MOTOROLA
SEMICONDUCTOR
TECHNICAL DATA

1M x 32 Bit Dynamic Random Access Memory Module

The MCM32(T)116 is a 32M dynamic random access memory (DRAM) module organized as 1,048,576 x 32 bits. The module is a 72–lead single–in–line memory module (SIMM) consisting of two MCM518160A DRAMs housed in standard 400 mil packages mounted on a substrate along with two 0.22 μF (min) decoupling capacitors mounted adjacent to each DRAM. The MCM518160A is a CMOS high speed dynamic random access memory organized as 1,048,576 sixteen–bit words and fabricated with CMOS silicon–gate process technology.

- Three–State Data Output
- Early–Write Common I/O Capability
- Fast Page Mode Capability
- TTL–Compatible Inputs and Outputs
- RAS–Only Refresh
- CAS Before RAS Refresh
- Hidden Refresh
- 1024 Cycle Refresh: = 16 ms (Max)
- Consists of Two 1M x 16 DRAMs and Four 0.22 μF (Min) Decoupling Capacitors
- Unlatched Data Out at Cycle End Allows Two Dimensional Chip Selection
- Fast Access Time (t_{RAC}): MCM32(T)116–60 = 60 ns (Max)
  MCM32(T)116–70 = 70 ns (Max)
- Low Active Power Dissipation: MCM32(T)116–60 = 2.04 W (Max)
  MCM32(T)116–70 = 1.71 W (Max)
- Low Standby Power Dissipation: TTL Levels = 22 mW (Max)
  CMOS Levels = 11 mW (Max)
- Also Available In Low–Height, Thin TSOP Module (MCM32T116)

MCM32116 MCM32T116
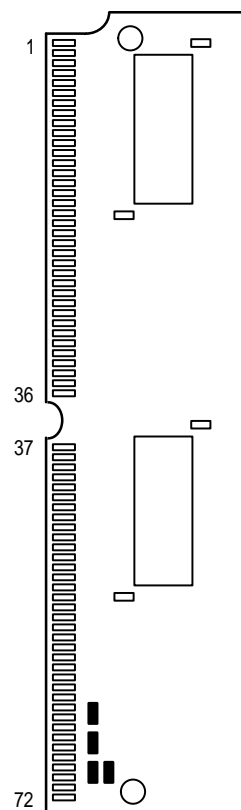
S PACKAGE (SOJ)
SIMM MODULE
CASE 866A–02
(MCM32116)

SH PACKAGE (TSOP)
LOW HEIGHT SIMM MODULE
CASE 866H–01
(MCM32T116)

TOP VIEW

PIN ASSIGNMENTS

Pin	Name	Pin	Name	Pin	Name	Pin	Name	Pin	Name	Pin	Name
1	V_{SS}	13	A1	25	DQ22	37	NC	49	DQ8	61	DQ13
2	DQ0	14	A2	26	DQ7	38	NC	50	DQ24	62	DQ30
3	DQ16	15	A3	27	DQ23	39	V_{SS}	51	DQ9	63	DQ14
4	DQ1	16	A4	28	A7	40	CAS0	52	DQ25	64	DQ31
5	DQ17	17	A5	29	NC	41	CAS2	53	DQ10	65	DQ15
6	DQ2	18	A6	30	V_{CC}	42	CAS3	54	DQ26	66	NC
7	DQ18	19	NC	31	A8	43	CAS1	55	DQ11	67	PD1
8	DQ3	20	DQ4	32	A9	44	RAS0	56	DQ27	68	PD2
9	DQ19	21	DQ20	33	NC	45	NC	57	DQ12	69	PD3
10	V_{CC}	22	DQ5	34	RAS2	46	NC	58	DQ28	70	PD4
11	NC	23	DQ21	35	NC	47	W	59	V_{CC}	71	NC
12	A0	24	DQ6	36	NC	48	NC	60	DQ29	72	V_{SS}

PIN NAMES

A0 – A9	Address Inputs
DQ0 – DQ31	Data Input/Output
CAS0 – CAS3	Column Address Strobe
PD1 – PD4	Presence Detect
RAS0 – RAS3	Row Address Strobe
W	Read/Write Input
V_{CC}	Power (+ 5 V)
V_{SS}	Ground
NC	No Connection

All power supply and ground pins must be connected for proper operation of the device.

MOTOROLA

BLOCK DIAGRAM
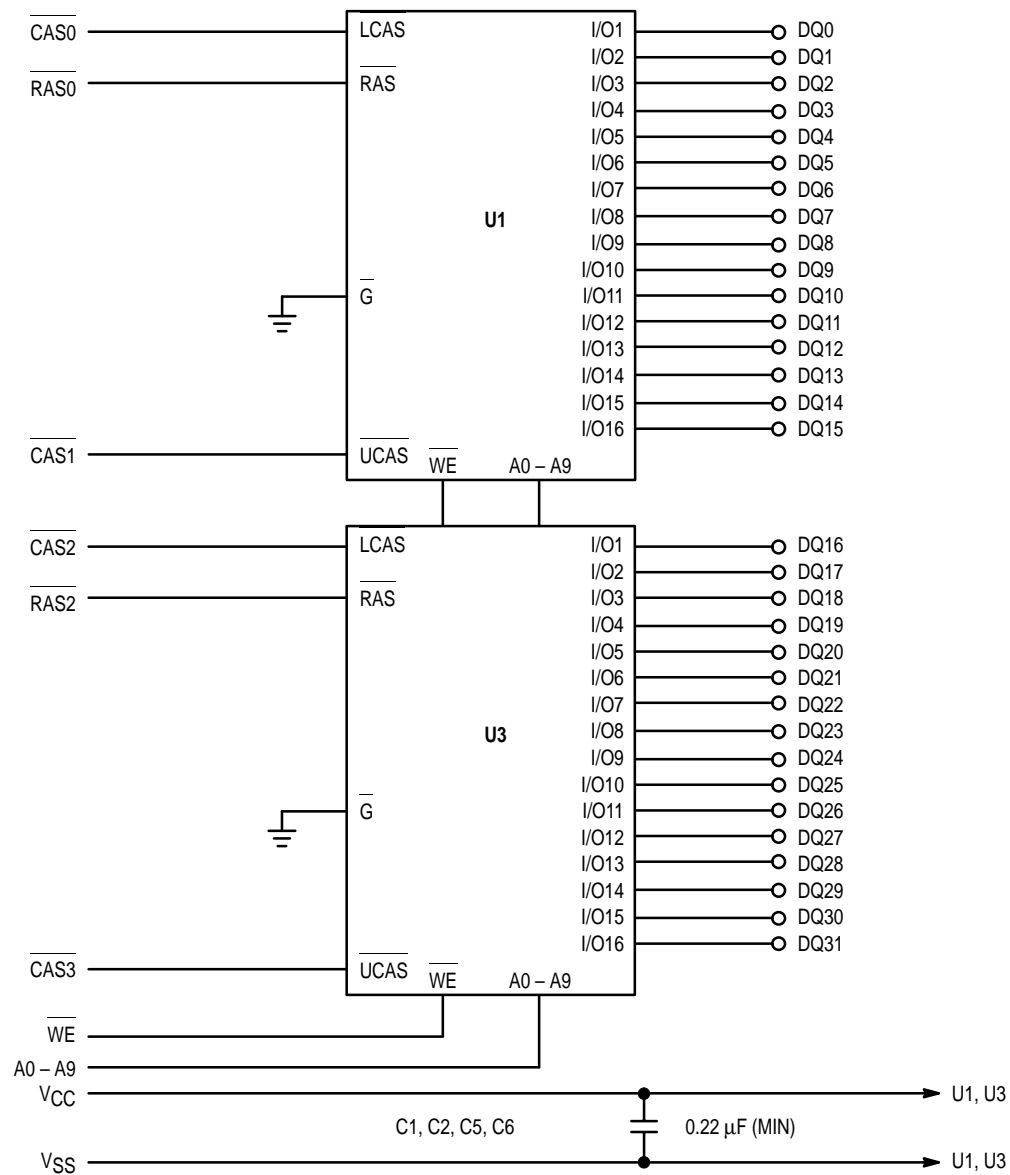
PRESENCE DETECT PIN OUT		
Pin Name	60 ns	70 ns
PD1	V_{SS}	V_{SS}
PD2	V_{SS}	V_{SS}
PD3	NC	V_{SS}
PD4	NC	NC

