSP16116/A when an FFT is starting and when each pass is complete. Fig.7 shows how these signals should be used and a commentary is provided below.

To commence the FFT, the signal <u>EOPSS</u> should be set high (where it will remain for the duration of the pass). <u>SOBFP</u> should be pulled low during the initial cycle when the first data words A and B are presented to the inputs of the butterfly processor. The following cycle <u>SOBFP</u> must be pulled high

where it should remain for the duration of the FFT. New data is presented to the processor each successive cycle until the end of the first pass of the FFT. On the last cycle of the pass, the signal <u>EOPSS</u> should be pulled low and remain low for a minimum of five cycles\*, the time required to clear the pipeline of the butterfly processor so that all the results from one pass are obtained before commencing the following pass. On the initial cycle of each new pass, the signal <u>EOPSS</u> should be pulled high and it should remain high until the final cycle of that pass, when it is pulled low again.

\* Should a longer pause be required between passes - to arrange the data for the next pass, for example, then <u>EOPSS</u> may be kept low as long as necessary - the next pass cannot commence until it is brought high again.

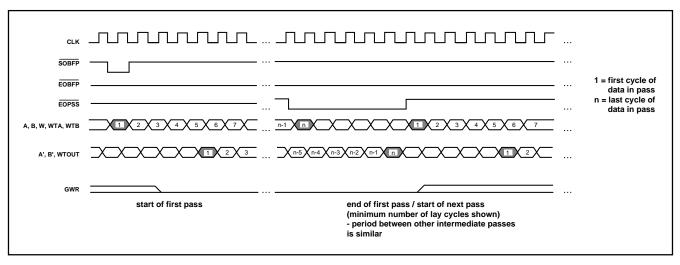


Fig.7 Use of the BFP Control Signals

# **FFT Output Normalisation**

When an FFT system outputs a series of FFT results for display, storage or transmission, it is essential that all results are compatible, i.e. with the binary point in the same position. However, in order to preserve the dynamic range of the data in the FFT calculation, the PDSP1601/A employs a range of different weightings. Therefore, data must be re-formatted at the end of the FFT to be pre-determined common weighting. This can be done by comparing the exponent of given data word with the pre-determined unversial exponent and then shifting the data word by the difference. The PDSP1601/A, with its multifunction 16 bit barrel shifter, is ideally suited to this task.

What value should the Unversal Exponent take? Well, according to theory, the largest possible data result from an FFT is N times the largest input data. This means that the binary point can move a maximum of log2(N) places to the right. Hence, if we choose the Unverisal Exponent to be log2(N) this should give us sufficient range to represent all data points faithfully.

In practice, data output may never approach the theoretical maximum. Hence, it may be worthwhile to try various Unverisal Exponents and choose the one best suited to the particular application.

Data is output from the butterfly processor with a two-part exponent: the 5-bit GWR applicable to all data words from a given FFT and a 2-bit WTOUT associated with each individual data word. To find the complete exponent for a given word, the GWR for that FFT must be modified by its WTOUT as shown in Table 6. The result is the number of places the binary point has shifted to the right during the course of the FFT.

This value must be compared with the Unversial Exponent to determine the shift required. This is done by subtracting it from the Unversial Exponent. The number of places to be shifted is equal to the difference between the two exponents. The shift can be implemented in a PDSP1601/A. The shift value is fed into the SV port.

As FFT data consists of real and imaginary parts, either two PDSP1601As must be used (controlled by the same logic) or a single PDSP1601/A could be used handling real and imaginary data on alternate cycles (using the same instructions for both cycles).

An example of an output normalisation circuit is shown in Fig.8. Only 4 bit data paths are used in calculating the shift. This means that we must be able to trap very small values negative of GWR and force a 15-bit right shift in such cases.

#### N.B.

It is easier to simply add the word tag to the exponent for the purpose of determing the shift required, instead of modifying it according to Table.6. To compensate for this, the Universal Exponent may be increased by one.