MCM32116•MCM32T116
2

MOTOROLA DRAM

ABSOLUTE MAXIMUM RATINGS (See Note)

Rating	Symbol	Value	Unit
Power Supply Voltage	V_{CC}	– 0.5 to + 7	V
Voltage Relative to V_{SS} (For Any Pin Except V_{CC})	V_{in}, V_{out}	– 0.5 to V_{CC} + 0.5	V
Data Output Current per DQ Pin	I_{out}	50	mA
Power Dissipation	P_D	2.2	W
Operating Temperature Range	T_A	0 to + 70	°C
Storage Temperature Range	T_{stg}	– 55 to + 125	°C

NOTE: Permanent device damage may occur if ABSOLUTE MAXIMUM RATINGS are exceeded. Functional operation should be restricted to RECOMMENDED OPERATING CONDITIONS. Exposure to higher than recommended voltages for extended periods of time could affect device reliability.

This device contains circuitry to protect the inputs against damage due to high static voltages or electric fields; however, it is advised that normal precautions be taken to avoid application of any voltage higher than maximum rated voltages to these high–impedance circuits.

DC OPERATING CONDITIONS AND CHARACTERISTICS

(V_{CC} = 5.0 V ± 10%, T_A = 0 to 70°C, Unless Otherwise Noted)

RECOMMENDED OPERATING CONDITIONS (All voltages referenced to V_{SS})

Parameter	Symbol	Min	Typ	Max	Unit
Supply Voltage (Operating Voltage Range)	V_{CC}	4.5	5.0	5.5	V
	V_{SS}	0	0	0	
Logic High Voltage, All Inputs	V_{IH}	2.4	—	V_{CC} + 0.5*	V
Logic Low Voltage, All Inputs	V_{IL}	– 0.5**	—	0.8	V

* V_{CC} + 2.0 V at pulse width ≤ 20 ns.
** – 2.0 V at pulse width ≤ 20 ns.

DC CHARACTERISTICS

Characteristic		Symbol	Min	Max	Unit	Notes
V_{CC} Power Supply Current	MCM32(T)116–60, t_{RC} = 110 ns MCM32(T)116–70, t_{RC} = 130 ns	I_{CC1}	— —	370 310	mA	1, 2
V_{CC} Power Supply Current (Standby) (RAS = CAS = V_{IH})		I_{CC2}	—	4	mA	
V_{CC} Power Supply Current During RAS only Refresh Cycles	MCM32(T)116–60, t_{RC} = 110 ns MCM32(T)116–70, t_{RC} = 130 ns	I_{CC3}	— —	370 310	mA	1, 2
V_{CC} Power Supply Current During Fast Page Mode Cycle	MCM32(T)116–60, t_{PC} = 45 ns MCM32(T)116–70, t_{PC} = 45 ns	I_{CC4}	— —	180 160	mA	1, 2
V_{CC} Power Supply Current (Standby) (RAS = CAS = V_{CC} – 0.2 V)		I_{CC5}	—	2	mA	
V_{CC} Power Supply Current During CAS Before RAS Refresh Cycle	MCM32(T)116–60, t_{RC} = 110 ns MCM32(T)116–70, t_{RC} = 130 ns	I_{CC6}	— —	370 310	mA	1
Input Leakage Current ($V_{SS} \leq V_{in} \leq V_{CC}$)		$I_{lkg(I)}$	– 20	+ 20	µA	
Output Leakage Current (CAS at Logic 1, $V_{SS} \leq V_{out} \leq V_{CC}$)		$I_{lkg(O)}$	– 10	10	µA	
Output High Voltage (I_{OH} = – 5 mA)		V_{OH}	2.4	—	V	
Output Low Voltage (I_{OL} = 4.2 mA)		V_{OL}	—	0.4	V	

NOTES:
1. Current is a function of cycle rate and output loading; maximum current is measured at the fastest cycle rate with the output open.
2. Column Address can be changed once or less while RAS = V_{IL} and CAS = V_{IH}.

MOTOROLA DRAM

MCM32116•MCM32T116

CAPACITANCE (f = 1.0 MHz, T_A = 25°C, V_{CC} = 5 V, Periodically Sampled Rather Than 100% Tested)

Characteristic		Symbol	Min	Max	Unit
Input Capacitance	A0 – A9	C_{I1}	—	20	pF
	W	C_{I2}		24	
	RAS0, RAS2	C_{I3}		17	
	CAS0 – CAS3	C_{I4}		17	
I/O Capacitance	DQ0 – DQ31	C_{DQ}	—	17	pF

NOTE: Capacitance measured with a Boonton Meter or effective capacitance calculated from the equation: $C = I\,\Delta t/\Delta V$.

AC OPERATING CONDITIONS AND CHARACTERISTICS
(V_{CC} = 5.0 V ± 10%, T_A = 0 to 70°C, Unless Otherwise Noted)

READ AND WRITE CYCLES (See Notes 1, 2, 3, and 4)

Parameter	Symbol		MCM32(T)116–60		MCM32(T)116–70		Unit	Notes
	Std	Alt	Min	Max	Min	Max		
Random Read or Write Cycle Time	t_{RELREL}	t_{RC}	110	—	130	—	ns	5
Fast Page Mode Cycle Time	t_{CELCEL}	t_{PC}	40	—	45	—	ns	
Access Time from RAS	t_{RELQV}	t_{RAC}	—	60	—	70	ns	6, 7
Access Time from CAS	t_{CELQV}	t_{CAC}	—	15	—	20	ns	6, 8
Access Time from Column Address	t_{AVQV}	t_{AA}	—	30	—	35	ns	6, 9
Access Time from Precharge CAS	t_{CEHQV}	t_{CPA}	—	35	—	40	ns	6
CAS to Output in Low–Z	t_{CELQX}	t_{CLZ}	0	—	0	—	ns	6
Output Buffer and Turn–Off Delay	t_{CEHQZ}	t_{OFF}	0	15	0	15	ns	10
Transition Time (Rise and Fall)	t_T	t_T	3	50	3	50	ns	
RAS Precharge Time	t_{REHREL}	t_{RP}	40	—	50	—	ns	
RAS Pulse Width	t_{RELREH}	t_{RAS}	60	10 k	70	10 k	ns	
RAS Pulse Width (Fast Page Mode)	t_{RELREH}	t_{RASP}	60	100 k	70	200 k	ns	
RAS Hold Time	t_{CELREH}	t_{RSH}	15	—	20	—	ns	
CAS Hold Time	t_{RELCEH}	t_{CSH}	60	—	70	—	ns	
RAS Hold Time from CAS Precharge (Fast Page Mode)	t_{CEHREH}	t_{RHCP}	35	—	40	—	ns	
CAS Pulse Width	t_{CELCEH}	t_{CAS}	15	10 k	20	10 k	ns	
RAS to CAS Delay Time	t_{RELCEL}	t_{RCD}	20	45	20	50	ns	11
RAS to Column Address Delay Time	t_{RELAV}	t_{RAD}	15	30	15	35	ns	12
CAS to RAS Precharge Time	t_{CEHREL}	t_{CRP}	5	—	5	—	ns	
CAS Precharge Time	t_{CEHCEL}	t_{CP}	10	—	10	—	ns	

NOTES: (continued)
1. V_{IH} (min) and V_{IL} (max) are reference levels for measuring timing of input signals. Transition times are measured between V_{IH} and V_{IL}.
2. An initial pause of 200 μs is required after power–up followed by 8 RAS cycles before proper device operation is guaranteed.
3. The transition time specification applies for all input signals. In addition to meeting the transition rate specification, all input signals must transition between V_{IH} and V_{IL} (or between V_{IL} and V_{IH}) in a monotonic manner.
4. AC measurements t_T = 5.0 ns.
5. The specification for t_{RC} (min) is used only to indicate cycle time at which proper operation over the full temperature range (0°C ≤ T_A ≤ 70°C) is ensured.
6. Measured with a current load equivalent to 2 TTL (– 200 μA, + 4 mA) loads and 100 pF with the data output trip points set at V_{OH} = 2.0 V and V_{OL} = 0.8 V.
7. Assumes that t_{RCD} ≤ t_{RCD} (max).
8. Assumes that t_{RCD} ≥ t_{RCD} (max).
9. Assumes that t_{RAD} ≥ t_{RAD} (max).
10. t_{OFF} (max) defines the time at which the output achieves the open circuit condition and is not referenced to output voltage levels.
11. Operation within the t_{RCD} (max) limit ensures that t_{RAC} (max) can be met. t_{RCD} (max) is specified as a reference point only; if t_{RCD} is greater than the specified t_{RCD} (max) limit, then access time is controlled exclusively by t_{CAC}.
12. Operation within the t_{RAD} (max) limit ensures that t_{RAC} (max) can be met. t_{RAD} (max) is specified as a reference point only; if t_{RAD} is greater than the specified t_{RAD} (max), then access time is controlled exclusively by t_{AA}.

| Parameter	Symbol		MCM32(T)116–60		MCM32(T)116–70		Unit	Notes
	Std	**Alt**	**Min**	**Max**	**Min**	**Max**		
Row Address Setup Time	t_{AVREL}	t_{ASR}	0	—	0	—	ns	
Row Address Hold Time	t_{RELAX}	t_{RAH}	10	—	10	—	ns	
Column Address Setup Time	t_{AVCEL}	t_{ASC}	0	—	0	—	ns	
Column Address Hold Time	t_{CELAX}	t_{CAH}	10	—	15	—	ns	
Column Address to RAS Lead Time	t_{AVREH}	t_{RAL}	30	—	35	—	ns	
Read Command Setup Time	t_{WHCEL}	t_{RCS}	0	—	0	—	ns	
Read Command Hold Time Referenced to CAS	t_{CEHWX}	t_{RCH}	0	—	0	—	ns	13
Read Command Hold Time Referenced to RAS	t_{REHWX}	t_{RRH}	0	—	0	—	ns	13
Write Command Hold Time Referenced to CAS	t_{CELWH}	t_{WCH}	10	—	15	—	ns	
Write Command Pulse Width	t_{WLWH}	t_{WP}	10	—	15	—	ns	
Write Command to RAS Lead Time	t_{WLREH}	t_{RWL}	15	—	20	—	ns	
Write Command to CAS Lead Time	t_{WLCEH}	t_{CWL}	15	—	20	—	ns	
Data In Setup Time	t_{DVCEL}	t_{DS}	0	—	0	—	ns	14
Data In Hold Time	t_{CELDX}	t_{DH}	10	—	15	—	ns	14
Refresh Period	t_{RVRV}	t_{RFSH}	—	16	—	16	ms	
Write Command Setup Time	t_{WLCEL}	t_{WCS}	0	—	0	—	ns	15
CAS Setup Time for CAS Before RAS Refresh	t_{RELCEL}	t_{CSR}	5	—	5	—	ns	
CAS Hold Time for CAS Before RAS Refresh	t_{RELCEH}	t_{CHR}	10	—	15	—	ns	
RAS Precharge to CAS Active Time	t_{REHCEL}	t_{RPC}	5	—	5	—	ns	
CAS Precharge Time for CAS Before RAS Counter Time	t_{CEHCEL}	t_{CPT}	20	—	30	—	ns	