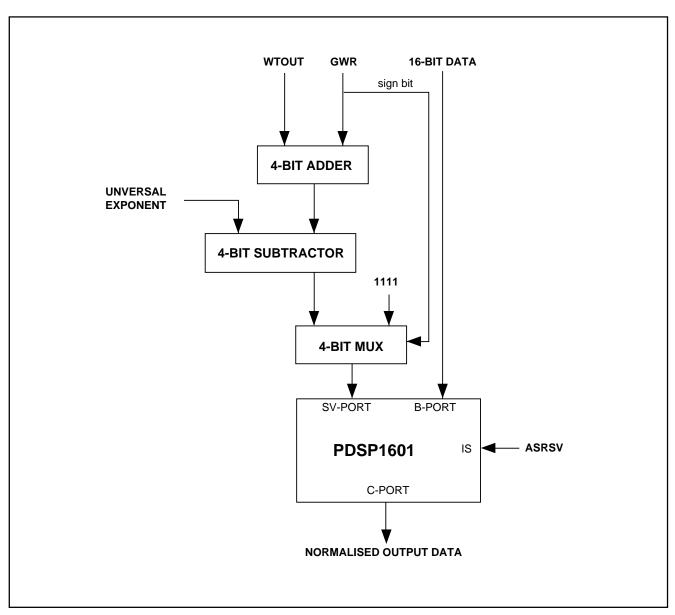


Fig.8 Output Normalisation Circuitry

# **ABSOLUTE MAXIMUM RATINGS (Note 1)**

Supply voltage Vcc	-0.5V to 7.0V				
Input voltage VIN	-0.5V to Vcc +0.5V				
Output voltage Vout	-0.5V to Vcc +0.5V				
Clamp diode current per lk (see note	2) 18mA				
Static discharge voltage (HBM)	500V				
Storage temperature range Ts	-65°C to +150°C				
Ambient temperature with power applied TAMB					
Military	-55°C to +125°C				
Industrial	-40°C to +85°C				
Junction temperature	150°C				
Package power dissipation	1000mW				
Thermal resistances					
Junction to case øuc	12°C/W				
Junction to case ØJA	29°C/W				

#### **NOTES**

- 1. Exceeding these ratings may cause permanent damage. Functional operation under these conditions is not implied.

  2. Maximum dissipation or 1 second should not be exceedeed, only
- one output to be tested at any one time.
- 3. Exposure to absolute maximum ratings for extended periods may affect device reliability.

# **ELECTRICAL CHARACTERISTICS**

Operating conditions (unless otherwise stated): Industrial:  $T_{AMB} = -40$ °C to +85°C,  $Vcc = 5.0V \pm 10$ %, GND = 0VMilitary:  $T_{AMB}$  = -55°C to +125°C, Vcc = 5.0V  $\pm$  10%, GND = 0V

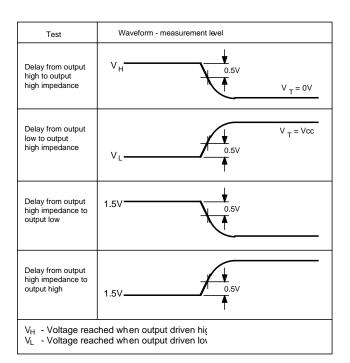
# **Static Characteristics**

Characteristic	Symbol	Value		Value		Value		Units	Conditions
		Min.	Тур.	Min.					
Output high voltage	Vон	2.4		-	V	Iон = 8mA			
Output low voltage	Vol	-		0.4	V	IoL = -8mA			
Input high voltage	Vін	3.0		-	V	CLK input only			
Input high voltage	Vін	2.2		-	V	All other inputs			
Input low voltage	VIL	-		0.8	V	GND <vin<vcc< td=""></vin<vcc<>			
Input leakage current	lin	-10		+10	μΑ				
Input capacitance	Cin		10		pF	GND <vin<vcc< td=""></vin<vcc<>			
Output leakage current	loz	-50		+50	μA	Vcc = Max			
Output S/C current	los	10		300	mA				