NOTES:
13. Either t_{RRH} or t_{RCH} must be satisfied for a read cycle.
14. These parameters are referenced to CAS leading edge in random write cycles.
15. t_{WCS} is not a restrictive operating parameter. It is included in the data sheet as an electrical characteristic only; if $t_{WCS} \geq t_{WCS}$ (min), the cycle is an early write cycle and the data out pin will remain open circuit (high impedance) throughout the entire cycle. If this condition is not satisifed, the condition of the data out (at access time) is indeterminate.
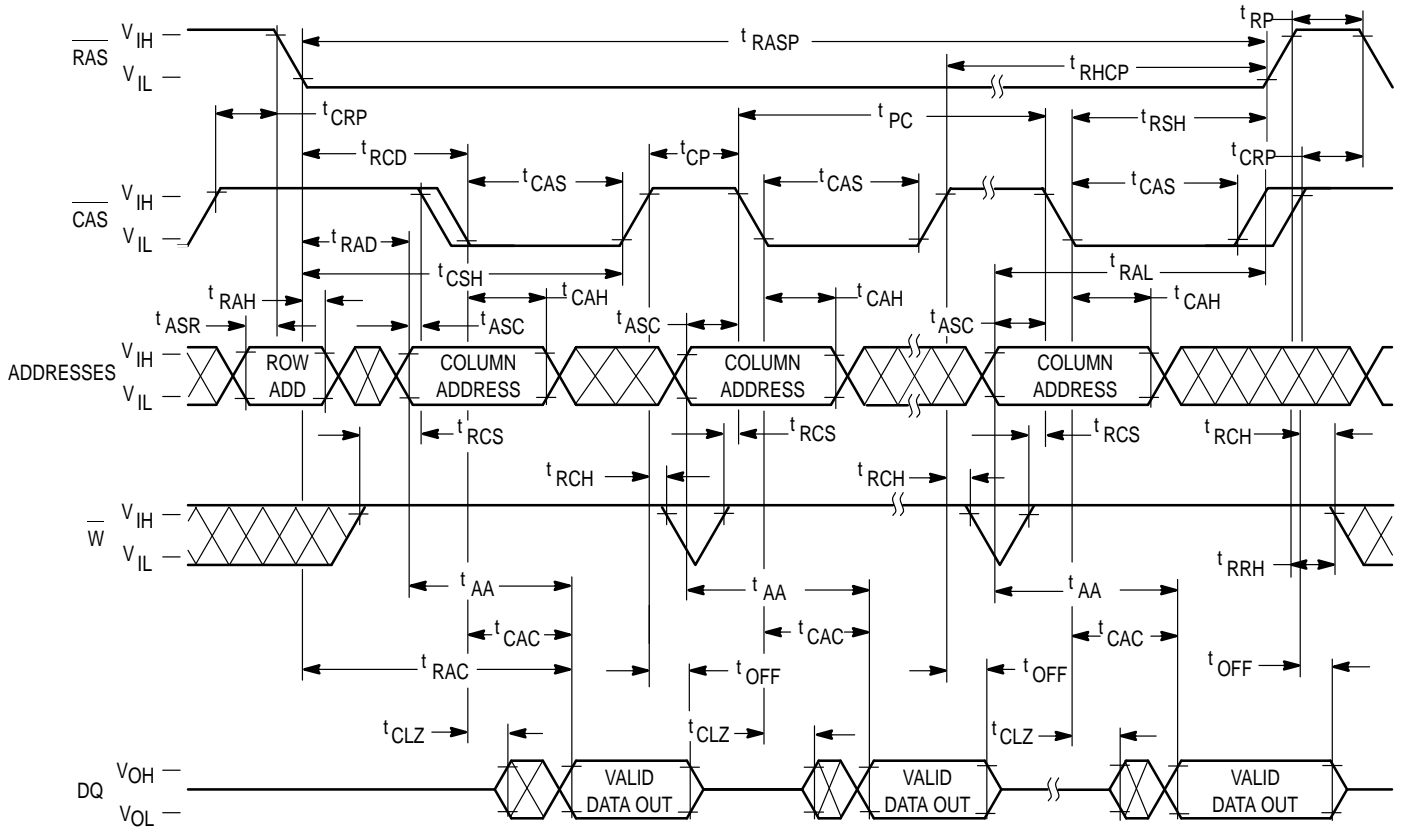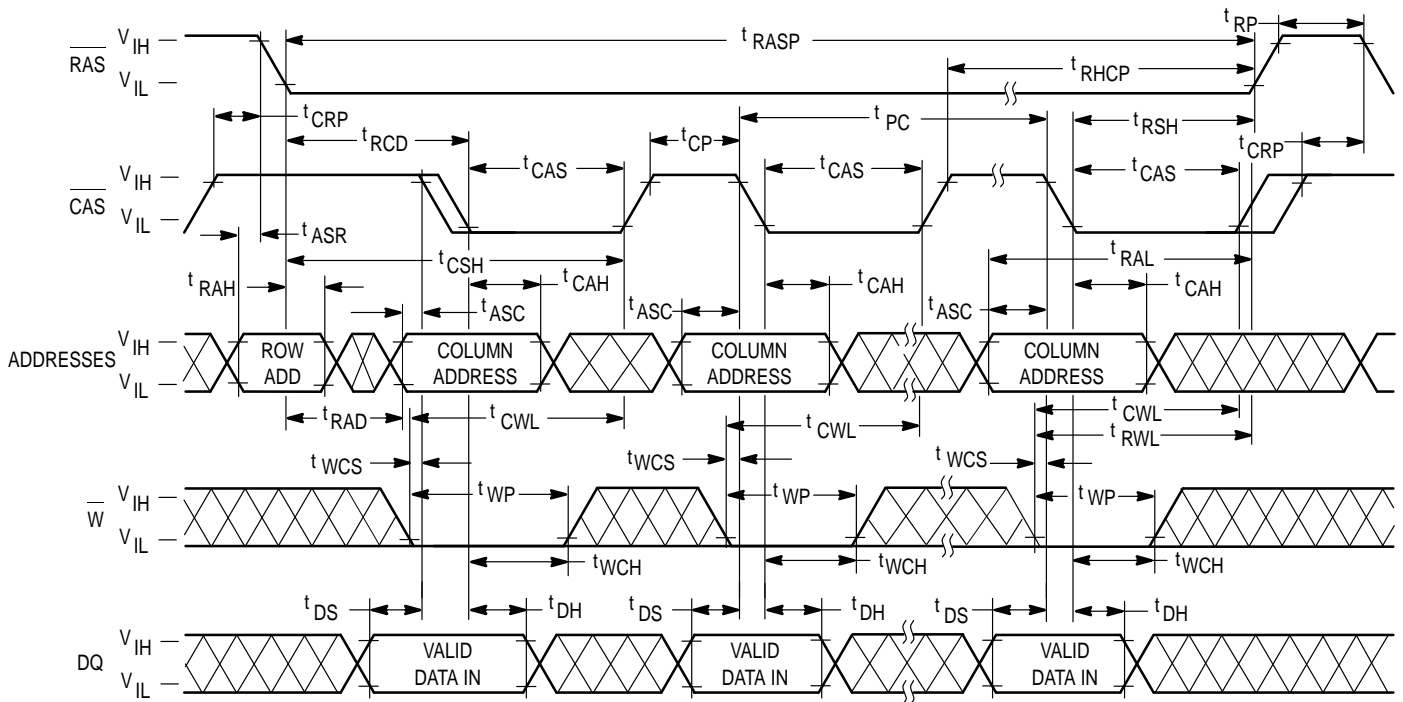
READ CYCLE
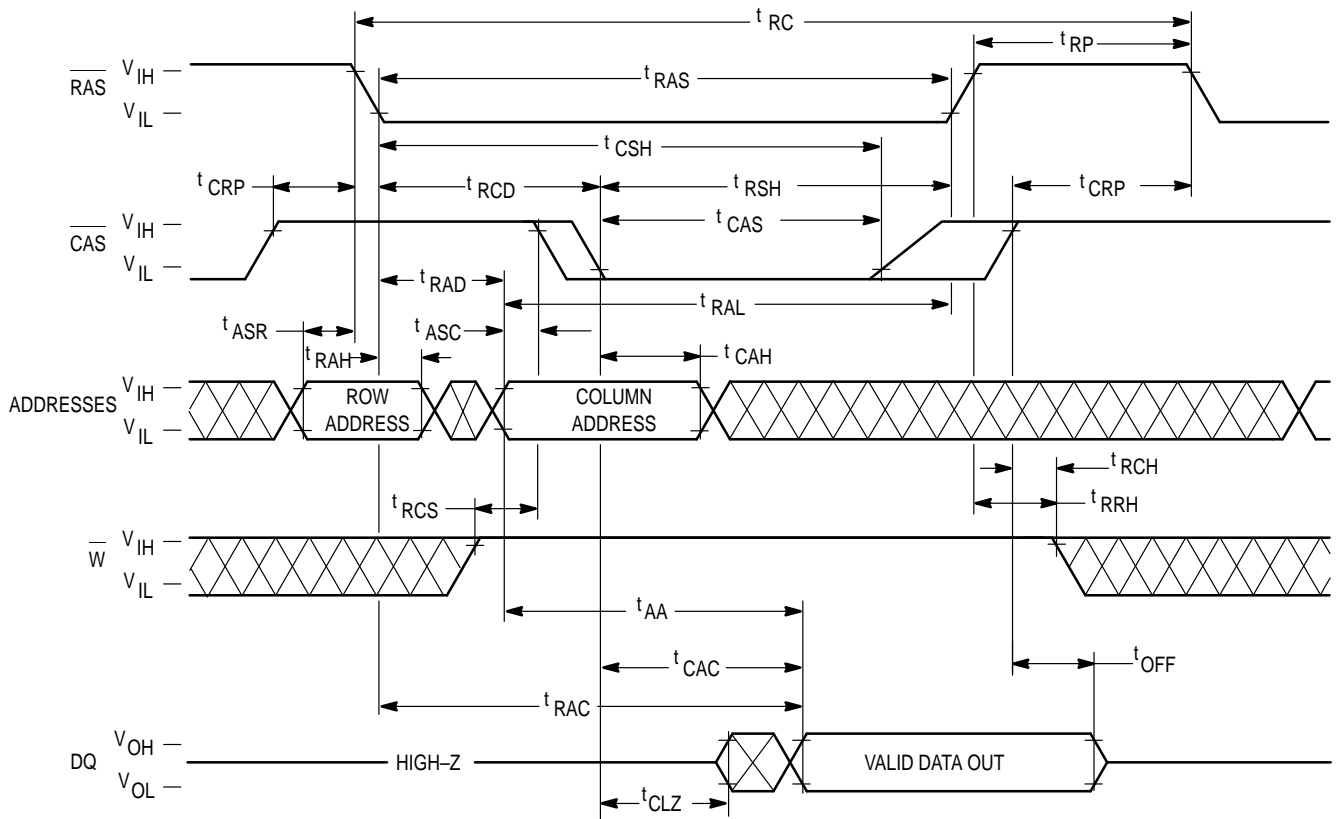
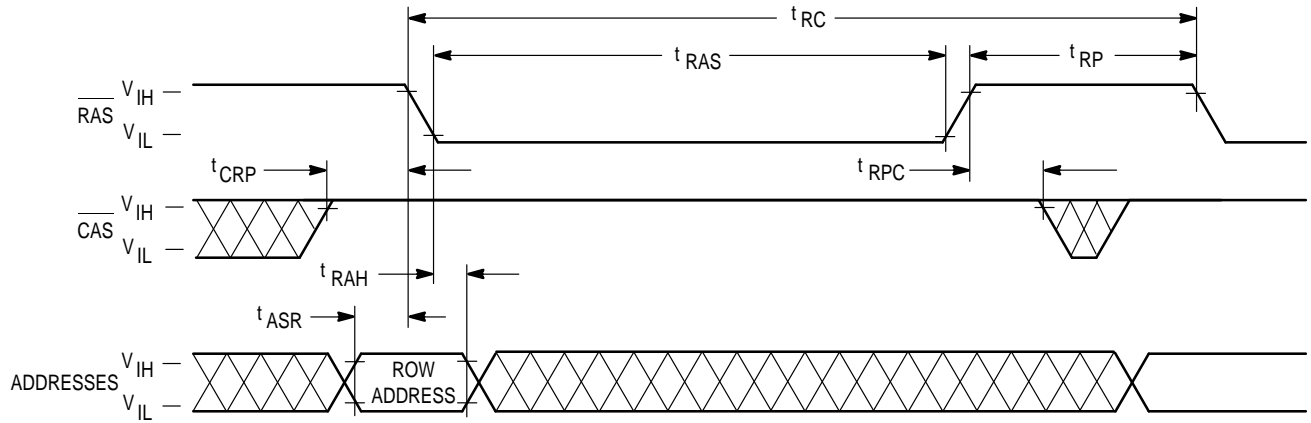

EARLY WRITE CYCLE
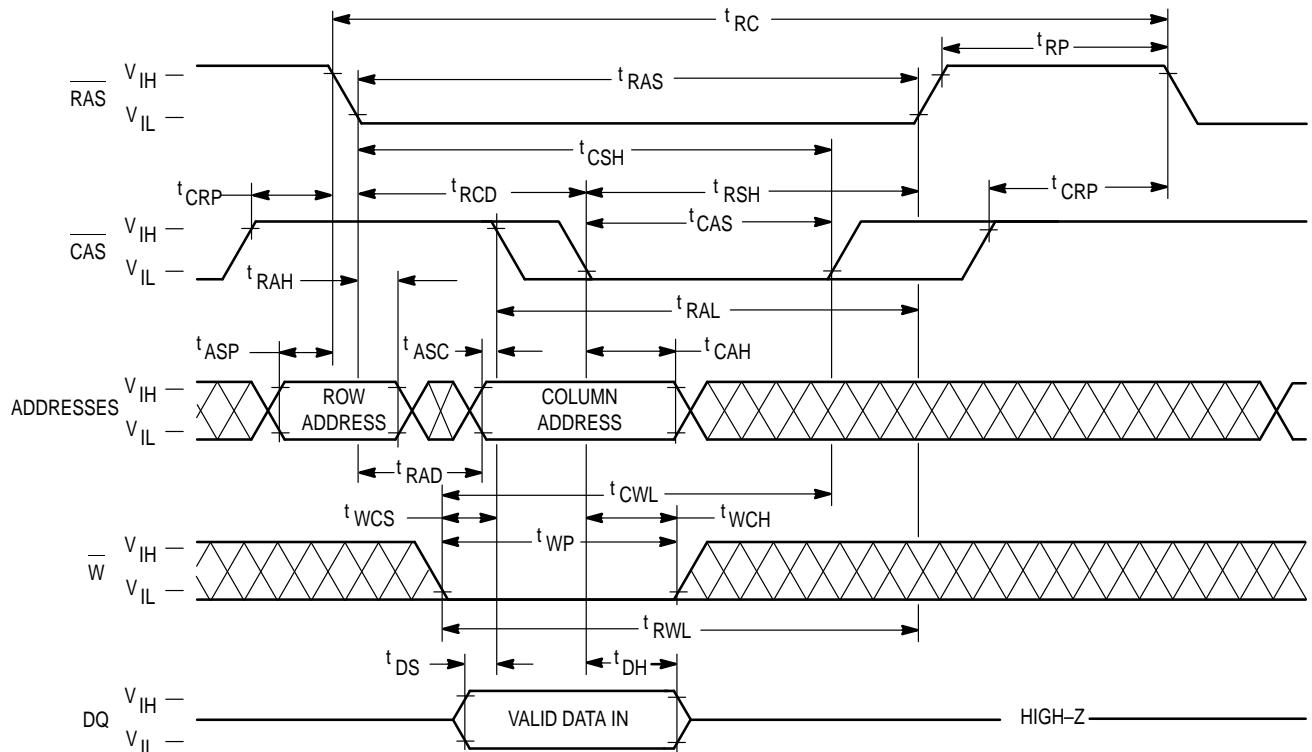
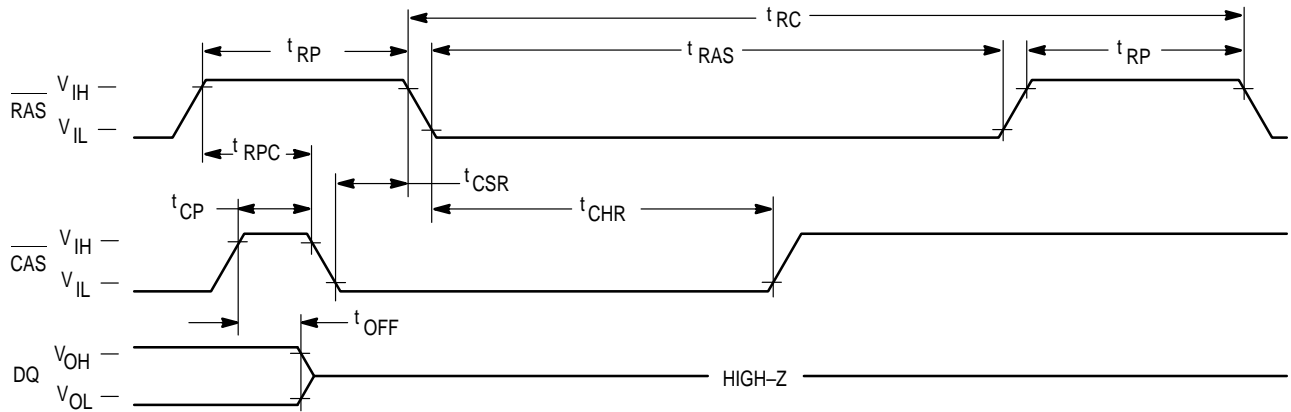## FAST PAGE MODE READ CYCLE