# **Switching Characteristics**

Characteristic		PDSP16116		PDSP16116A		Conditions
	Min.	Max.	Min.	Max.	Units	
CLK rising edge to P-PORTS	5	45	5	23	ns	2 x LSTTL + 20pF
CLK rising edge to WTOUT1:0	5	30	5	20	ns	2 x LSTTL + 20pF
CLK rising edge to GWR4:0	5	30	5	20	ns	2 x LSTTL + 20pF
CLK rising edge to SFTA1:0	5	60	5	30	ns	2 x LSTTL + 20pF
CLK rising edge to SFTR2:0	5	50	5	28	ns	2 x LSTTL + 20pF
Setup CEX or CEY to CLK rising edge	11	-	8	-	ns	
Hold CEX or CEY to CLK rising edge	-	0	-	0	ns	
Setup X or Y port inputs to CLK rising edge	11	-	8	-	ns	
Hold X or Y port inputs to CLK rising edge	-	2	-	0	ns	
Setup WTA1:0, WTB1:0, SOBFP or EOPSS inputs to CLK rising edge	14	-	8	-	ns	
Hold WTA1:0, WTB1:0, SOBFP or EOPSS inputs to CLK rising edge	-	0	-	0	ns	
Setup CONX or CONY inputs to CLK rising edge	14	-	8	_	ns	
Hold CONX or CONY inputs to CLK rising edge	-	0	_	Ιo	ns	
Setup AR15:13 or AI15:13 to CLK rising edge	14	-	-	_	ns	
Hold AR15:13 or Al15:13 to CLK rising edge	-	0	-	0	ns	
OPSEL to valid P-PORTS	-	35	-	20	ns	2 x LSTTL + 20pF
OER or OEI rising PR-PORT or PI-PORT high to Z	-	35	-	25	ns	see Fig.9
OER or OEI rising PR-PORT or PI-PORT low to Z		45	-	25	ns	see Fig.9
OER or OEI falling PR-PORT or PI-PORT Z to high	-	22	-	18	ns	see Fig.9
OER or OEI falling PR-PORT or PI-PORT Z to low	- 100	24	-	18	ns	see Fig.9
Clock period		-	50	-	ns	
Clock high time		-	12	-	ns	
Clock low time		-	12	-	ns	
Vcc Current (CMOS input levels)		60	-	80	mΑ	see Note 4
Vcc Current (TTL input levels)	-	100	-	130	mA	see Note 4

NOTE 4 :-  $V_{CC}$  = Max Outputs unloaded, clock freq = Max



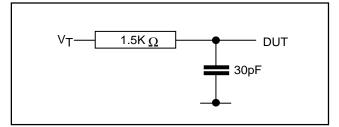


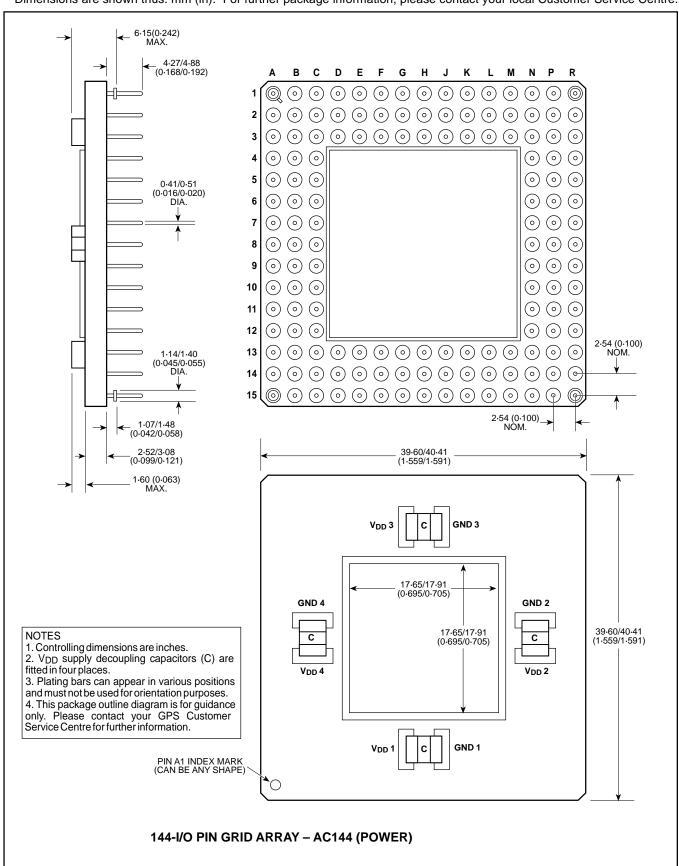
Fig.9 Three state delay measurement load

# **ORDERING INFORMATION**

PDSP16116 MC GGDR	10MHz	MIL-883 screened
PDSP16116A B0 AC	20MHz	Industrial
PDSP16116A A0 AC	20MHz	Military
PDSP16116A B0 GG	20MHz	Industrial
PDSP16116A A0 GG	20MHz	Military
PDSP16116A MC GGDR	20MHz	MIL-883 screened
PDSP16116B B0 AC	25MHz	Industrial