FAST PAGE MODE EARLY WRITE CYCLE

$\overline{RAS}$ ONLY REFRESH CYCLE
(W is "H" or "L")
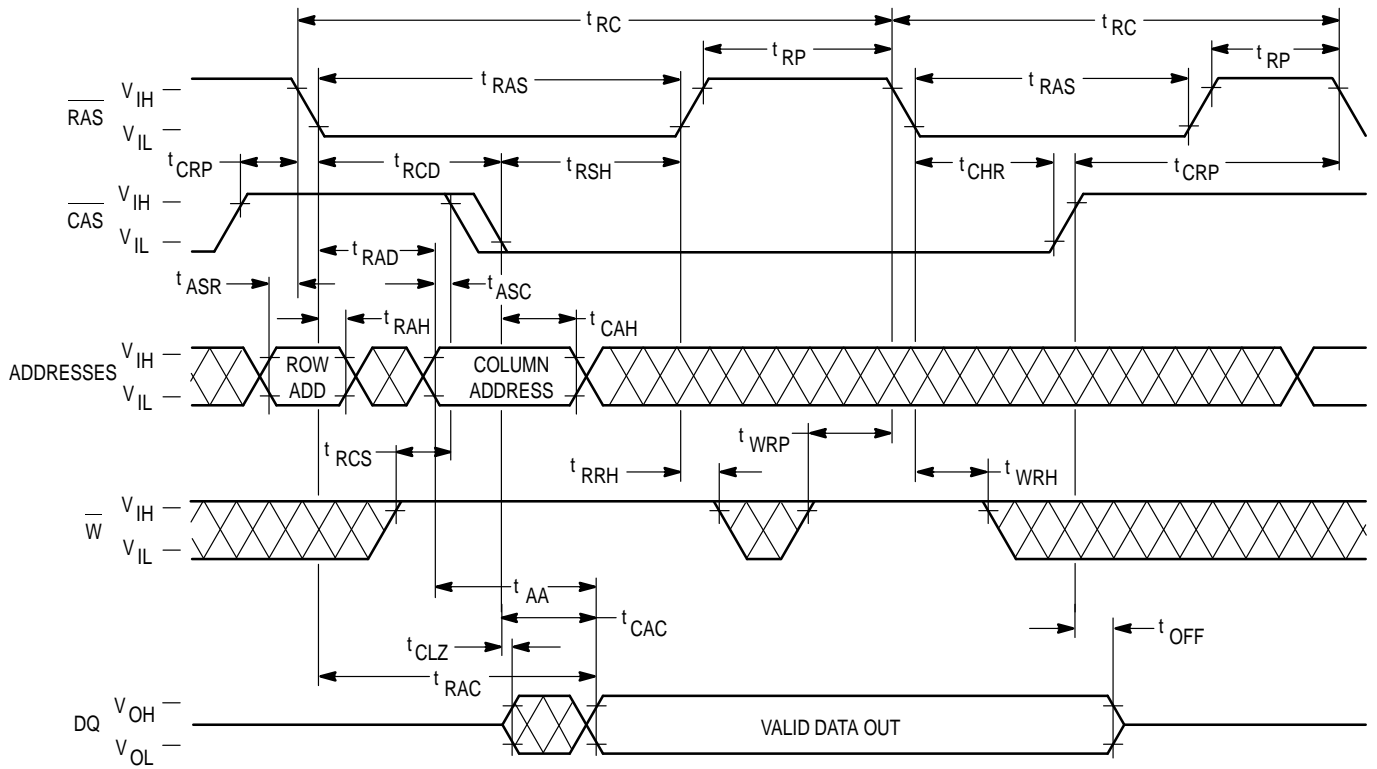

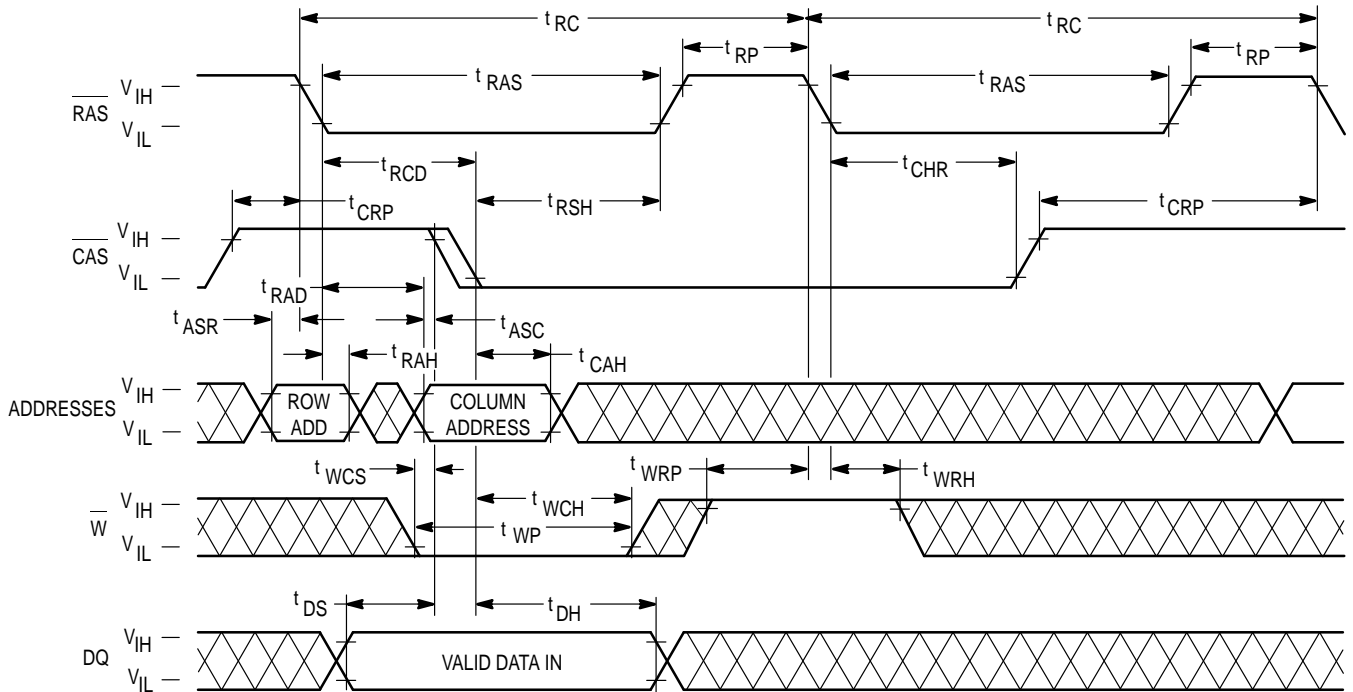## $\overline{CAS}$ BEFORE $\overline{RAS}$ REFRESH CYCLE
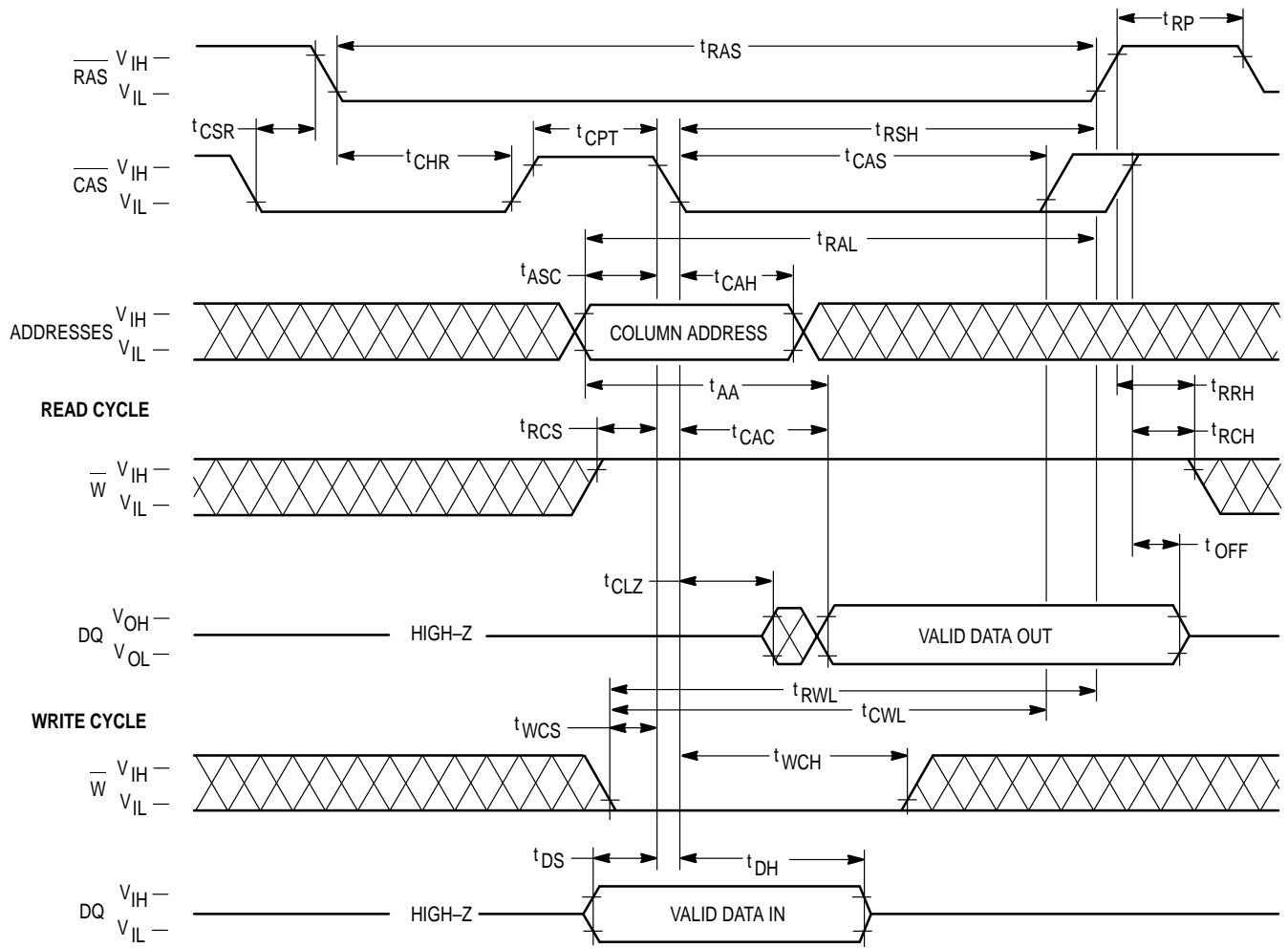### (W, A0 – A9 are "H" or "L")

HIDDEN REFRESH CYCLE (READ)

HIDDEN REFRESH CYCLE (WRITE)

CAS BEFORE RAS REFRESH COUNTER TEST CYCLE

DEVICE INITIALIZATION

On power–up, an initial pause of 200 microseconds is required for the internal substrate generator to establish the correct bias voltage. This must be followed by a minimum of eight active cycles of the row address strobe (clock) to initialize all dynamic nodes within the module. During an extended inactive state (greater than 16 milliseconds with the device powered up), a wake up sequence of eight active cycles is necessary to ensure proper operation.

ADDRESSING THE RAM

The ten address pins on the device are time multiplexed at the beginning of a memory cycle by two clocks, row address strobe (RAS) and column address strobe (CAS), into two separate 10–bit address fields. A total of twenty address bits, ten rows and ten columns, will decode one of the 1,048,576 word locations in the device. RAS active transition is followed by CAS active transition (active = V_{IL}, t_{RCD} minimum) for all read or write cycles. The delay between RAS and CAS active transitions, referred to as the **multiplex window**, gives a system designer flexibility in setting up the external addresses into the RAM.

The external CAS signal is ignored until an internal RAS signal is available. This "gate" feature on the external CAS clock enables the internal CAS line as soon as the row address hold time (t_{RAH}) specification is met (and defines t_{RCD} minimum). The multiplex window can be used to absorb skew delays in switching the address bus from row to column addresses and in generating the CAS clock.

There are three other variations in addressing the module: **RAS only refresh cycle, CAS before RAS refresh cycle,** and **page mode**. All three are discussed in separate sections that follow.

READ CYCLE

The DRAM may be read with either a "normal" random read cycle or a page mode read cycle. The normal read cycle is outlined here, while the page mode cycle is discussed in a separate section.

The normal read cycle begins as described in **ADDRESSING THE RAM**, with RAS and CAS active transitions latching the desired bit location. The write (W) input level must be high (V_{IH}), t_{RCS} (minimum) before the CAS active transition, to enable read mode.

Both the RAS and CAS clocks trigger a sequence of events which are controlled by several delayed internal clocks. The internal clocks are linked in such a manner that the read access time of the device is independent of the address multiplex window. CAS controls read access time: CAS must be active before or at t_{RCD} maximum to guarantee valid data out (DQ) at t_{RAC} (access time from RAS active transition). If the t_{RCD} maximum is exceeded, read access time is determined by the CAS clock active transition (t_{CAC}).

The RAS and CAS clocks must remain active for a minimum time of t_{RAS} and t_{CAS} respectively, to complete the read cycle. W must remain high throughout the cycle, and for time t_{RRH} or t_{RCH} after RAS or CAS inactive transition, respectively, to maintain the data at that bit location. Once RAS transitions to inactive, it must remain inactive for a minimum time of t_{RP} to precharge the internal device circuitry for the next active cycle. DQ is valid, but not latched, as long as the

CAS clock is active. When the CAS clock transitions to inactive, the output will switch to High–Z (three–state) t_{OFF} after the inactive transition.

WRITE CYCLE

The user can write to the DRAM with either an early write or a page mode early write cycle. Early write mode is discussed here, while page mode write operations are covered in a separate section.

A write cycle begins as described in **ADDRESSING THE RAM**. Write mode is enabled by the transition of W to active (V_{IL}). Early write mode is distinguished by the active transition of W, with respect to CAS. Minimum active time t_{RAS} and t_{CAS}, and precharge time t_{RP} apply to write mode, as in the read mode.

An early write cycle is characterized by W active transition at minimum time t_{WCS} before CAS active transition. Data in (DQ) is referenced to CAS in an early write cycle. RAS and CAS clocks must stay active for t_{RWL} and t_{CWL}, respectively, after the start of the early write operation to complete the cycle.

PAGE MODE CYCLES

Page mode allows fast successive data operations at all 1024 column locations on a selected row of the module. Read access time in page mode (t_{CAC}) is typically half the regular RAS clock access time, t_{RAC}. Page mode operation consists of keeping RAS active while toggling CAS between V_{IH} and V_{IL}. The row is latched by RAS active transition, while each CAS active transition allows selection of a new column location on the row.

A page mode cycle is initiated by a normal read or write cycle, as described in prior sections. Once the timing requirements for the first cycle are met, CAS transitions to inactive for minimum t_{CP}, while RAS remains low (V_{IL}). The second CAS active transition while RAS is low initiates the first page mode cycle (t_{PC}). Either a read or write operation can be performed in a page mode cycle, subject to the same conditions as in normal operation (previously described). These operations can be intermixed in consecutive page mode cycles and performed in any order. The maximum number of consecutive page mode cycles is limited by t_{RASP}. Page mode operation is ended when RAS transitions to inactive, coincident with or following CAS inactive transition.

REFRESH CYCLES

The dynamic RAM design is based on capacitor charge storage for each bit in the array. This charge will tend to degrade with time and temperature. Each bit must be periodically **refreshed** (recharged) to maintain the correct bit state. Bits in the MCM32(T)116 require refresh every 16 milliseconds.

This is accomplished by cycling through the 1024 row addresses in sequence within the specified refresh time. All the bits on a row are refreshed simultaneously when the row is addressed. Distributed refresh implies a row refresh every 15.6 microseconds. Burst refresh, a refresh of all 1024 rows consecutively, must be performed every 16 milliseconds.

A normal read or write operation to the RAM will refresh all the bits associated with the particular row decoded. Three other methods of refresh, **RAS–only refresh, CAS before RAS refresh,** and **hidden refresh** are available on this device for greater system flexibility.

RAS–Only Refresh

RAS–only refresh consists of RAS transition to active, latching the row address to be refreshed, while CAS remains high (V_{IH}) throughout the cycle. An external counter is employed to ensure all rows are refreshed within the specified limit.

CAS Before RAS Refresh

CAS before RAS refresh is enabled by bringing CAS active before RAS. This clock order actives an internal refresh counter that generates the row address to be refreshed. External address lines are ignored during the automatic refresh cycle. The output buffer remains at the same state it was in during the previous cycle (hidden refresh).