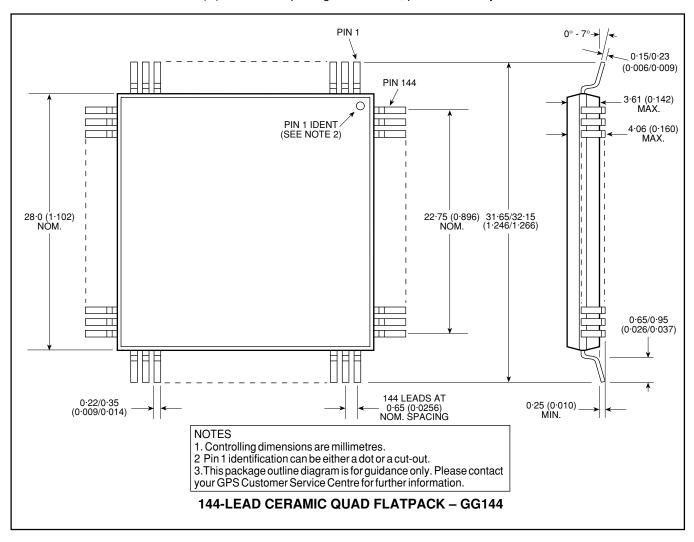
#### **PACKAGE DETAILS**

Dimensions are shown thus: mm (in). For further package information, please contact your local Customer Service Centre.



#### **PACKAGE DETAILS**

Dimensions are shown thus: mm (in). For further package information, please contact your local Customer Service Centre.





HEADQUARTERS OPERATIONS

# **GEC PLESSEY SEMICONDUCTORS**

Cheney Manor, Swindon,

Wiltshire SN2 2QW, United Kingdom.

Tel: (01793) 518000 Fax: (01793) 518411

#### **GEC PLESSEY SEMICONDUCTORS**

P.O. Box 660017

1500 Green Hills Road,

Scotts Valley, California 95067-0017,

United States of America.

Tel: (408) 438 2900

Fax: (408) 438 5576

#### CUSTOMER SERVICE CENTRES

- FRANCE & BENELUX Les Ulis Cedex Tel: (1) 69 18 90 00 Fax: (1) 64 46 06 07
- GERMANY Munich Tel: (089) 3609 06-0 Fax: (089) 3609 06-55
- ITALY Milan Tel: (02) 6607151 Fax: (02) 66040993
- JAPAN Tokyo Tel: (03) 5276-5501 Fax: (03) 5276-5510
- KOREA Seoul Tel: (2) 5668141 Fax: (2) 5697933
- NORTH AMERICA Scotts Valley, USA Tel (408) 438 2900 Fax: (408) 438 7023.
- SOUTH EAST ASIA Singapore Tel: 3827708 Fax: 3828872
- SWEDEN Stockholm Tel: (8) 702 97 70 Fax: (8) 640 47 36
- TAIWAN, ROC Taipei Tel: (2) 5461260. Fax: (2) 7190260
- UK, EIRE, DENMARK, FINLAND & NORWAY

Swindon Tel: (01793) 726666 Fax: (01793) 518582

These are supported by Agents and Distributors in major countries world-wide.

© GEC Plessey Semiconductors 1996 Publication No. DS3707 Issue No. 4.2 October 1996
TECHNICAL DOCUMENTATION - NOT FOR RESALE, PRINTED IN UNITED KINGDOM.

This publication is issued to provide information only which (unless agreed by the Company in writing) may not be used, applied or reproduced for any purpose nor form part of any order or contract nor to be regarded as a representation relating to the products or services concerned. No warranty or guarantee expresses or implied is made regarding the capability, performance or suitability of any product or service. The Company reserves the right to alter without prior notice the specification, design or price of any product or service. Information concerning possible methods of use is provided as a guide only and does not constitute any guarantee that such methods of use will be satisfactory in a specific piece of equipment. It is the user's responsibility to fully determine the performance and suitability of any equipment using such information and to ensure that any publication or data used is up to date and has not been superseded. These products are not suitable for use in any medical products whose faiture to perform may result in significant injury or death to the user. All products and materials are sold and services provided subject to the Company's conditions of sale, which are available on request.