Hidden Refresh

Hidden refresh allows refresh cycles to occur while maintaining valid data at the output pin. Holding CAS active at the end of a read or write cycle, while RAS cycles inactive for t_{RP} and back to active, starts the hidden refresh. This is essentially the execution of a CAS before RAS refresh from a cycle in progress (see Figure 1). W is subject to the same conditions with respect to RAS active transition (to prevent test mode cycle) as in CAS before RAS refresh.
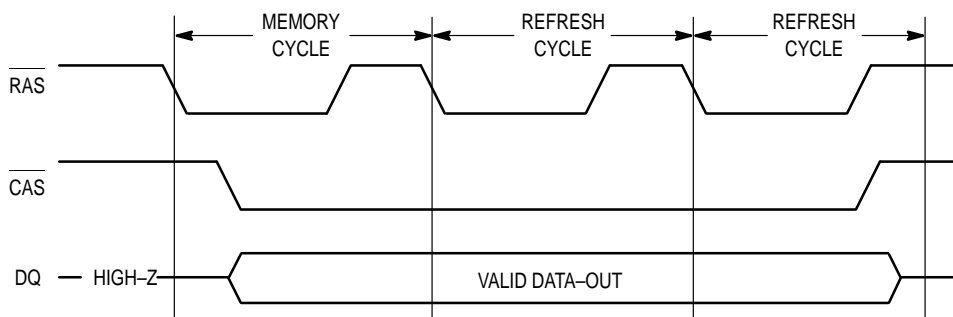
CAS BEFORE RAS REFRESH COUNTER TEST

The internal refresh counter of the device can be tested with a **CAS before RAS refresh counter test**. This refresh counter test is performed with read and write operations. During this test, the internal refresh counter generates the row address, while the external address input supplies the column address. The entire array is refreshed after 1024 test cycles, as indicated by the check data written in each row. See **CAS before RAS refresh counter test cycle** timing diagram.

The test can be performed only after a minimum of **8 CAS before RAS** initialization cycles. The test procedure is as follows:

1. Write 0s into all memory cells (normal write mode).
2. Select a column address, and read 0 out of the cell by performing **CAS before RAS refresh counter test, read cycle**. Repeat this operation 1024 times.
3. Select a column address, and write 1 into the cell by performing **CAS before RAS refresh counter test, write cycle**. Repeat this operation 1024 times.
4. Read 1s (normal read mode), which were written at step three.
5. Repeat steps one through four using complement data.

Figure 1. Hidden Refresh Cycle

ORDERING INFORMATION
(Order by Full Part Number)

MCM — Motorola Memory Prefix

32116 / 32T116 — Part Number

X — Package (SH = Short Height SIMM (TSOP Only), SHG = Short Height Gold Pad SIMM (TSOP Only), S = SIMM (SOJ Only), SG = Gold Pad SIMM (SOJ Only))

XX — Speed (60 = 60 ns, 70 = 70 ns)

Full Part Numbers —
MCM32116S60 MCM32116SG60
MCM32116S70 MCM32116SG70

MCM32T116SH60 MCM32T116SHG60
MCM32T116SH70 MCM32T116SHG70

PACKAGE DIMENSIONS

S PACKAGE
SIMM MODULE (SOJ)
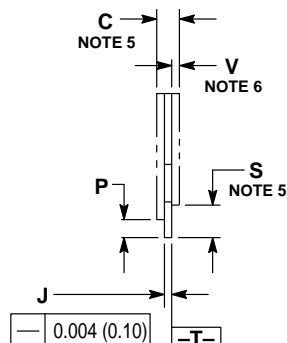**CASE 866A–02**
MCM32116
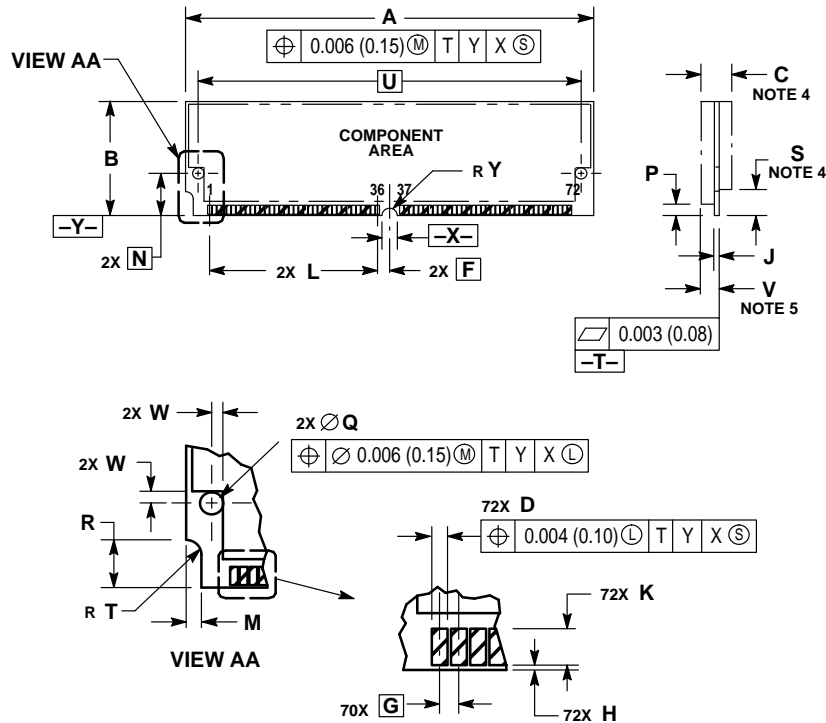
NOTES:
1. DIMENSIONING AND TOLERANCING PER ANSI Y14.5M, 1982.
2. CONTROLLING DIMENSION: INCH.
3. CARD THICKNESS APPLIES ACROSS TABS AND INCLUDES PLATING AND/OR METALIZATION.
4. DIMENSIONS C AND S DEFINE A DOUBLE–SIDED MODULE.
5. DIMENSION V DEFINES OPTIONAL SINGLE–SIDED MODULE.

DIM	INCHES		MILLIMETERS	
	MIN	MAX	MIN	MAX
A	4.245	4.255	107.82	108.08
B	1.245	1.255	31.62	31.88
C	—	0.360	—	9.14
D	0.040	0.042	1.02	1.07
F	0.125 BSC		3.18 BSC	
G	0.050 BSC		1.27 BSC	
H	—	0.010	—	0.25
J	0.047	0.053	1.19	1.35
K	0.100	—	1.19	—
L	1.750 REF		44.45 REF	
M	0.075	0.085	1.90	2.16
N	0.400 BSC		10.16 BSC	
P	0.125	—	3.18	—
Q	0.123	0.127	3.12	3.23
R	0.245	0.255	6.22	6.48
S	0.225	—	5.72	—
T	0.060	0.064	1.52	1.63
U	3.984 BSC		101.19 BSC	
V	—	0.208	—	5.28
W	0.044	—	1.12	—
Y	0.060	0.064	1.52	1.63

PACKAGE DIMENSIONS

SH PACKAGE
LOW HEIGHT SIMM MODULE (TSOP)
CASE 866H–01
MCM32T116

NOTES:
1. DIMENSIONING AND TOLERANCING PER ANSI Y14.5M, 1982.
2. CONTROLLING DIMENSION: INCH.
3. CARD THICKNESS APPLIES ACROSS TABS AND INCLUDES PLATING AND/OR METALIZATION.
4. DIMENSIONS C AND S DEFINE A DOUBLE–SIDED MODULE.
5. DIMENSION V DEFINES OPTIONAL SINGLE–SIDED MODULE.

DIM	INCHES		MILLIMETERS	
	MIN	MAX	MIN	MAX
A	4.245	4.255	107.82	108.08
B	0.995	1.005	25.27	25.53
C	—	0.157	—	4.00
D	0.040	0.042	1.02	1.07
F	0.125 BSC		3.18 BSC	
G	0.050 BSC		1.27 BSC	
H	—	0.010	—	0.25
J	0.047	0.053	1.19	1.35
K	0.100	—	2.54	—
L	1.750 REF		44.45 REF	
M	0.075	0.085	1.91	2.16
N	0.400 BSC		10.16 BSC	
P	0.125	—	3.18	—
Q	0.123	0.127	3.12	3.23
R	0.245	0.255	6.22	6.48
S	0.225	—	5.72	—
T	0.060	0.064	1.52	1.63
U	3.984 BSC		101.19 BSC	
V	—	0.106	—	2.70
W	0.044	—	1.12	—
Y	0.060	0.064	1.52	1.63

MOTOROLA

◊

MCM32116/D